

Soft Eng 3M04
Mid-Term II 2003
Dr. Jacques Carette

Name: _____

Student No.: _____

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- This midterm contains 5 questions on 3 double-sided pages (including this one).
 - This midterm will be marked out of 50. There are 54 total marks available.
 - Answer the question in the space provided.
 - Make sure that your name is on all sheets.
 - You may separate the last page with the MIS/MID and the definition of the mod function.
 - Make sure that you do not get stuck on one question; use your time wisely.
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1. Give the value of all state variables, output and/or exception after each call in the calling sequence for the MIS and MID on the last page. On the first line, write down the complete initial state. Assume that 'a=int. [32]

call	state		output		exceptions	
	MIS	MID	MIS	MID	MIS	MID
isempty()						
* init()						
addfront(1)						
addback(2)						
removefront()						
addback(3)						
addfront(4)						
isempty()						
* addback(5)						
* addback(6)						
removefront()						
removeback()						
removeback()						
removefront()						
isempty()						

2.

$$t \rightarrow ((t = (num \geq 0)) \wedge |s| = num \wedge$$
$$(\forall i : int. 0 \leq i \leq num - 1 \rightarrow s[|s| - i] = d[(start + i) \bmod Max]))$$

Evaluate the abstraction function given above for the starred (*) states indicated in question 1. Please show all relevant details. [8]

init()	
addback(5)	
addback(6)	

3. (true/false) It is possible to write an expression using the state variables of the MID on the last page which will always be equal to the MIS state variable *out*. [1]
4. Explain what your answer to question 3 means with respect to the abstraction function from question 2 from a verification and validation perspective. [2]

5. Given the declarations

```
val s : 'a seq
val f : ('a → bool) → (('a → 'b) → ('a seq → 'b seq) )
val g : 'a → bool
val h : 'a → 'b
```

(a) Write a formula that states that all the elements of s are different. [2]

(b) What is the type of $f(g)$? [2]

(c) $f(h)(g)(s)$ is an invalid expression - explain why. [1]

(d) Write a formula which asserts that all elements of s satisfy the predicate g . [1]

(e) Let $'a = \text{int}$ for this question. Write down a value for s which would make the following formula true. [1]

$$|s| \geq 5 \wedge (\forall i : \text{int}. 0 \leq i < |s| - 1 \rightarrow s[i] < s[i + 1]) \wedge (\exists j : \text{int}. 0 \leq j < |s| \wedge s[j] < -100)$$

(f) Write an expression which would be found in the MIS for f which would assert that $f(g)(h)(s)$ is a sequence of elements y_j of type $'b$, and where $y_j = h(x_i)$ for some $x_i \in s$, and where we select exactly those elements x_i for which $g(x_i)$ is true. [4]

The following MIS and MID will be used for several questions.

Common information:

Used External Functions: NONE

Used External Data Types: 'a

Exported Constants: Max:int = 4

Exported Functions:

Name	Input Types	Output Types	Exceptions
init			
addfront	'a		notinit, full
addback	'a		notinit, full
removefront		'a	notinit, empty
removeback		'a	notinit, empty
isempty		bool	notinit

MIS	MID
<p>State Variables:</p> <p>s : 'a seq</p> <p>t : bool := false</p> <p>out : 'a</p> <p>Transition Functions:</p> <p>init()</p> <p>Transition: s := <></p> <p>t := true</p> <p>addfront(n:'a)</p> <p>Exception: $\neg t \Rightarrow \text{notinit}$</p> <p>$s \geq \text{Max} \Rightarrow \text{full}$</p> <p>Transition: s := s n</p> <p>addback(n:'a)</p> <p>Exception: $\neg t \Rightarrow \text{notinit}$</p> <p>$s \geq \text{Max} \Rightarrow \text{full}$</p> <p>Transition: s := n s</p> <p>'a removefront()</p> <p>Exception: $\neg t \Rightarrow \text{notinit}$</p> <p>$s = 0 \Rightarrow \text{empty}$</p> <p>Transition: out := s[s - 1]</p> <p>s := s[0.. s - 2]</p> <p>Output: out</p> <p>'a removeback()</p> <p>Exception: $\neg t \Rightarrow \text{notinit}$</p> <p>$s = 0 \Rightarrow \text{empty}$</p> <p>Transition: out := s[0]</p> <p>s := s[1.. s - 1]</p> <p>Output: out</p> <p>bool isempty()</p> <p>Exception: $\neg t \Rightarrow \text{notinit}$</p> <p>Output: s = 0</p>	<p>Variables:</p> <p>'a d[0..Max-1]</p> <p>int num := -1, start := 0</p> <p>Exported Functions:</p> <p>init()</p> <p>num := 0;</p> <p>start := 0;</p> <p>addfront(n:'a)</p> <p>if num < 0 then ERROR(notinit)</p> <p>else if num ≥ Max then ERROR(full)</p> <p>else start := (start-1) mod Max;</p> <p>d[start] := n; num := num + 1</p> <p>addback(n:'a)</p> <p>if num < 0 then ERROR(notinit)</p> <p>else if num ≥ Max then ERROR(full)</p> <p>else d[(start+num) mod Max] := n;</p> <p>num := num + 1</p> <p>'a removefront()</p> <p>if num < 0 then ERROR(notinit)</p> <p>else if num = 0 then ERROR(empty)</p> <p>else start := (start+1) mod Max;</p> <p>num := num - 1;</p> <p>RETURN(d[(start-1) mod Max])</p> <p>'a removeback()</p> <p>if num < 0 then ERROR(notinit)</p> <p>else if num = 0 then ERROR(empty)</p> <p>else num := num - 1;</p> <p>RETURN(d[(start+num+1) mod Max])</p> <p>bool isempty()</p> <p>if num < 0 then ERROR(notinit)</p> <p>else RETURN(num=0)</p>

(see back of page for more information)

Convention: if s is of type 'a seq, then $s[a..b]$ with $b < a$ is understood to be $\langle \rangle$.

Reminder:

x	...	-4	-3	-2	-1	0	1	2	3	4	5	6	...
$x \bmod 4$...	0	1	2	3	0	1	2	3	0	1	2	...