## Two squares canonical factorization

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## Introduction

a configuration of two proportional squares $\boldsymbol{u}^{2}$ and $\boldsymbol{v}^{2}$

has been investigated in many different contexts:

- Smyth et al.: investigating three squares with intention to find a position for amortization argument for the runs conjecture a unique factorization of the type

$$
\boldsymbol{u}=\boldsymbol{u}_{1} \boldsymbol{u}_{2} \boldsymbol{u}_{1}, \boldsymbol{v}=\boldsymbol{u}_{1} \boldsymbol{u}_{2} \boldsymbol{u}_{1} \boldsymbol{u}_{1} \boldsymbol{u}_{2} \Longleftrightarrow \frac{3|\boldsymbol{u}|}{2}<|\boldsymbol{v}|<2|\boldsymbol{u}|
$$

- in a computational framework for computations of $\sigma_{d}(n)$ developed by Deza, F., and Jiang: such configurations are used in Liu's PhD thesis to speed up computation of certain values $\sigma_{d}(n)$ in the $(d, n-d)$ table
$\sigma_{d}(n)$ denotes the maximum number of distinct squares in a string of length $n$ with $d$ distinct symbols
- Lam: two rightmost squares $\boldsymbol{u}^{2} \triangleleft \boldsymbol{v}^{2}$ in $\boldsymbol{x}$ have a very particular structure $\boldsymbol{u}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{2}$ and $\boldsymbol{v}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{2} \boldsymbol{u}_{1}{ }^{e_{2}}$ a primitive $\boldsymbol{u}_{1}$ and a non-trivial proper prefix $\boldsymbol{u}_{2}$ of $\boldsymbol{u}_{1}$, and $e_{1} \geq e_{2} \geq 1$.
note that two rightmost squares in $\boldsymbol{x}$ are necessarily proportional
- Deza, F., Thierry: two proportional squares $\boldsymbol{u}^{2} \triangleleft \boldsymbol{v}^{2}$ form a factorizable double square if either $u$ or $v$ is primitive or $u^{2}$ is rightmost in $v^{2}$

A factorizable double square has a unique factorization $\left(\boldsymbol{u}_{\mathbf{1}}, \boldsymbol{u}_{\mathbf{2}}, \boldsymbol{e}_{1}, e_{2}\right): \boldsymbol{u}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{2}$ and $\boldsymbol{v}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{1}{ }^{e_{2}}$
where $\boldsymbol{u}_{1}$ is primitive, $\boldsymbol{u}_{2}$ a non-trivial proper prefix of $\boldsymbol{u}_{1}$, and $e_{1} \geq e_{2} \geq 1$.
moreover, there are some additional constrains to the structure of a factorizable double square
in this contribution we generalize and extend the factorization to all pairs of proportional squares starting at the same position

## Basic notions

$x$ is primitive $\Longleftrightarrow x \neq y^{p}$ for any string $y$ and any integer $p \geq 2$
Ex: $\underline{a a b} a a b$ is not primitive, while aabaaba is
primitive root of $x$ : the smallest $y$ s.t. $x=y^{p}$ for some integer
$p \geq 1$ (is unique and primitive)
$u^{2}$ is primitively rooted $\Longleftrightarrow u$ is a primitive string
$x$ and $y$ are conjugates if $x=u v$ and $y=v u$ for some $u, v$
$x \triangleleft y \Longleftrightarrow x$ is a proper prefix of $y$

A double square $\operatorname{DS}(\boldsymbol{u}, \boldsymbol{v}): \boldsymbol{u}^{2} \triangleleft \boldsymbol{v}^{2}$ and $|\boldsymbol{u}|<|\boldsymbol{v}|<2|\boldsymbol{u}|$.

## Note that in Deza, F., Thierry this would be called a balanced double square

## Lemma (Synchronization Principle)

Given a primitive string $\boldsymbol{x}$, a proper suffix $\boldsymbol{y}$ of $\boldsymbol{x}$, a proper prefix $\boldsymbol{z}$ of $\boldsymbol{x}$, and $m \geq 0$, there are exactly $m$ occurrences of $\boldsymbol{x}$ in $\boldsymbol{y} \boldsymbol{x}^{m} \boldsymbol{z}$.


## Lemma (Common Factor Lemma)

For any strings $x$ and $y$, if a non-trivial power of $x$ and a non-trivial power of $y$ have a common factor of length $|x|+|y|$, then the primitive roots of $x$ and $y$ are conjugates.

In particular, if $x$ and $y$ are primitive, then $x$ and $y$ are conjugates.

Note that both $x$ and $y$ must repeat at least twice
these lemmas are really a folklore, but we included proofs as we did not know of a published proof of the Common Factor Lemma

## A simple corollary of Common Factor Lemma:

## Corollary (Uniqueness Lemma)

$\boldsymbol{x}$ and $\boldsymbol{y}$ primitive strings, $p, q \geq 1$ :
(a) $\boldsymbol{x}^{p}=\boldsymbol{y}^{q} \Rightarrow \boldsymbol{x}=\boldsymbol{y} \& p=q$
(b) $p, q \geq 2, \quad x_{1} \triangleleft \boldsymbol{x}, \quad \boldsymbol{y}_{\mathbf{1}} \triangleleft \boldsymbol{y}$

$$
\boldsymbol{x}^{p} \boldsymbol{x}_{1}=\boldsymbol{y}^{q} \boldsymbol{y}_{1} \Rightarrow \boldsymbol{x}=\boldsymbol{y} \& \boldsymbol{x}_{1}=\boldsymbol{y}_{1} \& p=q
$$

(a) $x^{p}=y^{q}$

- $p=1$
then $\boldsymbol{x}=\boldsymbol{y}^{q}, \boldsymbol{x}$ primitive $\Rightarrow q=1$ and $\boldsymbol{x}=\boldsymbol{y}$
- $p, q \geq 2$
$\boldsymbol{x}^{p}$ and $\boldsymbol{y}^{q}$ have a common factor of length $\geq|\boldsymbol{x}|+|\boldsymbol{y}|$, by the Common Factor Lemma $\boldsymbol{x} \sim \boldsymbol{y}$, hence $\boldsymbol{x}=\boldsymbol{y}$
(b) $\boldsymbol{x}^{p} \boldsymbol{x}_{1}=\boldsymbol{y}^{q} \boldsymbol{y}_{1}, p, q \geq 2$
$\boldsymbol{x}^{p} \boldsymbol{x}_{1}=\boldsymbol{y}^{q} \boldsymbol{y}_{1}$ have a common factor of length $|\boldsymbol{x}|+|\boldsymbol{y}|$, hence $x=y$
the requirement $p, q \geq 2$ is essential - for instance:

$$
\begin{aligned}
& \boldsymbol{x}=a a b b, \boldsymbol{x}_{1}=a a \text { and } p=2: \\
& \boldsymbol{x}^{2} \boldsymbol{x}_{\mathbf{1}}=a a b b a a b b a a \\
& \boldsymbol{y}=a a b b a a b b a, \boldsymbol{y}_{\mathbf{1}}=a \text { and } q=1: \\
& \boldsymbol{y}^{1} \boldsymbol{y}_{\mathbf{1}}=\text { aabbaabbaa } \\
& \boldsymbol{x}^{2} \boldsymbol{x}_{1}=\boldsymbol{y}^{1} \boldsymbol{y}_{1}
\end{aligned}
$$

## Main results

## Lemma (Two Squares Factorization Lemma)

$\forall \quad \operatorname{DS}(\boldsymbol{u}, \boldsymbol{v}) \exists$ unique $\boldsymbol{u}_{1}, \boldsymbol{u}_{2}, e_{1}, e_{2}$ such that
$\boldsymbol{u}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}}$ and $\boldsymbol{v}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{\mathbf{1}}{ }^{e_{2}}$
$u_{1}$ primitive
$\boldsymbol{u}_{\mathbf{2}}$ is a possibly trivial proper prefix of $\boldsymbol{u}_{1}$
$e_{1} \geq e_{2} \geq 1$
Moreover,
(a) $\left|\boldsymbol{u}_{\mathbf{2}}\right|=0 \Rightarrow e_{1}>e_{2} \geq 1$
(b) $\left|\boldsymbol{u}_{\mathbf{2}}\right|>0 \Rightarrow \boldsymbol{v}$ primitive
(c) $\left|\boldsymbol{u}_{2}\right|>0 \& e_{1} \geq 2 \Rightarrow \boldsymbol{u}$ primitive


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$\exists k \geq 1$ s.t. $\boldsymbol{u}=\boldsymbol{z}^{k} \boldsymbol{z}^{\prime}$ for some proper prefix $\boldsymbol{z}^{\prime}$ of $\boldsymbol{z}$
$\boldsymbol{u}_{1}$ primitive root of $\boldsymbol{z}$
$\boldsymbol{z}=\boldsymbol{u}_{1}{ }^{e_{2}}$ for some $e_{2} \geq 1$
for some $e_{1} \geq e_{2} k$ and some prefix $\boldsymbol{u}_{2}$ of $\boldsymbol{u}_{1}$,
$\boldsymbol{u}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}}$ and $\boldsymbol{v}=\boldsymbol{u} \boldsymbol{z}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{\mathbf{1}}{ }^{{ }^{e_{2}}}$
now the uniqueness and other properties:
(i) $\left|\boldsymbol{u}_{\mathbf{2}}\right|=0$
$\boldsymbol{u}=\boldsymbol{u}_{1}{ }^{e_{1}}$ and $\boldsymbol{v}=\boldsymbol{u}_{1}{ }^{e_{1}+e_{2}} \Rightarrow \boldsymbol{v}=\boldsymbol{u}_{1}{ }^{2\left(e_{1}+e_{2}\right)}$
$|\boldsymbol{v}|<2|\boldsymbol{u}|$ and $e_{1} \geq e_{2} \quad \Rightarrow \quad e_{1}>e_{2}$
uniqueness of $\boldsymbol{u}_{1}$ is a consequence of Uniqueness
Lemma (a)
(ii) $\left|\boldsymbol{u}_{\mathbf{2}}\right|>0$
assume ( $\left.\boldsymbol{w}_{\mathbf{1}}, \boldsymbol{w}_{\mathbf{2}}, f_{1}, f_{2}\right)$

$$
\boldsymbol{u}=\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}}=\boldsymbol{w}_{\mathbf{1}}^{f_{1}} \boldsymbol{w}_{2} \& \boldsymbol{v}=\boldsymbol{u}_{1} e_{1} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{1}{ }^{e_{2}}=\boldsymbol{w}_{1}{ }_{1}^{f_{1}} \boldsymbol{w}_{\mathbf{2}} \boldsymbol{w}_{\mathbf{1}}^{f_{2}}
$$

$e_{1}=f_{1}=1 \Rightarrow \boldsymbol{v}=\boldsymbol{u} \boldsymbol{u}_{1}=\boldsymbol{u} \mathbf{w}_{1} \Rightarrow \boldsymbol{u}=\boldsymbol{v}$
WLOG assume that $f_{1}>e_{1}=1$
$\boldsymbol{u}=\boldsymbol{u}_{1} \boldsymbol{u}_{\mathbf{2}}=\boldsymbol{w}_{\mathbf{1}}{ }^{f_{1}} \boldsymbol{w}_{\mathbf{2}} \quad \& \quad \boldsymbol{v}=\boldsymbol{u}_{\mathbf{1}} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{\mathbf{1}}=\boldsymbol{w}_{\mathbf{1}}{ }^{f_{1}} \boldsymbol{w}_{\mathbf{2}} \boldsymbol{w}_{\mathbf{1}}{ }^{\mathrm{f}_{2}} \Rightarrow$ $\boldsymbol{u}_{1}=\boldsymbol{w}_{1}{ }^{\mathrm{f}_{2}}$
$\boldsymbol{u}_{1}$ primitive forces $f_{2}=1$ and $\boldsymbol{u}_{1}=\boldsymbol{w}_{1}$
$\boldsymbol{u}_{\mathbf{1}} \boldsymbol{u}_{\mathbf{2}}=\boldsymbol{w}_{\mathbf{1}}{ }^{f_{1}} \boldsymbol{W}_{\mathbf{2}}=\boldsymbol{u}_{\mathbf{1}}{ }^{f_{1}} \boldsymbol{W}_{\mathbf{2}}$, implies that $f_{1}=\mathbf{1}$
a contradiction
show that $\boldsymbol{v}$ is primitive
suppose the contrary: $\boldsymbol{v}=\boldsymbol{w}^{k}, k \geq 2$
$|\boldsymbol{W}| \leq \frac{|\boldsymbol{V}|}{2} \leq\left|\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}}\right|+\left|\boldsymbol{u}_{\mathbf{2}}\right|$
$\boldsymbol{w}^{2 k}=\boldsymbol{v}^{2}=\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}+e_{2}} \boldsymbol{u}_{\mathbf{2}} \boldsymbol{u}_{\mathbf{1}}{ }^{e_{2}}$
$\boldsymbol{w}^{2 k}$ and $\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}+e_{2}} \boldsymbol{u}_{\mathbf{2}}$ have a common factor $\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}+e_{2}} \boldsymbol{u}_{\mathbf{2}}$ of length
$\left(\left|\boldsymbol{u}_{1} e_{1}\right|+\left|\boldsymbol{u}_{\mathbf{2}}\right|\right)+\left|\boldsymbol{u}_{1}{ }^{e_{2}}\right| \geq|\boldsymbol{w}|+\left|\boldsymbol{u}_{\mathbf{1}}\right|$
apply Common Factor Lemma to conclude that $\boldsymbol{w} \sim \boldsymbol{u}_{1}$, thus $\boldsymbol{w}=\boldsymbol{u}_{\mathbf{1}}$
primitive string $\boldsymbol{u}_{1}=\boldsymbol{u}_{2} \overline{\boldsymbol{u}}_{2}$ aligns with $\boldsymbol{u}_{2} \boldsymbol{u}_{\mathbf{1}}$, and so $\overline{\boldsymbol{u}}_{2}$ is a prefix of $\boldsymbol{u}_{1}$, in contradiction to Synchronization Principle
let $e_{2} \geq 2$ show that $u$ is primitive
suppose the contrary: $\boldsymbol{u}=\boldsymbol{w}^{k}, k \geq 2$
Hence $|\boldsymbol{w}| \leq \frac{|\boldsymbol{U}|}{2}=\frac{\left(\left|\boldsymbol{u}_{1}{ }^{e_{1}}\right|+\left|\boldsymbol{U}_{\mathbf{2}}\right|\right)}{2}<\left|\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}-1}\right|+\left|\boldsymbol{u}_{\mathbf{2}}\right|$
$e_{2} \geq 1$ and $e_{2} \geq 2 \Rightarrow e_{1}+e_{2} \geq 3$
$\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}} \triangleleft \boldsymbol{u}^{2}=\boldsymbol{w}^{2 k}$
so $\boldsymbol{w}^{2 k}$ and $\boldsymbol{u}_{1}{ }^{e_{1}+e_{2}}$ have a common factor $\boldsymbol{u}_{1}{ }^{e_{1}} \boldsymbol{u}_{2}$
since $\left|\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}}\right| \geq|\boldsymbol{v}|+\left|\boldsymbol{u}_{\mathbf{1}}\right|$, applying Common Factor Lemma, $\boldsymbol{u}_{1}=\boldsymbol{w}$
this in turn implies $\boldsymbol{u}=\boldsymbol{u}_{\mathbf{1}}{ }^{e_{1}} \boldsymbol{u}_{\mathbf{2}}=\boldsymbol{u}_{\mathbf{1}}{ }^{k}$, impossible since $0<\left|\boldsymbol{u}_{\mathbf{2}}\right|<\left|\boldsymbol{u}_{\mathbf{1}}\right|$
observations:
$\left|\boldsymbol{u}_{\mathbf{2}}\right|>0$ if any one of the following conditions holds:
(a) $\boldsymbol{v}$ is primitive
(b) $\boldsymbol{u}$ is primitive
(c) $\boldsymbol{u}^{2}$ is rightmost in $\boldsymbol{v}^{2}$
moreover:
(d) $\left|\boldsymbol{u}_{\mathbf{2}}\right|>0 \Longleftrightarrow \boldsymbol{v}$ is primitive

## Applications

we concluded the paper with a comment and a sketch of how the canonical factorization could be applied to New Periodicity Lemma:
Lemma (2006, Fan, Puglisi, Smyth, and Turpin)
Let $\boldsymbol{x}=\mathrm{DS}(\boldsymbol{u}, \boldsymbol{v})$, where we require that $\boldsymbol{u}^{2}$ be regular and that $\boldsymbol{v}$ be primitive. There is no square $\boldsymbol{w}^{2}$ starting at position $i$, $1 \leq i<|\boldsymbol{v}|-|\boldsymbol{u}|$ with $|\boldsymbol{v}|-|\boldsymbol{u}|<|\boldsymbol{w}|<|\boldsymbol{v}|$ except possibly $|\boldsymbol{w}|=|\boldsymbol{u}|$.

## $v^{2}$ primitive, $u^{2}$ regular



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we can report that since the final submission to PSC2014, we were able to prove using the canonical factorization an extended NPL:

## Theorem

Consider a double square $\operatorname{DS}(\boldsymbol{u}, \boldsymbol{v})$ and let $\boldsymbol{u}^{\prime}$ be a suffix of $\boldsymbol{u}$ so that $\boldsymbol{v}=\boldsymbol{u} \boldsymbol{u}^{\prime}$. Let $\boldsymbol{w}^{2}$ be any square that is a factor of $\boldsymbol{v}^{2}$. Then exactly one of the following mutually exclusive cases holds:
(a) $\boldsymbol{w}=\boldsymbol{v}$, or
(b) $|\boldsymbol{w}|<|\boldsymbol{u}|$, or
(c) $|\boldsymbol{u}| \leq|\boldsymbol{w}|<|\boldsymbol{v}|$ and the primitive root of $\boldsymbol{w}$ is a conjugate of the primitive root of $\boldsymbol{u}^{\prime}$.


## Lemma (Crochemore-Rytter (1995), Fraenkel-Simpson (1998))

Let $\boldsymbol{u}^{2} \triangleleft \boldsymbol{v}^{2} \triangleleft \boldsymbol{w}^{2}$ and let $\boldsymbol{u}$ be primitive, then $|\boldsymbol{u}|+|\boldsymbol{v}| \leq|\boldsymbol{w}|$.
we can also report that since the final submission to PSC2014, we were able to prove using the canonical factorization a generalization of the above lemma:

## Theorem (Bai, Deza, and F.)

Let $\boldsymbol{u}^{2} \triangleleft \boldsymbol{v}^{2} \triangleleft \boldsymbol{w}^{2}$. Then either
(a) $|\boldsymbol{u}|+|\boldsymbol{v}| \leq|\boldsymbol{w}|$
or
(inclusive or)
(b) $\mathbf{u}, \boldsymbol{v}$, and $\boldsymbol{w}$ have the same primitive root

## $\mathcal{T H A N K} \mathcal{Y O U}$

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