Double-Pushout Graph Rewriting

Summary

- Applications of Graphs
- Graph rewriting
- Definition of a pushout
- Double-Pushout and rules
- Example
- Tools

Applications of graphs

- Visual Modeling and Specification
 - Representation of System States
 - Class/Object diagrams
 - Entity-relationship diagrams
 - Petri nets
- Model Transformation
 - Shift toward model driven code generation, analysis, computation and documentation
- Concurrency and Distribution
 - System behaviour represented by state changes is modeled by rule-based graph manipulations

- Software Development
 - Architectural structures and technical structures
 - Changes during software development process
 - Describe structural evolution in a rule-based way

What is a graph transformation?

The main idea of graph transformation is the rule based modification of graphs. The core of a rule (or production), p = (L, R) is a pair of graphs (L, R) known as the left-hand side L and the right-hand side R. Applying the rule p = (L, R) means

- to find a match L in the source graph and
- replacing L by R, thus leading to the target graph

Definition of a pushout

There are many approaches to graph transformations but here we discuss the algebraic approach which is based on pushout constructions. There are two variants of pushout:

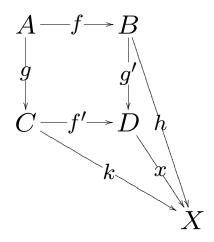
- double-pushout
- single-pushout

Definition: Given morphisms $f : A \to B$ and $g : A \to C$ in a category **C**, a pushout (D, f', g') over f and g is defined by

- a pushout object *D* and
- morphisms $f': C \to D$ and $g': B \to D$ with $f' \circ g = g' \circ f$

such that for all objects X and morphisms $h : B \to X$ and $k : C \to X$ with $k \circ g = h \circ f$ there is a unique morphism $x : D \to X$ such that $x \circ g' = h$ and $x \circ f' = k$

Definition of a pushout (continued) Visually:



We write $D = B +_A C$ for the pushout object D, where D is the gluing of B and C via (A, f, g).

Double-pushout approach

Two gluing constructions are used to model a graph transformation. In this approach a production is defined as p = (L, K, R), where:

- *L* describes the preconditions of the rule
- R describes the postconditions of the rule
- K describes a graph part which has to exist to apply the rule
- $L \setminus K$ describes the part to be deleted
- $R \setminus K$ describes the part to be created

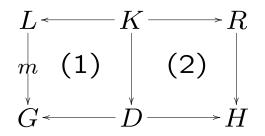
Double-pushout step by step process

Let m be the match of the left-hand side L in G such that m is structure preserving and let p be a production then:

- first create the context graph as $D := (G \setminus m(L)) \bigcup m(K)$ then
- then check for *gluing* condition
 - gluing condition points in L which are source or target in $G \setminus L(S_B)$ known as the boundary point and points given by K known as the gluing points (S_K) must satisfy $S_B \subseteq S_K$
- add R through K to D, to get H as $H = R +_K D$

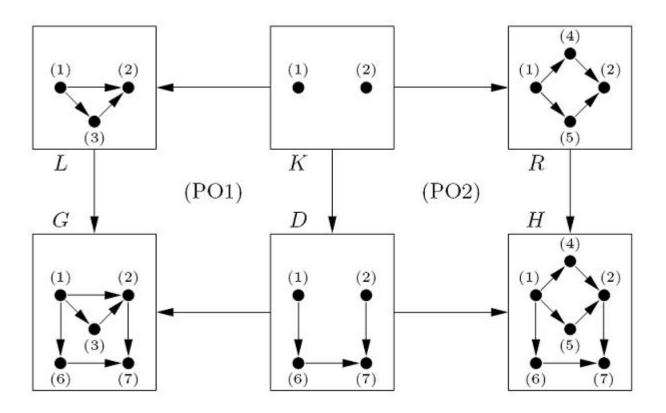
Since K is in L, R, D we express that there are morphisms $K \to L$, $K \to R$, $K \to D$ and pushout constructions $G = L +_K D$, $H = R +_K D$

Double-pushout step by step process (continued)



(1) construct D such that the gluing of L and D via K is equal to G(2) construct the gluing of R and D via K leading to the graph H

Simple Example



Tools

- AGG (http://tfs.cs.tu-berlin.de/agg/down.html)
- Progres (PRogrammed Graph REwriting Systems)
- Grace
- Fujaba

References:

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