A Brief History

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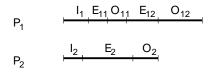
Systems that execute programs serially with no direct interaction between the user and the computer.

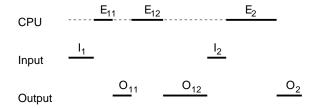
- The OS always resides in memory.
- Single user, no execution and I/O overlap.
- OS = a program to load and run user jobs, take dumps,

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• The CPU is often idle because I/O devices are much slower than the CPU.

Example



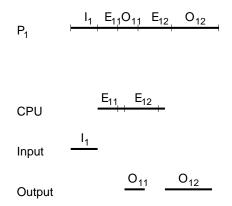


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Overlap the I/O of one job with its own computation. **I/O Interrupts** The interrupt handler signals the user process upon the completion of I/O.

- Synchronous. The user process makes an I/O request and waits until the I/O completion. The user process must know the I/O latency, so it knows how long it should wait.
- Asynchronous. An I/O request returns without waiting for the I/O completion. An I/O interrupt is scheduled at the completion of I/O. The interrupt handler signals the user process when the I/O is done. This allows concurrent I/O operations to several devices.

Example



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Systems that are designed to concurrently execute more than one task.

$$P_{1} \qquad \begin{array}{c} I_{1} \\ F_{1} \\ F_{2} \\ P_{2} \\ \end{array} \qquad \begin{array}{c} I_{2} \\ F_{2} \\ F_{2} \\ \end{array} \qquad \begin{array}{c} I_{2} \\ F_{2} \\ F_{2} \\ F_{2} \\ \end{array} \qquad \begin{array}{c} I_{2} \\ F_{2} \\ F_{2} \\ F_{2} \\ \end{array} \qquad \begin{array}{c} I_{1} \\ F_{2} \\ F_{$$

CPU
$$\underbrace{E_{11} \ E_{12}}_{\text{Input}} \underbrace{E_2}_{\text{Input}}$$

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$$P_{1} \xrightarrow{I_{1} E_{11} O_{11} E_{12} O_{12}} P_{2} \xrightarrow{I_{2} E_{2} O_{2}} P_{2}$$

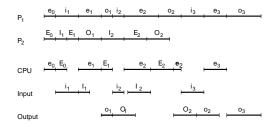
CPU
$$\underbrace{E_{11} \ E_{12}}_{\text{Input}} \underbrace{E_2}_{\text{Input}}$$

Issues:

- Memory protection + relocation.
- OS began to be an important science.

Timesharing or multitasking

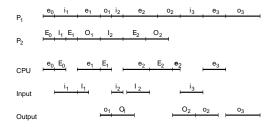
Systems that allow multiple users (programs) to run concurrently. The system switches from one user to another. Examples, MULTICS at MIT and UNIX at Bell Lab (1970). Example.



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Issues:

- Response time.
- Thrashing.

High speed, better performance/price ratio, fault tolerant. Symmetric multiprocessing (SMP). All processors run the same OS, no master-slave relationship, share memory.

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High speed, better performance/price ratio, fault tolerant. Symmetric multiprocessing (SMP). All processors run the same OS, no master-slave relationship, share memory.

Issues:

- data consistency.
- Ioad balancing.
- I/O bottle-neck.

Cache coherency An update to a variable in one cache must be immediately reflected in all other caches that hold the variable.

Operations have time limits on the operation of a processor or the flow of data. For example, sensor control. Hard real-time system. Hard time limit on critical tasks. Soft real-time system. A critical task get priority over other tasks and retains that priority until it completes.