

OBJECT-ORIENTED LITERATE PROGRAMMING

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Abstract

During the past decades, programming methodology has seen an improvement by structured programming and object-oriented programming (OOP), leading to software that is more reliable and easier to develop. However, software engineers are still dealing with problems in processing associated documentation. Literate programming was introduced by Donald Knuth in the early 80's as an approach to produce programs together with their documentation in a way that is aimed at consumption by humans rather than by compilers. However, dated and complex features, dependence on formatting and programming language, and a lack of methodology prevented the method from gaining in popularity.

In this thesis, we propose an approach to “integrate” OOP with literate programming in order to present and document the whole design in a consistent and maintainable way. In our approach, both program code and corresponding documentation are generated from the same source. The resulting documentation consists of code chunks with detailed explanations, as well as automatically generated class diagrams with varying levels of detail. A tool, *Spark*, has been developed and applied to the design of a Transit Information System from requirement analysis to testing. *Spark* was also used in its own development.

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Chapter 1

Introduction

1.1 Why Object-Oriented Literate Programming

Although object-oriented programming (OOP) is common in the area of computer software development, it is only one of evolutionary extensions to get to a software revolution [10]. The past several decades saw the development from early imperative programming through to OOP and from unstructured programming to structured one. Each programming paradigm and every progress has shortened the gap between human beings and machines, the real world and computer software [36, 40]. However, none of them escape the limitation of instructing a computer what to do. To make a further progress, we propose a new approach, *object-oriented literate programming* (OOLP), which combines the benefits of both OOP and literate programming and is assured with automated support. Such an approach is desirable for the following reasons:

- **Language-independence:** Donald Knuth’s literate programming encourages programmers to concentrate on explaining to human beings what they want a computer to do, which does lead to significantly better documentation. However, the fact that it fails to employ new programming techniques and its doomed complex features and language-dependence made one still wait for a better alternative [38]: in spite of the support of tools such as CWEB [25], noweb [38], FunnelWeb [1], FWEB [27], and OCAMAWEB [29], the dependence on certain programming language or documentation formatting language

is predeterminate. Our approach is more flexible: software engineers can choose their favorites, both programming languages, like Eiffel [32], Lime [41], Java, or C#, and formatting languages, such as L^AT_EX [26, 33], AsciiDoc [2], and DocBook [45].

- **Consistent, accurate, and readable documentation:** it is hard to say that any software application without qualified documentation is valuable, because documentation absences, errors and even omissions can undoubtedly lead to disasters for both end-users and successive software engineers. In fact, all software development projects must generate a large amount of associated documentation; a high proportion of software process costs is incurred in producing documentation. In our approach, any change in program code can affect its documentation directly and every algorithm, even a single variable, is discussed properly. This kind of work of literature makes reading technical documentation fun.
- **Automated tools support:** one of the main benefits of this approach is the ability to use software tools to analyze program code and generate design diagrams. With such tool support, we not only hope for an increase in development speed, but also for progression towards a more standardized form of documentation.
- **Object-oriented:** programming in an object-oriented language, like Java, is neither a necessary nor sufficient condition for being object-oriented; the key is the object-oriented design technique [5]. This approach maintains the concepts and techniques of OOP, so programs still may be seen as a collection of cooperating objects, which makes the code characterized by flexibility and maintainability, as opposed to traditional view in which a program may be seen as a list of instructions to the computer.
- **Simplicity:** it is because of the feature of language-independence that little extra training is needed. Experienced programmers can begin their OOLP immediately with their favorite OOP language, like Eiffel, and documentation formatting language, say L^AT_EX; as to beginners, they can choose the simplest

but powerful enough ones and get jobs done perfectly in a short term. In addition, there is no extra cost needed on other application software or equipments; a text editor is good enough. Configurable, platform-independent and open-source text editors such as VIM [35] and Emacs [7], are perfect for our job.

1.2 Contributions

My contributions include:

- design a scheme to combine OOP and literate programming,
- design and implement Spark, a set of tools supporting this technique,
- propose a reference development process,
- develop a case study, *Transit Information System*, in the way of OOLP.

OOLP, as well as Spark, is an ongoing research project and many aspects are likely to evolve over time. Therefore, the design of schemes and the implementation of Spark should be as general as possible. The code of Spark is written entirely in *perl* [11, 44], which is good at text processing. *Graphviz* [14, 20] is employed to produce diagrams, since we want to avoid looking deeply into layout algorithms.

1.3 Structure of the Thesis

The remainder of this thesis is organized as follows:

- Chapter 2 surveys the related work on object-oriented techniques, and literate programming. Both advantages and limitations of them are discussed.
- Chapter 3 introduces OOLP by describing its features and predominance as well as its supporting tools, Spark.
- Chapter 4 studies the case of Transit Information System, which is developed entirely in the way of OOLP.
- Chapter 5 focuses on the implementation and testing of Spark.

- Chapter 6 draws the conclusions of our work, in addition to discussing future work.

Chapter 2

Related Work

In this chapter, a survey of the existing work on both object-oriented techniques and literate programming implementations is developed where both its advantages and disadvantages are analyzed.

2.1 Object-Oriented Programming

After the 1990s, OOP became a mainstream technique in software development. It is widely used successfully in various applications including compiler, graphics, user interfaces, databases, object-oriented languages, computer-aided design systems, games, and control system [39]. Basing on abstraction, encapsulation, and polymorphism, OOP has more predominance than traditional approaches on *reliability*, *modularity*, *compatibility*, *portability*, and *efficiency* [31]. As studied in detail in subsequent sections, object-oriented techniques, some of which come from previously established paradigms, can significantly improve these quality factors, which is why it is so attractive.

2.1.1 Class and Object

In OOP, rather than just a list of instructions to the computer, an executing program may be seen as a collection of cooperating objects, which distinguishes object-oriented approach from other non-object-oriented methods [12]. Objects could be anything, including not only the entities, such as a student, an university, a hospital and a car,

but also the abstract concept, such as a plan and an event, so the models or programs conceived in such a natural way are more understandable.

As the smallest module of programs, an object is equipped with essential attributes and behaviors and becomes active by executing one of its methods, in which it may change its state or send messages to other objects, which in turn invokes the execution of the corresponding methods of those objects. Compare to structured methods, they focus on functional decomposition; once a complex program has been decomposed into some programmable sub-functions, the software certainly will work mechanically just like an assembly line. It is very likely that in order to produce some new “parts”, developers have to reconstruct a new “assembly line” for both new “parts” and old ones or maintain two or more “assembly lines” for all at the same time. The following workload and risk of new bugs could be a disaster for everybody concerned. While for object-oriented methods, developers only need to prepare some new class(es) for such an extension. Unfortunately, in practice the changes of whole workflow are more common than the ones of concrete methods used to process certain object. Hence, object-oriented programs can enjoy better extendibility and stability.

However, rather than the object, the class is the central concept of OOP. A class is a kind of abstract data type equipped with a possibly partial implementation, through which classes establish the necessary link with software construction — design and implementation. Much of the power of the object-oriented method, such as inheritance, encapsulation, and polymorphism, derives from this concept.

2.1.2 Inheritance

In OOP, inheritance is a strong conceptual relation that can hold between classes: a class may be an extension, specialization or combination of others. It is common that new software expands on previous developments, but problems of repetition and variation are largely ignored by traditional design methods. Writing the same code over and over again means not only waste of time, but also the risk of incoming inconsistencies and errors. With the support of inheritance, one class (subclass) can introduce itself by capturing the striking commonalities that exist within one or more mature classes (superclass) and adding the many differences that characterize individual cases.

The advantages of inheritance also cover a faster modifying mechanism: by inheritance, any modification in high level can affect the other related parts of system immediately. In addition, inheritance divides a system into different abstract levels, where developers can focus on them one by one without the bother of trivial details at the very beginning.

When a class inherits its members from more than one ancestor class, this is called *multiple inheritance*, which is a debatable feature. Generally speaking, multiple inheritance make the ancestor relation complex, so any change made in a certain ancestor may result in some unexpected effects on its successors and a compiler has to face the problem of processing those identical members. Nevertheless, there are still some object-oriented programming languages, such as Eiffel, C++, and perl, that support multiple inheritance with different ways. Eiffel will automatically join the members with same name and implementation together if they are not renamed by the programmer explicitly. C++ requires the programmer to state the inheritance path in detail. Perl uses the list of classes to inherit from as an ordered list. These contrast with Java and C#, which allow classes to inherit from multiple interfaces rather than multiple inheritance; this results in no ambiguity.

2.1.3 Assertion

For software programmers, reliability refers to the correctness and robustness of software. In order to improve reliability, assertions, boolean expressions, usually written as annotations are employed to specify what a system behavior is supposed to do rather than how it does. The use of assertion dates back to Hoare's 1969 paper on formal verification [22]. A correctness formula (also called Hoare triple) is an expression of the form

$$P \quad \{Q\} \quad R \quad (2.1)$$

where Q denotes a program; P and R , the properties of Q , are called precondition and postcondition respectively. However, it is only a mathematical notation used to constrain the properties of programs.

Design by Contract (DBC) proposed by Meyer as a trademark of Eiffel is a formal technique for dynamically checking specification violation during runtime [31]. The main idea behind DBC is that a client and a supplier, the elements of a software

system, collaborate with each other according to a “contract”. For example, suppose method M of class C provides a certain functionality needed by class D. Then, class D, the client, must fulfil a certain requirement, the precondition of method M, before invoking method M and as a result, class C, the provider, must ensure a certain property on the exit of method M, its postcondition. That is:

- precondition constrains the client, so it is an obligation for the client and a benefit for the supplier.
- postcondition constrains the supplier, so it is an obligation for the supplier and a benefit for the client.

Only precondition and postcondition are not enough, because they are used to describe the properties of individual methods. For the global states of a class, Eiffel supports the concept of class invariant. An assertion I is a correct class invariant for class C if and only if it meets the following two conditions:

- every constructor of C, when applied to arguments satisfying its precondition in a state where the attributes have their default values, yields a state satisfying I
- every exported method of C, when applied to arguments satisfying its precondition and a state satisfying I, yields a state satisfying I

Assertions used in preconditions, postcondition, and invariant express the semantic constraints on a class, so it is possible for us to define formally what it means for the class to be correct.

2.1.4 Garbage Collection

Garbage collection (GC) as one of automatic memory management techniques is used in most good object-oriented environments. The main idea is that a garbage collector, a facility included in the runtime system for a programming language, takes care of both detecting and reclaiming unreachable objects. With this technique, software developers do not need to worry about memory wasted on useless objects created by their software any more, so the reliability and timeliness of software products will surely benefit from it. GC was invented by John McCarthy around 1959 to solve

the problems of manual memory management in his Lisp programming language [30]. The basic principle of how a mark and sweep garbage collector works is:

- mark phase, starting from the origins, follows references recursively to traverse the active part of the structure, marking as reachable all the objects it encounters
- sweep phase traverses the whole memory structure, reclaiming unmarked elements and erasing all the marks

Classical garbage collectors are inactive as long as there is enough memory available for the application. Its advantage is that it causes no overhead before the collector is triggered and a serious potential drawback is that a complete mark-sweep cycle may take a long time — especially in a virtual memory environment. Therefore, GC is rarely used on embedded or real-time systems.

To pursue better performance, some techniques were employed. First of all, endow developers with some control over the activation and deactivation of collector cycles. If a system contains a time-critical section, which must not be subject to any unpredictable delay, the developer may put a “stop sign” at the beginning of the section and show a “green light” at the end; and at any point where the application is known to be idle, the developer may ask collector to work immediately. In addition, ones also use *generation scavenging*, the philosophy behind which is that the more garbage collection cycles an object has survived, the better chance it has of surviving many more cycles or even remaining forever reachable. Although this technique helps through lessen the frequency of collector cycles on “old” objects, there remains a need to perform full collections occasionally. Parallel collection, one of the practical solutions for GC, requires two separate threads: the application and the collector. During the execution of an object-oriented system, the application creates as many new objects as it needs; the collector free them continuously according to the principle mentioned above.

2.1.5 Object-Oriented Approaches

In contrast with structured approaches, which focus on functional decomposition from the perspective of “input-process-output”, many object-oriented approaches

have been derived from these exceedingly popular object-oriented techniques discussed above. Each of them has introduced a set of new modelings or notations. The rest of this section presents summarily five popular ones: Responsibility-Driven Design (RDD) [46], Object Modeling Technique (OMT) [39], Business Object Notation (BON) [43], Catalysis [13], and Vienna Development Method (VDM++) [16].

RDD

RDD, conceived by Rebecca Wirfs-Brock in 1990, is a shift from thinking about objects as data plus algorithms to thinking about objects as roles plus responsibilities. In a responsibility-based model, objects play specific roles and occupy well-known positions in the application architecture, which is a smoothly-running community of objects. Each object is accountable for a specific portion of the work and all objects collaborate in clearly-defined ways, contacting with each other to fulfill the larger goals of the application. By creating such a “community of objects,” and assigning specific responsibilities to each, developers build a collaborative and flexible model of their application. On the other hand, developing consistent control styles for similar parts of the application may reduce the maintenance costs and incorporating flexibility into the design at specific “flex points” allows for planned extensions. Therefore, responsibilities are a good way to think about the behaviors of complex software systems. RDD consists of the following steps:

- create a CRC (Class, Responsibility, Collaborator) card for each candidate class, which usually is a noun with a small well-defined set of responsibilities
- identify and assign the responsibilities to candidate classes by asking what this class’s objects need to know and what steps towards accomplishing each goal these objects should be responsible for
- find collaborations for candidate classes by asking what other objects need those result
- build inheritance hierarchies for all classes if necessary.
- identify subsystems by drawing the collaborations graph and then looking for strongly coupled classes

- construct protocols for each class
- implementing design

OMT

OMT, developed circa 1991 by Rumbaugh, Blaha, Premerlani, Eddy and Lorensen, is one of popular object-oriented development methods today. It is primarily used by system and software developers supporting full life-cycle development, targeting object-oriented implementations. Because of its simple core notation, OMT has proven easy to understand, to draw, and to use. So it continues to be successful in various application domains, such as telecommunication, transportation, and compilers. OMT consists of the following phases:

- analysis phase: understand and model the application and the domain within which it operates by formal models: the object model specifies what it happens to, the dynamic model specifies when it happens, and functional model specifies what happens.
 - object model: capture the static structure of a system by showing the objects in the system, relationships between these objects, and the attributes and operations that characterize each class of objects
 - dynamic model: describe the control flow, interactions, and operating sequences of the system and consist of multiple state diagrams
 - functional model: describe computations within a system
- system design phase: determine the overall architecture of the system
 - organize the system into subsystems
 - identify concurrency
 - allocate subsystems to processors and tasks
 - handle the boundary conditions and the system resources, especially the permanent data.
 - choose software control implementation

- object design phase: determine the full definitions of the classes and associations used in the implementation, as well as the interfaces and algorithms of the methods used to implement operations.
- implementation phase: discuss the specific details for implementing a system using programming languages and database management systems.

BON

BON, developed in the early 1990s by Jean-Marc Nerson and Kim Walden, is a means of extending the higher-level concepts of Eiffel into the realm of analysis and design aided by a set of graphical notations. These graphical notations do not include the associations, multiplicities, and state-charts that can be found in nearly all object-oriented analysis and design notations today. BON consists of informal charts, static architecture, class interfaces, dynamic scenarios and nine standard tasks are grouped into three parts:

- gather analysis information
 - delineate system borderline
 - list candidate classes
 - select classes and group into clusters
- describe the gathered structure
 - define classes
 - sketch system behavior
 - define public features
- design a computational model
 - refine system
 - generalize
 - complete and review system

Catalysis

Catalysis coined by Desmond D’Souza and Alan Cameron Wills is a method for component-based and object-oriented software development. It provides a strongly coherent set of techniques for business analysis and system development using Unified Modeling Language (UML) and is characterized by following:

- **Traceability** from business models to code
- **Precision**, with clear unambiguous models and documents
- **Component Based Development**
- **Reuse** of designs, specifications, problem domain models, and even architectures
- **Scalability** from small to large teams and projects
- **Process** that is flexible yet repeatable, with multiple “routes”

Catalysis believes that there is no single process that fits every project: each one has different starting points, goals, and constraints. Therefore, it provides a list of process patterns that help developers plan a project appropriately to their situation [13].

VDM++

VDM++ is extended by Nico Plat, Paul Mukherjee and, later, Marcel Verhoef from VDM. It employs a formal notation to complement and enhance object-oriented class models and its development process consists of the following ten steps:

1. determine the purpose of the model.
2. read the requirements.
3. analyze the functional behavior from the requirements.
4. extract a list of possible classes or data types and operations. Create a dictionary.

5. sketch out representations for the classes using UML class diagrams. This includes the attributes and the associations between classes. Transfer this model to VDM++ and check its internal consistency.
6. sketch out signatures for the operations. Again, check the model's consistency in VDM++. The development is continued by adding operation signatures (the formal parameters and the result) at the class diagram level.
7. complete the class or data type definitions by determining potential invariant properties from the requirements and formalizing them. To make the model more comprehensive, it is a good idea to review the model to check coverage of the requirement. Document important properties or constraints as invariants. Before being able to validate the model created so far it is also necessary to consider how to construct instances of the different classes. In VDM++, constructors are simply written as operations with the same name as the class in which they are defined.
8. complete the operation definition by determining pre- and postcondition and operation body, modifying the type definition if necessary.
9. validate the specification using systematic testing and rapid prototyping. Three methods are used here:
 - (a) integrity properties are formal descriptions of system properties that can be generated automatically by VDMTools.
 - (b) VDMTools supports validation using conventional testing techniques, including features to enable test coverage documentation directly at the VDM++ level.
 - (c) validation can be made executing models together with other code, e.g., a graphical front end.
10. implement the model manually or using automatic code generators that produce directly compilable code in C++ or Java.

2.1.6 Critiques

We have introduced some object-oriented techniques and five object-oriented approaches. However, it is not necessarily followed by a reusable, robust, modifiable, and maintainable software applications. The rest of this section explains several other issues that contribute to a satisfying software applications.

Design Patterns

The work of designing a good object-oriented software is easy to say, but difficult to do [18, 19]. Although design patterns may introduce some more classes through delegation and inheritance, they do provide partial solutions to some common problems, including analysis [17, 28], system design [6], middleware [34], process modeling [3], dependency management [15], and configuration management [4]. Let us take the *strategy design pattern* for instance. Suppose that there is a requirement of a duck pond simulation game, which can show a large variety of duck species swimming and making quacking sounds. Basing the standard object-oriented techniques and approaches discussed above, developers may naturally define one Duck superclass from which all other duck types inherit as shown in Figure 2.1. Since all ducks quack and swim, the superclass takes care of their implementations, while every subclass has to be responsible for implementing its own display function.

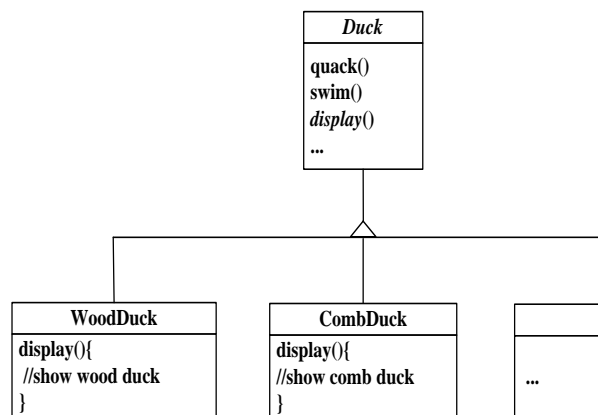


Figure 2.1: The initial structure of the duck game

Unfortunately, the extendibility of such a design structure is not satisfying. What will happen if here come two new requirements: let all existed ducks fly and add some rubber ducks, which can neither quack nor fly? Apparently, inheritance and overriding can not be the answer, because the specification will keep changing and developers will be forced to keep an eye on and possibly override fly method and quack method for every new subclass; trying to declare some interfaces, such as Flyable and Quackable, for the changes must lead to a mass of duplicate code.

The strategy design pattern seems to be a key to such a problem. The main idea of it is to decouple a policy-deciding class from a set of mechanisms so that different mechanisms can be changed transparently from a client. In other words, all the “problematic” behaviors, such as quack and fly, should be taken out of the super-class and then assigned to the specific duck according to concrete circumstances (see Figure 2.2). In this way, all the concrete strategies like FlyWithWings, FlyNoWay, Quark, and Mute can be substituted at runtime and new behaviors also can be added without modifying the other parts.

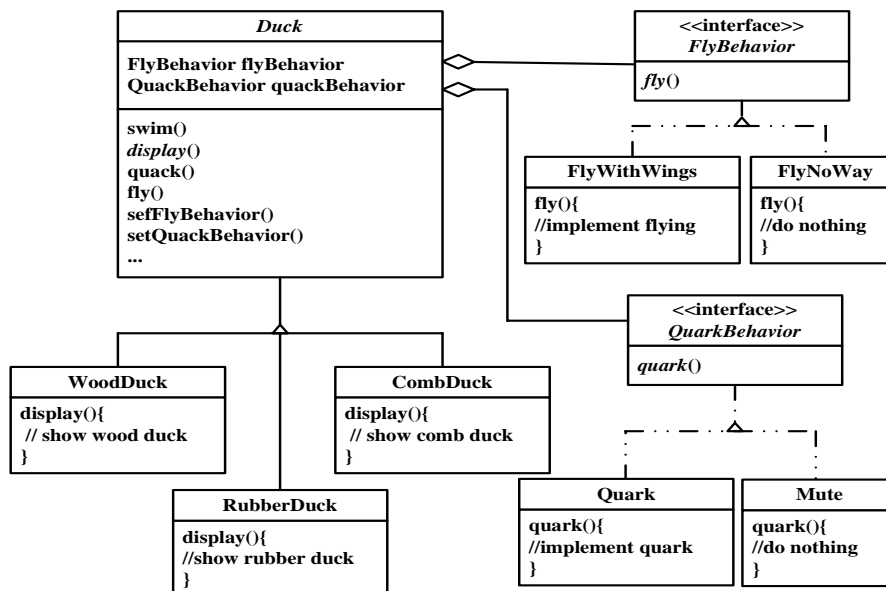


Figure 2.2: The refined structure using strategy design pattern

Through the example above, we can say that knowing the object-oriented basics,

abstraction, encapsulation, polymorphism, and inheritance, as well as some popular approaches do not make one a good object-oriented designer. As practical object-oriented experience, design patterns show designers how to build systems with better qualities: reusability, extensibility, and maintainability, not concrete program code. More and more good patterns are going to be discovered by the following principles:

- Encapsulate what varies.
- Favor composition over inheritance.
- Program to interfaces, not implementations.
- Strive for loosely coupled designs between objects that interact.
- Classes should be open for extension but closed for modification.
- Depend on abstractions rather than concrete classes.
- A class should have only one reason to change.

Software Documentation

Software documentation is written text that accompanies and explains computer software. Its absence, insufficiency, or inconsistency means the loss of the partial or even total previous effort, because the program will undergo modifications due to errors or changes of requirements and reuses in other software applications. All large software development projects, irrespective of application, generate a large amount of associated documentation, mainly including the project plan, quality plan, requirements specification architecture description, design documentation, technical documentation, user manuals and test plan [42].

OOP did achieve a major improvement in the analysis and design of software, but it also suffers the problems coming from the consistency and readability of software documentation. The reason of that comes mainly from the documentation mechanism itself. Like traditional programming paradigms, OOP separates most documentation, such as design documentation and architecture description, from code, so it is hard to keep all this pivotal documentation up-to-date and synchronized. Especially for large projects and plus the time pressure, the inconsistency of documentation gets worse.

On the other hand, although technical documentation, which is used to explain class as well as its members, data structures and algorithms, is embedded within the source code as comments and may be supported by automatic documentation tools, such as doxygen [21], Javadoc, and TwinText [37], it is always short and organized in an order suitable for compilers rather than human beings.

In addition, software documentation without graphical notations is neither expressive nor appealing. In fact, most popular programming approaches today have their own set of graphical notations used to create an abstract model for their target software systems, which, especially the complex ones, enhance the importance of graphical notion. Usually, these notions are constructed by developers manually and then included into the corresponding software documentation. So incomplete changes may lead to inconsistency, which is the reason for other readers' misunderstanding.

2.2 Literate Programming

Introduced by Donald Knuth in the early 80's, literate programming is an approach that combines a programming language with a formatting language, thereby making programs more robust, more portable, and more easily to maintain than programs written only in one high-level language [9, 24, 38]. Its main idea is to treat a program as a *work of literature*, which is used to explain to human beings what it let a computer do rather than to instruct a computer what to do. The program is also viewed as a hypertext document, rather like the World Wide Web. By contrast with other programming paradigms, the program source code is embedded into documentation rather than the other way and the practitioner of literate programming needs to manipulate two kinds of languages simultaneously, neither of which can provide significantly better documentation of programs by itself.

The first published literate programming environment is WEB [23], which uses Pascal as its underlying programming language and T_EX [26] for typesetting of the documentation. Pascal makes it possible to specify the algorithms formally and unambiguously, while T_EX provides typographic tools to explain the local structure of such parts. The structure of WEB program may be thought of as a “web” that is made up of many interconnected *modules*, which may contain the actual program source code, abbreviations for the code, and description of the code. All the modules

should be subdivided until their functionality is easily understandable. In WEB, the “bilinguist” writes such a program that serves as the source language for two different system routines as shown in Figure 2.3. Besides WEB, other implementations of this concept are CWEB [25], FWEB [27], noweb [38], FunnelWeb [1], and OCA-MAWEB [29]. Some of them are different versions of WEB for documenting specific programming languages, such as C++ and Fortan, while others are documentation formatting language independent, such as noweb, and FWEB.

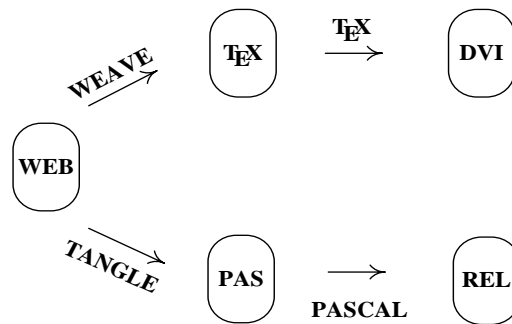


Figure 2.3: Dual usage of a WEB file (adapted from [24])

2.2.1 Advantages

As an efficient way to combine source code and its documentation, literate programming enhances the quality of programs. Its programs are characterized by flexible programming order, lossless factoring, better readability and better maintainability.

Flexible Programming Order

In literate programming, a program consists of some modules, which can be organized in arbitrary order without the constrain coming from compilers. So a programmer can choose the order best suited to explaining to human beings what he or she want a computer to do. In other words, this principle encourages the author of a literate program to take the time to consider each fragment of the program in its proper

position. The reordering is especially useful for encapsulating tasks such as input validation, debugging, and printing output fit for humans.

Factoring

Traditionally, a function is the smallest ordered list of computer instructions and the compiler requires the full text of its algorithm to be held together continuously. This is the reason that overweighed code chunks can be found everywhere. To improve this situation, Knuth introduced a decomposition facility into the meta-language. With this technique, the definition can be broken into constituent parts without the extra cost for defining new functions. Therefore, every part of any algorithm can be discussed in detail sufficiently.

Readability

Knuth believes that a program should be regarded as a work of literature. By such a literary style of writing, programmers enjoy the freedom to discuss the design decisions as well as constraints that may lead to certain intricacies in their implementation. A program presented in book form is certainly characterized by better readability.

Maintainability

Since factoring and literary style endow programmers with the ability to describe their algorithms as well as the trade-offs in detail fully, every reader, including the author, can understand the program totally at any time. When an alteration is required, it should be fairly obvious which part of the “book” needs to be modified. Similarly, the description concerning such a alteration will be used as a reference for other maintenance or development later.

In a word, maximized factoring, detailed description and literate sequence bring literate programs better readability, which in turn makes the programs easier to be maintained.

2.2.2 Disadvantages

We have introduced Knuth's literate programming as well as its advantages. It turns out to be a good approach to produce better documentation and to improve the quality of software, but literate programming has not become a mainstream technique in software development yet. The reasons for this reside in the fact that writing literate programs requires additional time in comparison to writing "illiterate" programs and the limitation of language-dependence.

Time Overhead

There are following several issues that contribute to the time overhead. Literate programmers need longer time to learn before starting to work than traditional programmers do. Besides specific programming language and compare to Javadoc and TwinText, which are Source Code Documentation Tools rather than literate programming tools, they have to learn to install and configure the set of applications that support literate programming. Additionally, the harder part is learning how to properly write literate programs. So the longer learning curve of literate programming is a challenge to the beginners' patience.

Literate programming forces programmers to develop software applications using a completely different perspective, where the developers should first make their thoughts clear to everybody. In order to fulfil this requirement, only the ability of programming is obviously insufficient. It is because there are too many choices of expression way, order, and factoring extend to choose that literate programmers have to sacrifice time for the best.

In addition to programming errors, two new types of errors are introduced by this technique: WEB structural errors and formatting errors. The former are those caused by the incorrect use of the WEB's own language required to define the structure of a program. Since both Weave and Tangle routines use such structure as an input, this kind of error can be propagated into programming and formatting language errors. Formatting errors are those caused by the misuse of formatting language. Similarly, these errors could affect other parts of a program.

For literate programmers, there is only one way to obtain the executable program. They have to run the Tangle routine over the WEB file first and then compile the

output. If there exist any programming errors, they can not be found until executable program is built. In order to correct them, developers have to go back to the WEB file, make changes, then run the Tangle, and compile the output again.

Language-dependence

The first published literate programming environment is WEB [24], introduced by Donald Knuth in 1981; this system uses $\text{T}_{\text{E}}\text{X}$ as the document formatting language and PASCAL as the programming language. It is true that as long as a person knows both of the underlying languages, there is no trick at all to learning WEB, but what does it mean for those who do not know these two languages or for the circumstance that the underlying languages do not suit for the programming of the target project?

In the section “The WEB System” of his Computer Journal article, Knuth addressed that the same principles would apply equally well if other languages were substituted: instead of $\text{T}_{\text{E}}\text{X}$ one could use a language like Scribe or Troff; instead of PASCAL, one could use ADA, ALGOL, LISP, COBOL, FORTRAN, APL, C, or even assembly language. However, all the literate programming systems derived from WEB depend on one or both underlying languages. CWEB is created by Donald Knuth and Silvio Levy as a follow up to Knuth’s WEB, using the C programming language instead of PASCAL. OCAMAWEB is a CWEB like literate programming tool, which is a combination of the MATLAB [8] language and $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X}$. Although noweb, FunnelWeb, and FWEB, can work with multiple programming languages, they still depends on their document formatting language respectively.

2.3 Summary

As a software application development approach in mainstream, object-oriented programming improve the quality of program, which includes reliability, modularity, compatibility, portability, and efficiency. The reason for this resides in the object-oriented techniques, such as class, object, inheritance, polymorphism, and abstraction. Its improvement on documentation and design approach as well as supporting tool is not satisfying. On the other hand, literate programming could produce significantly better documentation and improve the quality of software. However, its inevitable

time overhead and language-dependence prevent literate programming from being a mainstream technique in software development.

Chapter 3

Object-Oriented Literate Programming

The previous chapters explored the the goals of OOLP and related research. In this chapter, we take a closer look at OOLP from the perspectives of its key aspects and supporting tools.

3.1 Overview

Nothing concerning OOLP is intrinsically new; what we have done is combined a number of ideas that have been in the field for a time. All of these techniques have their own advantages. By applying them systematically in a slightly new way, we propose a new programming paradigm — OOLP, which is anatomized in the following sections.

The practitioner of OOLP can be regarded as an essayist whose main tasks are to break the whole program into little pieces and to order or reorder them for pursuing the best suited to explaining what this software is doing. Every algorithm, even a single variable, is discussed properly in its natural place. In this way, the program and its documentation, including diagrams, are always consistent with each other. On the other hand, it still can be viewed as a collection of loosely connected objects, each of which is responsible for a certain specific task, which is a natural way for human beings to cognise the world. Therefore, this kind of *works of literature* is

characterized by readability, reusability, flexibility, and maintainability.

Class diagrams are used in nearly all object-oriented analysis and design methods today. They can present readers a clear and intuitive view of the system structure. All existing literate programming tools would require developers to draw them by hand. In OOLP, Spark allows them (see Figure 3.1) to be generated automatically and inserted around the corresponding code part. Such automatic feature of Spark not only lightens developers' workload, but also ensures the consistence of class diagrams with the program code.

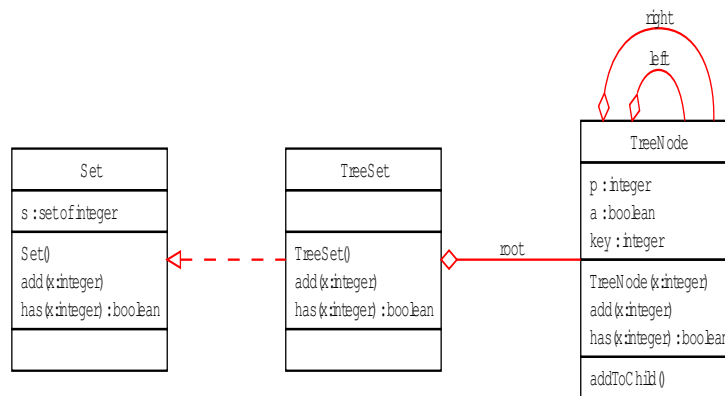


Figure 3.1: An example of automatic class diagram.

Like WEB, the programming environment of OOLP itself is chiefly a combination of two other languages: (1) a documentation formatting language and (2) a programming language. The difference is that programmers can choose their favorite or most suitable combination of these two kinds of languages. The main point is to let the practitioners of OOLP enjoy the power of the inherently bilingual tool, and get rid of the limitation of language-dependence.

In OOLP, the documentation formatting language provides tools to explain the local structure of documentation parts and to build the documentation that describes the program clearly and that facilitates program maintenance, while the programming language makes it possible to specify the algorithms formally and to obtain a machine-executable program. In addition, the supporting tool, Spark, is responsible

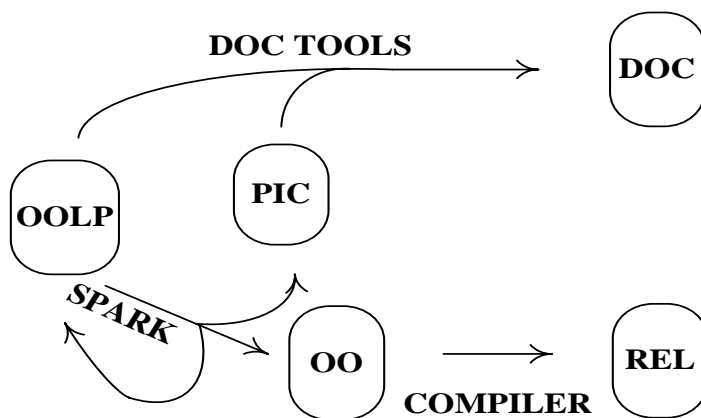


Figure 3.2: Overview of OOLP environment

for reconstructing compiler-acceptable code files, inserting continuous program code back to documentation, and building graphical notations.

3.2 Document Formatting Languages

Since OOLP does not fix on any specific document formatting language, a programmer can choose anyone from the popular text based document generation systems, such as AsciiDoc [2], \LaTeX , and DocBook [45]. The main point is that the target source file can be edited by arbitrary text editor platform-independently and that Spark can parse and process the target source file.

Let us look at this process in slightly more detail. Suppose AsciiDoc is used as the document formatting language and we have written an OOLP program and put it into a computer text file called EXAMPLE.TXT. The concrete syntax of AsciiDoc can be found in Appendix E. To generate hardcopy documentation for the program, we can run *asciidoc.py*, which is an executable program that takes the file EXAMPLE.TXT as input and produces another file as output. By setting different command line parameters, we can ask AsciiDoc to produce several predefined back end outputs, including DocBook, HTML, LinuxDoc, and \LaTeX . Take the \LaTeX output for example, after running the following command, we can have a file EXAMPLE.TEX as output.

```
asciidoc.py -unsafe -backend=latex EXAMPLE.TXT
```

Then we run the \LaTeX processor, which takes `EXAMPLE.TEX` as input and produces `EXAMPLE.PDF` as output.

```
pdflatex EXAMPLE.TEX
```

By default, AsciiDoc produces plain HTML 4.01 file. We can simply run the following command,

```
asciidoc.py EXAMPLE.TXT
```

Then, we will get a file named `EXAMPLE.HTML`.

To use AsciiDoc, we need to setup the environment first (see Appendix A).

The process is the same for other document formatting languages, but the features of OOLP are denoted a little differently. For example, a piece of program code is listed in AsciiDoc as following:

```
-----  
feature {ANY}  
  STAFF...match (id: INTEGER): BOOLEAN is  
    require  
      id >= 0  
    do  
      Result := id = number  
    end  
-----
```

In DocBook, the same code goes as following:

```
<programlisting>  
feature {ANY}  
  STAFF...match (id: INTEGER): BOOLEAN is  
    require  
      id >= 0  
    do  
      Result := id = number  
    end  
</programlisting>
```

In \LaTeX , it is listed as following:

```
\begin{verbatim}
feature {ANY}
  STAFF...match (id: INTEGER): BOOLEAN is
    require
      id >= 0
    do
      Result := id = number
    end
\end {verbatim}
```

Graphic notations are included in different ways too. For example, in AsciiDoc, a picture is included as following:

```
image::hpta_transit_info.jpg[Object Model]
// $ HPTA_TRANSIT_INFO DATABASE FILE_DATABASE @VERTICAL
```

“hpta_transit_info.jpg” is the picture’s name and followed by its attribute, “Object Model”. “//” denotes a comment line, which will be omitted by AsciiDoc compiler, but Spark considers it as a setting of the diagram: HPTA_TRANSIT_INFO, DATABASE and FILE_DATABASE are explained as the classes included in this diagram; @VERTICAL means that the diagram must be drawn vertically. Other settings are discussed in Chapter 5.

In DocBook, the same picture is included as following:

```
<figure><title>Object Model</title>
<graphic fileref="hpta_transit_info.jpg"></graphic>
<!-- $ HPTA_TRANSIT_INFO DATABASE FILE_DATABASE @VERTICAL -->
</figure>
```

In L^AT_EX, it is included as following:

```
\includegraphics[width=100mm, height=65mm]{hpta_transit_info.jpg}
% $ HPTA_TRANSIT_INFO DATABASE FILE_DATABASE @VERTICAL
```

3.3 Programming Languages

In OOLP, programmers can also choose their programming language from multiple popular candidates, such as Java, C#, Eiffel, Lime, and C++. In principle, any programming language, like PASCAL, Basic, and even assembly language, is eligible for being such a candidate, but in this paper, we only focus on object-oriented programming languages.

Since in literate programming, the continuous program written in certain programming language has been broken into sections and ordered best for explaining to human beings, the traditional process of “compile, load, and go” has been slightly lengthened to “reassemble, compile, load, and go”.

3.4 Spark

Spark consists of two parts: front end and back end. The front end takes an OOLP program as input and produces a number of program source code files as well as one graphic notation script file; the back end takes the graphic notation script file as input and produces a number of DOT files, which are used to feed GraphViz. GraphViz generates all the diagram files upon the request (see Figure 3.3). This structure decouples the programming language parsing from the algorithm of diagram layout so that different mechanisms can be changed transparently from each other.

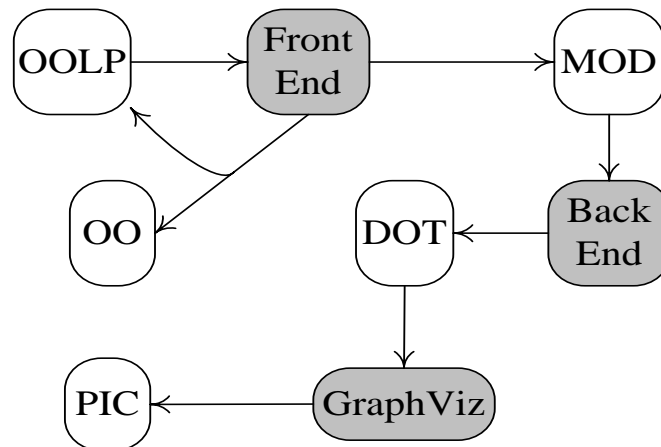


Figure 3.3: The Workflow of Spark

Two issues contribute to the fact that Spark gets rid of the limitation of the language-dependence in all existing literate programming tools. First, Spark focuses only on the code blocks and picture blocks and considers the other parts in OOLP programs as comments. In other words, as long as the document formatting language can work well by itself, Spark can accept it. Second, by providing different front

ends, Spark can be easily extent to adapt to various programming languages (see Appendix D).

It is because Spark parses program code partially that it can help to debug the program. Spark generates graphic notation automatically. This not only lightens the workload of developers, but also ensures the consistence of diagrams with code. The usage of Spark is simple, i.e. the front end followed by the OOLP program file. For example, the front end for Eiffel is chose and the program is still EXAMPLE.TXT, then the command is:

```
perl sparkf-eiffel.pl EXAMPLE.TXT
```

3.5 Editors

All the supporting tools are independent software and can be either embedded into any extensible edit platform as plug-ins or called under OS shells, So there is no specific requirement for its editor. In this paper, as an example we choose VIM, which is a highly configurable text editor built to enable efficient text editing platform-independently. In what it follows we will show how to build a integrated development environment (IDE) by using the supporting tools as well as VIM.

```
vim EXAMPLE.TXT
```

To make the work easier, we can define a new command for VIM as following:

```
:command Spark :!perl sparkf.pl %
```

Then as long as finishing editing the OOLP program in VIM, we can switch to the command mode and input the new command set above as following:

```
:Spark
```

Whenever the command, Spark, is invoked, it begins to parse the current source file, and then both program code files and graphic notation files have been generated immediately if no errors. In this way, we can call the compiler to compile the source code and use other document formatting language tools to produce the consistent document.

3.6 Reference Developing Process

Basing the existing object-oriented approaches such as RDD, OMT, BON, Catalysis and VDM++ (see Section 2.1.5), we propose a reference developing process, of which each step is discussed in details in the rest of this section.

3.6.1 Analysis

In object-oriented software development, this phase takes the input of a fuzzy, minimal, possibly inconsistent target specification and produces the output of a understanding, complete, consistent description of essential characteristics and behavior. The final product, object, distinguishes object-oriented analysis from other approaches, such as structured analysis and Jackson's method [12].

Creating a Dictionary

The correctness of understanding the main terms used in the requirements is the key to get the correct model of the target system, so the dictionary must be as detailed and rigorous as possible. The potential classes and types identified in the dictionary could then form the basis of a class diagram, whereas the potential operations might be described as use cases. This idea comes from VDM++.

Identifying Classes

Object-oriented software consists of classes, which describes a group of objects with similar properties, common behavior, common relationships to other objects. So the main task of this step is to find out all classes from the dictionary constructed and keep the number of entities in the initial model as small as possible at the same time. The principles are listed as following:

- Omit those nouns, if they are irrelevant with the purpose of the system.
- Model those nouns as attributes, if they have only trivial functionality.
- Create an overall class to represent the entire system so that the precise relationships between the different classes and their associations can be expressed there.

- Whenever an association is introduced consider its multiplicity and give it a role name.
- Try to keep encapsulation by the modifiers such as private and protected.
- Document important properties or constraints as invariants.

Sketching Operations

An operation is a function or transformation that may be applied to or by objects in a class. The aim of this step is to try to describe all the operations listed in the dictionary with signature (parameters and result) and formal specification (pre- and postcondition). Then, assign them to the classes identified respectively. This idea comes from BON, Catalysis, and VDM++.

Constructing Initial Model

An object model captures the static structure of a system by showing the objects in the system, relationships between the objects and the attributes and operations that characterize each class of objects. This model provides an intuitive graphic representation of a system and is valuable for communicating with customers. This idea comes from OMT and VDM++; Spark supports the automatic generation of this model.

3.6.2 Design

During analysis, the focus is on what needs to be done. During design, decisions are made about how the problem will be solved better. This goal can be approached more efficiently by employing the successful experience such as existing business models and design patterns.

Consulting Existing Business Models

A business model is the object model that has been employed successfully in a certain actual project. Such well-found business model can bring us not only the speed of development, but also the quality of software product, after all it have passed

the enough arguments and testing. So as long as a business model can meeting the requirements of the target system entirely or partially, we should replace the initial model generated in analysis with it entirely or partially. This idea comes from Catalysis.

Applying Design Patterns

We can not expect to find out everything in our library; in more cases, we need construct a new one. Then, what kind of design is good design? This question is ignored by most existing object-oriented approaches. We recommend design patterns, which can provide the answer (see Section 2.1.6). In this step, many new classes may be introduced into the current model, but this is we have to pay for the design with better flexibility, extensibility and reliability.

3.6.3 Implementing

The goal of Analysis and design is to bridge the gap between the real world and computer domain; the goal of this portion is discuss the specific details for implementing a system using programming languages. By the aid of Spark and literate tools, one can debug the program and view the document freely.

3.6.4 Testing

Testing is the process used to measure the quality of developed computer software. Since software is developed by human beings, it is inevitable that there exist some errors. Therefore, testing must be conducted for every computer software.

In order to cover the correctness, completeness, security, reliability, usability and portability of software, one needs to perform the following tests:

- **Usability testing**, which tries to find faults in the user interface design of the system.
- **Unit testing**, which tries to find faults in participating objects.
- **Integration testing**, which is the activity of finding faults when testing the individually tested components together.

- **System testing**, which tests the entire system.
 - **Functional testing**, which tests the requirements.
 - **Performance testing**, which checks the design goals.
 - **Acceptance testing**, which check the system against the project agreement and is done by the customer.

3.7 Summary

This chapter has introduced the following major features of OOLP and Spark:

- Including program code and graphic notations in various document formatting languages.
- Setting graphic notations
- Setting up the developing environment.
- Using Spark to generate graphic notation files and program code.
- Constructing OOLP IDE with VIM.
- A reference developing process

Chapter 4

Transit Information System Case Study

In this chapter, by an example, Transit Information System, we demonstrate how to use the technique of OOLP in developing software. The source code is listed in Appendix B. The rest of this chapter that follow is the actual output of an OOLP program file.

4.1 Transit Information System

4.1.1 Requirements

In this project, we are asked to develop an information system for a local train and bus service. Our customer, HPTA (Happy Passenger Transit Authority), has no clear picture what it should do, except to increase customer satisfaction and make traveling more attractive. All the information we have goes as follows:

- It will be used by passengers as well as by HPTA staff.
- Selected staff members would be allowed to update the information.
- Passengers should be able to enter their start and destination, a desired time, and get a bunch of possible connections.
- Connections can be direct or with changing busses or trains.

- For each bus and train station, the information like opening hours and accessibility is maintained.
- Users can browse a list of all bus and train routes or check the details of a certain route..
- Some bus stops and train stops are conjoint, but some not.
- Trains have two-digit numbers and busses have three-digit numbers.
- Connections between trains and busses must have at least five minutes for the change.

For simplicity, we assume that detours and delays do not occur, stops are never skipped.

4.1.2 An Overview

The following picture (Figure 4.1) is the object model of transit information system. As the root class, HPTA_TRANSIT_INFO controls the whole system from the beginning to the end. Class DATABASE is a deferred class, whose subclasses, such as class FILE_DATABASE, are responsible for maintaining system data. Class CONNECTION_FINDER is also a deferred class, whose subclasses, such as class PRIME_FINDER, are responsible for finding the possible connections.

The purpose of the application is to maintain the system information, including local train or bus service and the status of staffs, and provide users current public transit service information, including possible connections and routes.

4.1.3 Dictionary

To understand the main terms used in the requirements, we create a dictionary as following:

- *passenger*: a person, who want to get his or her destination by bus or train.
- *staff*: a person, who works for HPTA.

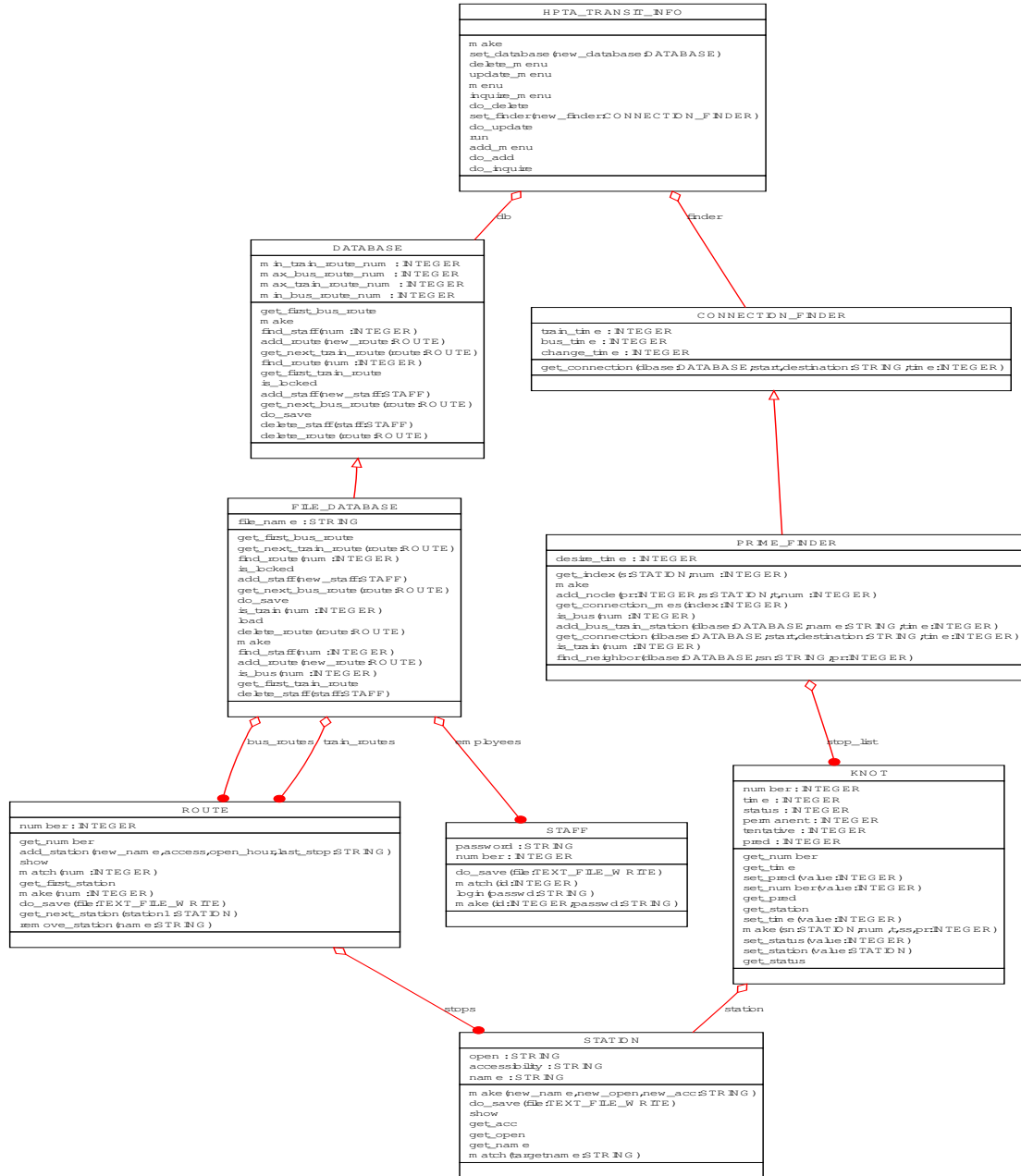


Figure 4.1: Object model of transit information system.

- *start*: a station, where a passenger begin his or her journey.
- *destination*: a station to which a passenger is going or directed.
- *desired time*: an interval, within which one want get to the destination from the start.
- *connection*: a sequence of stations.
- *bus*: a long motor vehicle for carrying passengers, usually along a fixed route.
- *train*: a series of connected railroad cars pulled or pushed by one or more locomotives.
- *route*: a course for buses or trains to travel from one station to another.
- *opening hour*: a time, at which the first vehicle departs.
- *accessibility*: a description of the running status of a station.
- *update*: a change of system information.
- *browse*: a display of the information of all routes.
- *check*: a detail show of a certain route information.

4.1.4 Identifying Class

The following classes are identified from the requirements.

```
class HPTA_TRANSIT_INFO
end
```

Class HPTA_TRANSIT_INFO is identified as a class of the entire system.

```
class STAFF
feature {NONE}
  number: INTEGER
  password: STRING
end
```

STAFF is a class with attributes employee number and password. The requirements state that selected staff members would be allowed to update the system.

```

class STATION
feature {NONE}
  name: STRING
  open: STRING
  accessibility: STRING
end

```

STATION is a class with attributes name, opening hour, and accessibility.

```

class ROUTE
feature {NONE}
  number: INTEGER
  stops: LINKED_LIST [STATION]
end

```

ROUTE is a class with attributes station list and route number.

4.1.5 Identifying Operations

All three operations listed in the dictionary belong naturally in the class HPTA_TRANSIT_INFO, because they are dependent on the interface of the system.

- login should belong in class STAFF, because it keep the secret of a certain staff.

4.1.6 Consulting The Library of Model

There is no suitable business model in our existing library, so we have to build this system from the beginning.

4.1.7 Applying Design Patterns

According to the requirements, our application needs to keep all system information and to calculate possible connections. There exist so many different methods for these two tasks. Hence, we apply the strategy design pattern. We declare two deferred classes

```

deferred class DATABASE
end

```

and

```

deferred class CONNECTION_FINDER
end

```

Then, we define two private members in class HPTA_TRANSIT_INFO denoted by the class name followed by three dots as following:

```

feature {NONE}
  HPTA_TRANSIT_INFO... db: DATABASE
  finder: CONNECTION_FINDER

```

i.e.

HPTA_TRANSIT_INFO
db :DATABASE finder:CONNECTION_FINDER

Figure 4.2: The attributes of class HPTA_TRANSIT_INFO

In this way, we can add new algorithms easily and even change mechanisms at runtime with the following private methods:

```

feature {NONE}
  HPTA_TRANSIT_INFO... set_finder(new_finder: CONNECTION_FINDER) is
    require
      new_finder /= Void
    do
      finder := new_finder
    ensure
      finder = new_finder
    end

```

and

```

feature {NONE}
  HPTA_TRANSIT_INFO... set_database(new_database: DATABASE) is
    require
      new_database /= Void
    do
      db := new_database
    ensure
      db = new_database
    end

```

Their preconditions require that the new comers are not invalid and their post-conditions ensure that the private member db and finder are set correctly.

Class CONNECTION_FINDER describes the interface that is common to all concrete mechanisms as following:

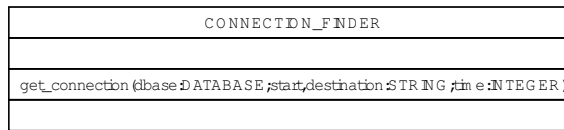


Figure 4.3: Class CONNECTION_FINDER

```

feature {HPTA_TRANSIT_INFO}
  CONNECTION_FINDER...
  get_connection (dbase: DATABASE; start, destination: STRING; time: INTEGER): STRING is
    require
      start /= Void
      destination /= Void
      time >= 0
      dbase /= Void
    deferred
  end

```

Class DATABASE describes the interface that is common to all concrete data maintain mechanisms as following:

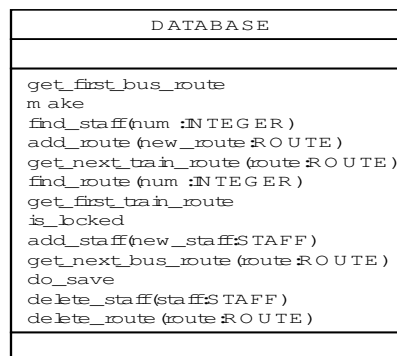


Figure 4.4: Class database

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_first_bus_route: ROUTE is
    deferred
  end

```

This method can return the first bus route object and is used by class HPTA_TRANSIT_INFO and class CONNECTION_FINDER only. Together with the following method, its clients can browse all bus routes one by one.

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_next_bus_route(route: ROUTE): ROUTE is
    deferred
  end

```

Similarly, we can browse all train routes by the following two methods:

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_first_train_route: ROUTE is
    deferred
  end

```

and

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_next_train_route(route: ROUTE): ROUTE is
    deferred
  end

```

Browsing all staff information is not necessary, but we need to find given staff object by the following method.

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... find_staff (num: INTEGER): STAFF is
    require
      num > 0
    deferred
  end

```

This method can return an STAFF object, whose employee number equals to the parameter *num*. It is because all employee number start from 1 that the precondition is added.

For convenience, we also provide a route finding method as follows:

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... find_route (num: INTEGER): ROUTE is
    require
      num >= min_train_route_num
      num <= max_bus_route_num
    deferred
  end

```

The following method is the creation of class DATABASE and invoked by class HPTA_TRANSIT_INFO only.

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... make is
    deferred
  end

```

In order to update system information, class DATABASE also requires the interfaces of adding and deleting as following:

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... add_route(new_route: ROUTE) is
    require
      new_route /= Void
    deferred
  end

```

and

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... delete_route (route: ROUTE) is
    require
      route /= Void
    deferred
  end

```

These two methods can add or delete a certain route to or from this system respectively and is called by class HPTA_TRANSIT_INFO only.

Similarly, class HPTA_TRANSIT_INFO also can add or delete a certain staff by the following two methods:

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... add_staff(new_staff: STAFF) is
    require
      new_staff /= Void
    deferred
  end

```

and

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... delete_staff (staff: STAFF) is
    require
      staff /= Void
    deferred
  end

```

As long as some system information is updated, DATABASE object must be informed to save the change by the following method.

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... do_save is
    deferred
  end

```

According to the requirements, only selected members can update the system. We define that when the database is locked, only the user, who knows both employee number and password, can conduct an update.

```

feature {HPTA_TRANSIT_INFO}
  DATABASE... is_locked: BOOLEAN is
    deferred
  end

```

The subclasses of these two deferred classes implement each concrete behavior mentioned above.

The following four constants are used to point out the bound of route number

```
feature {NONE}
  DATABASE... max_bus_route_num: INTEGER is 999
  min_bus_route_num: INTEGER is 100
  max_train_route_num: INTEGER is 99
  min_train_route_num: INTEGER is 10
```

4.1.8 Algorithms Design

File Database

For simplicity, we save the system information in a file named "sys_info.txt". So we define a subclass of class DATABASE, FILE_DATABASE as following:

```
class FILE_DATABASE
  inherit DATABASE
  feature {NONE}
    file_name: STRING is "sys_info.txt"
  end
```

i.e.

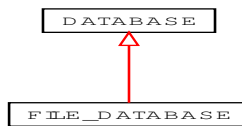


Figure 4.5: The hierarchy of databases

class FILE_DATABASE keep bus routes, train routes and staffs with LINKED_LIST as following:

```
feature {NONE}
  FILE_DATABASE... train_routes: LINKED_LIST [ROUTE]
  bus_routes: LINKED_LIST [ROUTE]
  employees: LINKED_LIST [STAFF]
```

now, class FILE_DATABASE becomes:

The creation of FILE_DATABASE is method make

```
create FILE_DATABASE... make
```

The main task of make is to initialize this three list

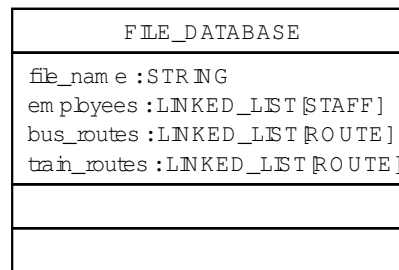


Figure 4.6: The attributes of FILE_DATABASE

```

feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... make is
  do
    create employees.make
    create bus_routes.make
    create train_routes.make
    load
  ensure
    employees /= Void
    bus_routes /= Void
    train_routes /= Void
end

```

and to load the system information for that file:

```

feature {NONE}
FILE_DATABASE... load is
  local
    input_string : STRING
    text_file_read : TEXT_FILE_READ
    text_file_write : TEXT_FILE_WRITE
    split : ARRAY [STRING]
    new_staff : STAFF
    route : ROUTE
  do
    create text_file_read.connect_to(file_name)
    if text_file_read.is_connected then
      from text_file_read.read_line
      until text_file_read.end_of_input
      loop
        if text_file_read.last_string.upper = 1 then
          inspect text_file_read.last_string.first.to_upper
          when 'S' then
            text_file_read.read_line
            input_string := text_file_read.last_string.twin
            split := input_string.split
            create new_staff.make (split.first.to_integer, split.last)
            employees.add_last(new_staff)
          when 'B', 'T' then
            text_file_read.read_line
            input_string := text_file_read.last_string.twin
            split := input_string.split
            route := find_route(split.item(4).to_integer)
            if route = Void then
              create route.make(split.item(4).to_integer)
            end
          end
        end
      end
    end
  end

```

```

route.add_station(split.first, split.item(2), split.item(3), split.last)
if split.item(4).to_integer > max_train_route_num then
    bus_routes.add_last(route)
else
    train_routes.add_last(route)
end
else
    route.add_station(split.first, split.item(2), split.item(3), split.last)
end
else
end
end
text_file_read.read_line
end
text_file_read.disconnect
else
    create text_file_write.connect_to(file_name)
    if text_file_write.is_connected then
        text_file_write.disconnect
    end
end
end
end

```

By the following method, one can get the specific route object.

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
FILE_DATABASE... find_route (num: INTEGER): ROUTE is
    local
        i: INTEGER
        route: ROUTE
    do
        if num > max_train_route_num then
            from i := bus_routes.lower
            until i > bus_routes.upper or else bus_routes.item(i).match(num)
            loop
                i := i+1
            end
            if i <= bus_routes.upper then
                route := bus_routes.item(i)
            end
        else
            from i := train_routes.lower
            until i > train_routes.upper or else train_routes.item(i).match(num)
            loop
                i := i+1
            end
            if i <= train_routes.upper then
                route := train_routes.item(i)
            end
        end
        Result := route
    end
end

```

Similarly, using the following method, one can get the staff with such employee number:

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
FILE_DATABASE... find_staff (num: INTEGER): STAFF is
    local
        i: INTEGER
        staff: STAFF
    do
        from i := employees.lower
        until i > employees.upper or else employees.item(i).match(num)
        loop

```

```

        i := i+1
    end
    if i <= employees.upper then
        staff := employees.item(i)
    end
    Result := staff
end

```

By the following four methods, one can browse all train routes and bus routes:

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
FILE_DATABASE... get_first_bus_route: ROUTE is
    local
        route: ROUTE
    do
        if not bus_routes.is_empty then
            route := bus_routes.first
        end
        Result := route
    end
get_next_bus_route(route: ROUTE): ROUTE is
    require
        bus_routes.index_of(route) > 0
    local
        next_route: ROUTE
    do
        if bus_routes.index_of(route) < bus_routes.upper then
            next_route := bus_routes.item(bus_routes.index_of(route)+1)
        end
        Result := next_route
    end
get_first_train_route: ROUTE is
    local
        route: ROUTE
    do
        if not train_routes.is_empty then
            route := train_routes.first
        end
        Result := route
    end
get_next_train_route(route: ROUTE): ROUTE is
    require
        train_routes.index_of(route) > 0
    local
        next_route: ROUTE
    do
        if train_routes.index_of(route) < train_routes.upper then
            next_route := train_routes.item(train_routes.index_of(route)+1)
        end
        Result := next_route
    end
end

```

By the following method, HPTA_TRANSIT_INFO object can add an arbitrary route to this database

```

feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... add_route(new_route: ROUTE) is
    do
        if is_bus(new_route.get_number) then
            bus_routes.add_last(new_route)
        elseif is_train(new_route.get_number) then
            train_routes.add_last(new_route)
        end
    end
end

```

By the following method, HPTA_TRANSIT_INFO object can add a staff to this database

```
feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... add_staff (new_staff: STAFF) is
do
    employees.add_last(new_staff)
end
```

By the following method, HPTA_TRANSIT_INFO object can remove an arbitrary route from this database

```
feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... delete_route (route: ROUTE) is
do
    if is_bus(route.get_number) then
        bus_routes.remove(bus_routes.index_of(route))
    elseif is_train(route.get_number) then
        train_routes.remove(train_routes.index_of(route))
    end
end
```

By the following method, HPTA_TRANSIT_INFO object can remove a staff from this database

```
feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... delete_staff (staff: STAFF) is
do
    employees.remove(employees.index_of(staff))
end
```

In FILE_DATABASE, as long as employees is not empty, this database is locked, which means you have to log in before updating.

```
feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... is_locked: BOOLEAN is
do
    Result := not employees.is_empty
end
```

Whenever the database is changed, it have to save the new data to the specific file by the following method:

```
feature {HPTA_TRANSIT_INFO}
FILE_DATABASE... do_save is
local
    file_2_write: TEXT_FILE_WRITE
    i: INTEGER
do
    create file_2_write.connect_to(file_name)
    if file_2_write.is_connected then
        from i := employees.lower
        until i > employees.upper
        loop
            employees.item(i).do_save(file_2_write)
            i := i + 1
        end
        from i := bus_routes.lower
```

```

    until i > bus_routes.upper
    loop
        bus_routes.item(i) .do_save(file_2_write)
        i := i + 1
    end
    from i := train_routes.lower
    until i > train_routes.upper
    loop
        train_routes.item(i) .do_save(file_2_write)
        i := i + 1
    end
    file_2_write.disconnect
else
    io.put_string("Update_database_failed!%N")
end
end

```

For convenience, we define the following two methods to tell if the current route is bus or train route:

```

feature {NONE}
FILE.DATABASE... is_train(num: INTEGER): BOOLEAN is
do
    Result := num >= min_train_route_num and num <= max_train_route_num
end
is_bus(num: INTEGER): BOOLEAN is
do
    Result := num >= min_bus_route_num and num <= max_bus_route_num
end

```

Conditional Shortest Path

According to the requirement that connections between trains and busses must have at least five minutes for the change, we have to consider bus station and train station as two different stations even they share the same name. In addition, we define a constant `change_time` in class `CONNECTION_FINDER`, whose subclasses need it.

```

feature {NONE}
CONNECTION_FINDER... change_time: INTEGER is 5

```

For convenience, we assume that a bus needs 2 minutes to get to the second stop and a train needs only 1 minute. So we also define the following two members in class `CONNECTION_FINDER`.

```

feature {NONE}
CONNECTION_FINDER... train_time: INTEGER is 1
bus_time: INTEGER is 2

```

`PRIME_FINDER` is one of the subclasses of `CONNECTION_FINDER`

```

inherit
PRIME_FINDER... CONNECTION_FINDER

```

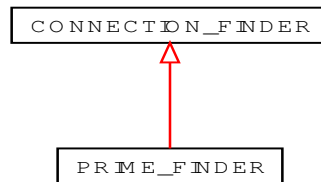


Figure 4.7: The hierarchy of class CONNECTION_FINDER

i.e.

Our first algorithm, PRIME_FINDER, is that starting from the start stations, including both bus station and train station, we search for all direct neighbors one after another and calculate their time respectively. In this way, as long as we found the destination as the next neighbor or no more new neighbors before get to the destination, our searching work is done.

To implement this algorithm, we declare list in class PRIME_FINDER

```

feature {NONE}
PRIME_FINDER... stop_list: LINKED_LIST [KNOT]
  
```

Every node of this list record the following information:

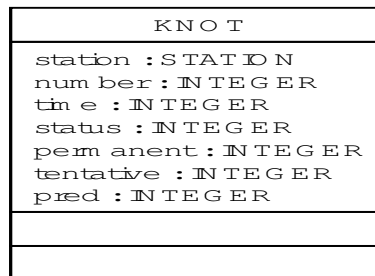


Figure 4.8: The attributes of class KNOT

```

feature {NONE}
KNOT... station: STATION
  
```

Form the start down to the destination, as long as the station is found as a valid neighbor, it will be set in a KNOT object by the following method.

```

feature {PRIME_FINDER}
  KNOT...set_station(value: STATION) is
    do
      station := value
    end

```

Of course, class KNOT requires PRIME_FINDER object give a non Void value.

```

feature {PRIME_FINDER}
  KNOT...get_station: STATION is
    do
      Result := station
    end

```

After searching, PRIME_FINDER object can get the record of station by the above method.

```

feature {NONE}
  KNOT...number: INTEGER

```

The number of KNOT object keeps the route number of the station and is set by the following method:

```

feature {PRIME_FINDER}
  KNOT...set_number(value: INTEGER) is
    require
      value >= 0
      value <= 999
    do
      number := value
    end

```

According to the requirement that train route number is a two-digit number and bus route number is a three-digit number, we set a precondition like that for this method.

```

feature {PRIME_FINDER}
  KNOT...get_number: INTEGER is
    do
      Result := number
    end

```

The above method can tell PRIME_FINDER object the route, to which this station belongs.

```

feature {NONE}
  KNOT...time: INTEGER

```

Member time records the total time needed from start and is set by the following method

```

feature {PRIME_FINDER}
  KNOT...set_time(value: INTEGER) is
    require
      value >= 0
    do
      time := value
    end

```

The time of start node is 0 and the time of destination is desire time plus one, so here KNOT object requires a nonnegative number.

```

feature {PRIME_FINDER}
  KNOT...get_time: INTEGER is
    do
      Result := time
    end

```

The above method is used to provide time for PRIME_FINDER object.

```

feature {NONE}
  KNOT...pred: INTEGER

```

This member is used to record the index of last stop in this list. The *pred* of start is -1. That the *pred* of two destination are all -1 means that there is no possible connection between the start and the destination.

PRIME_FINDER object set this member by the following method:

```

feature {PRIME_FINDER}
  KNOT...set_pred(value: INTEGER) is
    do
      pred := value
    end

```

and get the value of this member by the following method:

```

feature {PRIME_FINDER}
  KNOT...get_pred: INTEGER is
    do
      Result := pred
    end

```

Then, how can we judge if this node should be check for new neighbors? we define the member status in class KNOT.

```

feature {NONE}
  KNOT...status: INTEGER

```

If there is no more new neighbors can be found for the current station, this member should be set as permanent, which is a constant of class KNOT;

```

feature {PRIME_FINDER}
  KNOT...permanent: INTEGER is 1

```

otherwise, member status should be set as tentative, which is another constant of class KNOT.

```
feature {PRIME_FINDER}
  KNOT...tentative: INTEGER is 0
```

This member can be set by the following method

```
feature {PRIME_FINDER}
  KNOT...set_status(value: INTEGER is
    require
      value >= tentative
      value <= permanent
    do
      status := value
    end
```

and get by the following method

```
feature {PRIME_FINDER}
  KNOT...get_status: INTEGER is
    do
      Result := status
    end
```

Method make is the creation of class KNOT

```
creation {PRIME_FINDER}
  KNOT...make
```

and its main task is to initialize this object with the given parameters as following:

```
feature {PRIME_FINDER}
  KNOT...make(sn: STATION; num, t, ss, pr: INTEGER is
    do
      set_station(sn)
      set_number(num)
      set_time(t)
      set_status(ss)
      set_pred(pr)
    end
```

Every node is added into the list by the following method:

```
feature {NONE}
  PRIME_FINDER...add_node(pr: INTEGER; s: STATION; t, num: INTEGER is
    require
      t >= 0
    local
      node: KNOT
    do
      create node.make(s, num, t, node.tentative, pr)
      if s = Void then
        node.set_status(node.permanent)
      end
      stop_list.add_last(node)
    end
```

If the station is Void, then the new node will be considered as dead.

The logic of possible connection finding is implemented mainly in the following method.

```

feature {HPTA_TRANSIT_INFO}
PRIME_FINDER...
get_connection(dbase: DATABASE; start, destination: STRING; time: INTEGER): STRING is
  require else
    stop_list.upper = 0
  local
    connection, cur_station: STRING
    node: KNOT
    i, monitor: INTEGER
    is_end, break: BOOLEAN
  do
    connection := ""
    desire_time := time
    add_bus_train_station(dbase, destination, desire_time+1)
    add_bus_train_station(dbase, start, 0)

    i := 3
    cur_station := start.twin

  from
  until is_end or else cur_station = Void
  loop
    monitor := stop_list.upper
    find_neighbor(dbase, cur_station, i)
    if monitor = stop_list.upper then
      if stop_list.item(i) /= Void then
        stop_list.item(i) .set_status(node.permanent)
      end
    end
    is_end := True

  from
  until break or else i > stop_list.upper
  loop
    if stop_list.item(i) /= Void then
      node := stop_list.item(i)
      if node.get_status = node.tentative and node.get_station /= Void then
        cur_station := node.get_station.get_name
        is_end := False
        break := True
      end
    end
    if not break then
      i := i + 1
    end
  end

  if break then
    break := False
  end
end

connection := get_connection_mes(1)
connection := connection + get_connection_mes(2)

if connection.same_as("") then
  connection := "There_is_no_connection_from_your_start"
  + "_to_your_destination_in_such_time."
end

Result := connection
ensure
Result /= Void
end

```

The first parameter provides the source of data; the second and third parameters are the names of start station and destination station respectively; the last parameter is the desire time, which will be used to set the private member `desire_time`:

```
feature {NONE}
  PRIME_FINDER... desire_time: INTEGER
```

At the beginning of searching, we initialize the `stop_list` of a `PRIME_FINDER` object with four nodes, i.e. bus and train stations of destination followed by bus and train stations of start, using the following method:

```
feature {NONE}
  PRIME_FINDER... add_bus_train_station(dbase: DATABASE; name: STRING; time: INTEGER) is
  require
    name /= Void
    time >= 0
  local
    route: ROUTE
    station: STATION
    is_end: BOOLEAN
    num: INTEGER
  do
    route := dbase.get_first_bus_route
    from
    until is_end or route = Void
    loop
      station := route.get_first_station
      from
      until is_end or station = Void
      loop
        if name.same_as(station.get_name) then
          is_end := True
        end
        if not is_end then
          station := route.get_next_station(station)
        end
        if not is_end then
          route := dbase.get_next_bus_route(route)
        end
      end
    end
    if not is_end then
      station := Void
    end
    if route /= Void then
      num := route.get_number
    else
      num := 0
    end
    add_node(-1, station, time, num)

    station := Void
    is_end := False
    route := dbase.get_first_train_route
    from
    until is_end or route = Void
    loop
      station := route.get_first_station
      from
      until is_end or station = Void
```

```

    loop
      if name.same_as(station.get_name) then
        is_end := True
      end
      if not is_end then
        station := route.get_next_station(station)
      end
    end
    if not is_end then
      route := dbase.get_next_train_route(route)
    end
    if not is_end then
      station := Void
    end
    if route /= Void then
      num := route.get_number
    else
      num := 0
    end
    add_node(-1, station, time, num)
  end
end

```

Then from the bus station of start, we try to find its direct neighbor by the following method:

```

feature {NONE}
  PRIME_FINDER... find_neighbor(dbase: DATABASE; sn: STRING; pr: INTEGER) is
    require
      sn /= Void
    local
      cost, index, switch: INTEGER
      p_node, node: KNOT
      route: ROUTE
      station, last: STATION
      name: STRING
      break: BOOLEAN
    do
      from switch := 0
      until switch > 1
      loop
        if switch = 0 then
          cost := bus_time
        else
          cost := train_time
        end
        if pr >= stop_list.lower and pr <= stop_list.upper then
          p_node := stop_list.item(pr)
        end
        if p_node /= Void then
          if p_node.get_station /= Void then
            if switch = 0 then
              if is_train(p_node.get_number) then
                cost := change_time + cost
              end
              route := dbase.get_first_bus_route
            else
              if is_bus(p_node.get_number) then
                cost := change_time + cost
              end
              route := dbase.get_first_train_route
            end
          from
          until route = Void
          loop
            station := route.get_first_station
          end
        end
      end
    end

```

```

last := station
from
until station = Void or break
loop
name := station.get_name.twin
if name /= Void and name.is_equal(sn) then
if not last.get_name.is_equal(name) then
index := get_index(last, route.get_number)
if index >= 0 then
node := stop_list.item(index)
if node.get_station /= Void then
if is_train(node.get_number) then
if node.get_time > p_node.get_time + cost then
node.set_pred(pr)
node.set_time(p_node.get_time + cost)
node.set_number(route.get_number)
end
end
end
else
add_node(pr, last, p_node.get_time+cost, route.get_number)
end
end
last := route.get_next_station(station)
if last /= Void then
index := get_index(last, route.get_number)
if index >= 0 then
node := stop_list.item(index)
if node.get_station /= Void then
if is_train(node.get_number) then
if node.get_time > p_node.get_time + cost then
node.set_pred(pr)
node.set_time(p_node.get_time + cost)
node.set_number(route.get_number)
end
end
end
else
add_node(pr, last, p_node.get_time+cost, route.get_number)
end
end
break := True
else
last := station;
station := route.get_next_station(station)
end
end
break := False
if switch = 0 then
route := dbase.get_next_bus_route(route)
else
route := dbase.get_next_train_route(route)
end
end
end
end
switch := switch + 1
end
end

```

For convenience, we define the following two methods to tell if the current route is train or bus:

```

feature {NONE}
PRIME.FINDER... is_train(num: INTEGER): BOOLEAN is
do

```

```

Result := num >= 10 and num <= 99
end

```

and

```

feature {NONE}
  PRIME.FINDER... is_bus(num: INTEGER): BOOLEAN is
  do
    Result := num >= 100 and num <= 999
  end

```

The following method is used to get the index of a certain station in the list; if the target station is not in the list, -1 will be return.

```

feature {NONE}
  PRIME.FINDER... get_index(s: STATION; num: INTEGER): INTEGER is
  require
    s /= Void

  local
    ind, i: INTEGER
    node: KNOT
    name: STRING

  do
    ind := -1
    from i := stop_list.lower
    until i > stop_list.upper
    loop
      node := stop_list.item(i)
      if node.get_station /= Void then
        name := node.get_station.get_name
        if name.is_equal(s.get_name) then
          if is_bus(num) and is_bus(node.get_number) then
            ind := i
          elseif is_train(num) and is_train(node.get_number) then
            ind := i
          end
        end
      end
      i := i + 1
    end
  end
  Result := ind
end

```

When the searching is done, we can get the information of possible connections by the following method:

```

feature {NONE}
  PRIME.FINDER... get_connection_mes(index: INTEGER): STRING is
  require
    index >= 0

  local
    node: KNOT
    mes: STRING

  do
    mes := ""
    node := stop_list.item(index)
    if node /= Void then
      if node.get_station /= Void then
        if node.get_pred /= -1 and node.get_time <= desire_time then
          mes := "-No." + node.get_number.to_string + "->"
            + node.get_station.get_name + "_in_"
            + node.get_time.to_string + "_minutes%N"
        end
      end
    end
  end

```

```

node := stop_list.item(node.get_pred)
from
until node = Void or else node.get_station = Void or else node.get_pred = -1
loop
mes := "-No." + node.get_number.to_string + "->"
      + node.get_station.get_name + mes
node := stop_list.item(node.get_pred)
end
if node /= Void then
if node.get_station /= Void then
mes := "%N" + node.get_station.get_name + mes
end
else
mes := ""
end
end
end
Result := mes
end

```

The creation of PRIME_FINDER is method make

```
creation {ANY} PRIME.FINDER... make
```

it is defined as following:

```
feature {HPTA_TRANSIT_INFO}
PRIME_FINDER... make is
do
create stop_list.make
ensure
stop_list /= Void
end

```

Now, let us talk about the root class HPTA_TRANSIT_INFO.

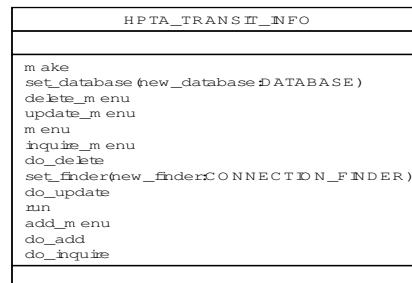


Figure 4.9: The methods of class HPTA_TRANSIT_INFO

The creation of class HPTA_TRANSIT_INFO is make

```
create HPTA_TRANSIT_INFO... make
```

Its main task is to initialize the database and connection finder, and then run the whole system:

```
feature {ANY}
  HPTA_TRANSIT_INFO...make is
    local
      prime_finder: PRIME_FINDER
      file_database: FILE_DATABASE
    do
      create file_database.make
      set_database(file_database)
      create prime_finder.make
      set_finder(prime_finder)
      run
    end
```

In order to increase customer satisfaction, we run the system by a series of menus

```
feature {NONE}
  HPTA_TRANSIT_INFO...run is
    do
      from
      until io.last_character.to_upper = 'Q'
      loop
        menu
        io.read_character
        io.put_new_line
        inspect io.last_character.to_upper
        when 'U' then do_update
        when 'I' then do_inquire
        else
        end
      end
    end
```

In order to use OS command, we let class HPTA_TRANSIT_INFO be a subclass of class SYSTEM, which is a predefined class in Eiffel.

```
inherit
  HPTA_TRANSIT_INFO...SYSTEM
```

Method menu is the main menu of the interface of this system and

```
feature {NONE}
  HPTA_TRANSIT_INFO...menu is
    do
      execute_command_line("cls")
      io.put_string("[
*****
-----Welcome_to_HPTA
*****
-----U_Update_System_Information
-----I_Inquire_about_Transit_Information
-----Q_Quit
-----Enter_menu_choice:
-----]")
    end
```

This is the main menu and there are two items in it, through which users can either update or inquire system information. The first line of the method body is

used to clear the screen.

If users chose the first menu item, they are going to enter the following menu, i.e. update_menu:

```
feature {NONE}
  HPTA_TRANSIT_INFO... update_menu is
    do
      execute_command_line("cls")
      io.put_string("[
-----*****
-----Welcome_to_HPTA
-----*****
-----A_Add
-----D_Delete
-----G_Go_back

-----Enter_menu_choice:
-----]")
    end
```

In this menu, users can add new information, such as staffs and stations, as follow:

```
feature {NONE}
  HPTA_TRANSIT_INFO... add_menu is
    do
      execute_command_line("cls")
      io.put_string("[
-----*****
-----Welcome_to_HPTA
-----*****
-----S_Add_a_station
-----E_Add_a_staff
-----G_Go_back

-----Enter_menu_choice:
-----]")
    end
```

Follows the logic of method do_add:

```
feature {NONE}
  HPTA_TRANSIT_INFO... do_add is
    local
      employee: STAFF
      id: INTEGER
      input, name, password, open, access, last: STRING
      is_end: BOOLEAN
      route: ROUTE
    do
      from
      until is_end
      loop
        add_menu
        io.read_line
        input := io.last_string.twin
        io.put_new_line
        if not input.is_empty then
          inspect input.first.to_upper
          when 'G' then is_end := True
          when 'S' then
            io.put_string("%NEnter_station_name:_")
            io.read_line
            name := io.last_string.twin
            io.put_string("%NEnter_open_hour:_")
```

```

        io.read_line
        open := io.last_string.twin
        io.put_string("%NEnter_its_accessibility:_")
        io.read_line
        access := io.last_string.twin
        io.put_string("%NEnter_route_number:_")
        io.read_line
        id := io.last_string.to_integer
        io.put_string("%NEnter_the_name_of_its_last_station:_")
        io.read_line
        last := io.last_string.twin
        route := db.find_route(id)
        if route = Void then
            create route.make(id)
            route.add_station(name, access, open, last)
            db.add_route(route)
        else
            route.add_station(name, access, open, last)
        end
    when 'E' then
        io.put_string("%NEnter_your_ID:_")
        io.read_line
        id := io.last_string.to_integer
        io.put_string("%NEnter_your_password:_")
        io.read_line
        password := io.last_string.twin
        create employee.make(id, password)
        db.add_staff(employee)
    else
    end
end
end
end
end

```

they can also delete those information as follow:

```

feature {NONE}
    HPTA_TRANSIT_INFO... delete_menu is
    do
        execute_command_line("cls")
        io.put_string(" [
        *****
        -----Welcome_to_HPTA
        *****
        -----S_Delete_a_station
        -----E_Delete_a_staff
        -----R_Delete_a_route
        -----G_Go_back
        -----Enter_menu_choice :
        -----] ")
    end

```

Follows the logic of method do_delete:

```

feature {NONE}
    HPTA_TRANSIT_INFO... do_delete is
    local
        is_end: BOOLEAN
        num: INTEGER
        staff: STAFF
        route: ROUTE
        input, name: STRING
    do
        from
        until is_end

```

```

loop
  delete_menu
  io.read_line
  input := io.last_string.twin
  io.put_new_line
  if not input.is_empty then
    inspect input.first.to_upper
    when 'G' then is_end := True
    when 'S' then
      io.put_string ("%NEnter_route_number:_")
      io.read_line
      num := io.last_string.to_integer
      io.put_string ("%NEnter_station_name:_")
      io.read_line
      name := io.last_string.twin
      route := db.find_route(num)
      if route /= Void then
        route.remove_station(name)
      else
        io.put_string ("%NNo_such_a_station%N")
        io.read_line
      end
    when 'R' then
      io.put_string ("%NEnter_route_number:_")
      io.read_line
      num := io.last_string.to_integer
      route := db.find_route(num)
      if route /= Void then
        db.delete_route(route)
      else
        io.put_string ("%NNo_such_a_station%N")
        io.read_line
      end
    when 'E' then
      io.put_string ("%NEnter_ID:_")
      io.read_line
      num := io.last_string.to_integer
      staff := db.find_staff(num)
      if staff /= Void then
        db.delete_staff(staff)
      else
        io.put_string ("%NNo_such_a_staff%N")
        io.read_line
      end
    else
      end
    end
  end
end
end
end

```

According to the requirement, only authorized staffs can do such things, so this system will ask the user to log in the system before he or she enter the update menu. The following method `do_update` has the logic to require the user to enter his or her employee number and password first.

```

feature {NONE}
  HPTA_TRANSIT_INFO... do_update is
    local
      id: INTEGER
      passed, is_end: BOOLEAN
      password, input: STRING
      staff: STAFF
    do
      io.read_line

```

```

    if db.is_locked then
        io.put_string ("%NEnter_employee_ID:")
        io.read_line
        id := io.last_string.to_integer
        staff := db.find_staff(id)
        if staff /= Void then
            io.put_string ("%NEnter_password:")
            io.read_line
            password := io.last_string.twin
            passed := staff.login(password)
        end
    else
        io.put_string (" [
        -----The_list_of_authorized_staff_is_not_empty,
        -----so_please_set_authorization_as_soon_as_possible...
        -----] ")
        passed := True
        io.read_line
    end
    if passed then
        from
        until is_end
        loop
            update_menu
            io.read_line
            input := io.last_string.twin
            io.put_new_line
            if not input.is_empty then
                inspect input.first.to_upper
                when 'A' then do_add
                when 'D' then do_delete
                when 'G' then is_end := True
            else
                end
            end
        end
        db.do_save
    else
        io.put_string ("%NLogin_failed!%N")
        io.read_line
    end
end
end

```

The actual logging responsibility is assigned to class STAFF as public feature to class HPTA_TRANSIT_INFO:

```

feature {HPTA_TRANSIT_INFO}
STAFF... login (passwd: STRING): BOOLEAN is
    require
        passwd /= Void
    do
        Result := password.is_equal (passwd)
    end
end

```

If the result is True, the user can continue his or her update, otherwise, this system will remain on the main menu.

If users chose the second menu item of the main menu, they will enter the following query menu without any bother, because the requirement says that any one can have access to the transit information.

```

feature {NONE}
  HPTA_TRANSIT_INFO... inquire_menu is
    do
      execute_command_line(" cls")
      io.put_string(" [
*****
*****Welcome_to_HPTA
*****
*****F_Find_a_possible_connection
*****
*****S_Show_a_route
*****B_Browse_all_routes
*****G_Go_back
*****
*****Enter_menu_choice:
*****] ")
    end

```

The first item of this menu is used for users to find a possible connection. Following the logic of method `do_inquire`, users are required to enter their start, destination, as well as their desire time.

```

feature {NONE}
  HPTA_TRANSIT_INFO... do_inquire is
    local
      input, start, dest: STRING
      is_end: BOOLEAN
      num, time: INTEGER
      route: ROUTE
    do
      from
        until is_end
        loop
          inquire_menu
          io.read_line
          input := io.last_string.twin
          io.put_new_line
          if not input.is_empty then
            inspect input.first.to_upper
            when 'B' then
              from route := db.get_first_bus_route
              until route = Void
              loop
                route.show
                route := db.get_next_bus_route(route)
              end
            from route := db.get_first_train_route
            until route = Void
            loop
              route.show
              route := db.get_next_train_route(route)
            end
          io.put_string ("%N%NStrike_any_key_to_continue...")
          io.read_line
          when 'F' then
            io.put_string ("%NEnter_the_station_name_of_your_start:_")
            io.read_line
            start := io.last_string.twin
            io.put_string ("%NEnter_the_station_name_of_your_destination:_")
            io.read_line
            dest := io.last_string.twin
            io.put_string ("%NEnter_your_desire_time(in_minutes):_")
            io.read_line
            time := io.last_string.to_integer
            io.put_string (finder.get_connection(db, start, dest, time))
          end
        end
      end

```

```

io.put_string ("%N%NStrike_any_key_to_continue...")
io.read_line
when 'S' then
io.put_string ("Input_the_route_number_(10--999):_")
io.read_line
num := io.last_string.to_integer
route := db.find_route(num)
if route /= Void then
route.show
else
io.put_string ("Sorry_there_is_no_such_a_route")
end
io.put_string ("%N%NStrike_any_key_to_continue...")
io.read_line
when 'G' then is_end := True
else
end
end
end
end
end

```

Now, it is time to implement the methods of class ROUTE

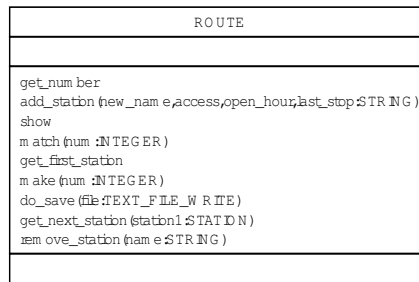


Figure 4.10: The methods of class ROUTE

The creation of ROUTE is make, which can be invoke by class HPTA_TRANSIT_INFO

```

creation ROUTE... make

```

The main task of make is initialize the route number and station list

```

feature {HPTA_TRANSIT_INFO}
ROUTE... make (num: INTEGER) is
require
num > 9
num < 1000
do
number := num
create stops.make
ensure
number = num
stops /= Void
end

```

According to the requirement, route number must be two- or three-digit number, so we define the following invariant for class ROUTE.

```
invariant
    ROUTE... number > 9
    number < 1000
end
```

At any time, its client gets route number by the following method:

```
feature {ANY}
    ROUTE... get_number : INTEGER is
    do
        Result := number
    end
```

also, by the following method to tell if the current route is which we want:

```
feature {ANY}
    ROUTE... match (num: INTEGER): BOOLEAN is
    do
        Result := num = number
    end
```

By the following method, its client adds new stations for this ROUTE object and at the same time sets the name, the accessibility, the opening hour, and last station for this new station.

```
feature {HPTA_TRANSIT_INFO, DATABASE}
    ROUTE... add_station(new_name, access, open_hour, last_stop: STRING) is
    local
        new_station: STATION
        i: INTEGER
        last: STRING
    do
        from i := stops.lower
        until i > stops.upper or else stops.item(i) .match(new_name)
        loop
            i := i + 1
        end
        if i > stops.upper then
            create new_station.make(new_name, access, open_hour)
            last := last_stop.twin
            last.to_upper
            if last.same_as("NONE") then
                stops.add_first(new_station)
            else
                from i := stops.lower
                until i > stops.upper or else stops.item(i) .match(last_stop)
                loop
                    i := i + 1
                end
                if i <= stops.upper then
                    stops.add(new_station, i+1)
                else
                    create new_station.make(last_stop, access, open_hour)
                    stops.add_last(new_station)
                    stops.add_last(new_station)
                end
            end
        end
    end
end
```

HPTA_TRANSIT_INFO object removes a certain station by the following method, whose only parameter is the name of the target station.

```

feature {HPTA_TRANSIT_INFO}
ROUTE... remove_station(name: STRING) is
  local
    i: INTEGER
  do
    from i := stops.lower
    until i > stops.upper or else stops.item(i) .match(name)
    loop
      i := i + 1
    end
    if i <= stops.upper then
      stops.remove(i)
    end
  end
end

```

The subclasses of CONNECTION_FINDER use the following two methods to visit all stations in this route

```

feature {CONNECTION_FINDER}
ROUTE... get_first_station: STATION is
  local
    station: STATION
  do
    if stops.upper > 0 then
      station := stops.first
    end
    Result := station
  end
get_next_station(station1: STATION): STATION is
  require
    station1 /= Void
  local
    station: STATION
  do
    if stops.index_of(station1) < stops.upper then
      station := stops.item(stops.index_of(station1)+1)
    end
    Result := station
  end
end

```

Class ROUTE keep the secret of saving itself, so DATABASE object can call this method to fulfill the task. Actually, such assignment is worth to discuss. Maybe should move to the subclasses of DATABASE, because only they know exactly how to save those data.

```

feature {DATABASE}
ROUTE... do_save(file: TEXT_FILE_WRITE) is
  require
    file.is_connected
  local
    i: INTEGER
    tag, last: STRING
  do
    if number > 99 then
      tag := "b"
    else
      tag := "t"
    end
  end
end

```

```

last := "None"
from i := stops.lower
until i > stops.upper
loop
  file.put_string(tag+"%N")
  stops.item(i) .do_save(file)
  file.put_string("_" + number.to_string + "_" + last + "%N")
  last := stops.item(i) .get_name.twin
  i := i + 1
end
end

```

Similarly, the following method is responsible for showing the details of this route, but only class HPTA_TRANSIT_INFO know exactly how to display with interface, so this method should be move to class HPTA_TRANSIT_INFO.

```

feature {HPTA_TRANSIT_INFO}
ROUTE... show is
  local i: INTEGER
  do
    if number > 99 then
      io.put_string ("%NBus_route_No.")
    else
      io.put_string ("%NTrain_route_No.")
    end
    io.put_integer (number)
    io.put_string (":_")
    from i := stops.lower
    until i > stops.upper
    loop
      stops.item(i) .show
      if i < stops.upper then
        io.put_string (">")
      end
      i := i+1;
    end
    io.put_new_line
  end
end

```

Same problem can be found on the method show of class STATION

```

feature {ROUTE}
STATION... show is
  do
    io.put_string (name)
  end
end

```

Now, let us look at the class STATION, whose creation is method make too,

```

create STATION... make

```

and defined as following:

```

feature {ROUTE}
STATION... make (new_name, new_open, new_acc: STRING) is
  require
    new_name /= Void
    new_open /= Void
    new_acc /= Void
  do
    name := new_name.twin
    open := new_open.twin
  end
end

```

```

accessibility := new_acc.twin
end

```

The main task of it is to initial these three features of class STATION. At any time, its client can visit these three features by the following methods:

```

feature {ROUTE, CONNECTION_FINDER}
STATION...get_name: STRING is
  do
    Result := name.twin
  end
get_acc: STRING is
  do
    Result := accessibility.twin
  end
get_open: STRING is
  do
    Result := open.twin
  end
end

```

Similar with the method do_save of class ROUTE, this method should be moved into the subclasses of DATABASE.

```

feature {ROUTE}
STATION...do_save(file: TEXT_FILE_WRITE) is
  require
    file.is_connected
  do
    file.put_string(name + "\n" + accessibility + "\n" + open)
  end
end

```

The same problem can be found on class STAFF

```

feature {DATABASE}
STAFF...do_save(file: TEXT_FILE_WRITE) is
  require
    file.is_connected
  do
    file.put_string("%s\n" + number.to_string + "\n" + password + "%N")
  end
end

```

We identify station with name only, i.e. if two stations share the same name, we assume they are the same station. Here case is insensitive.

```

feature {ROUTE}
STATION...match (targetname: STRING): BOOLEAN is
  require
    targetname /= Void
  do
    Result := name.same_as (targetname)
  end
end

```

Now, let us talk about the implementation of class STAFF.

The creation of class STAFF is make

```

creation{ANY} STAFF...make

```

it is defined as following:

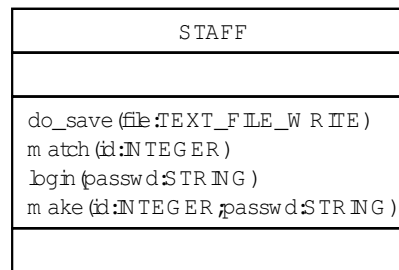


Figure 4.11: The methods of class STAFF

```
feature {ANY}
  STAFF...make (id: INTEGER; passwd: STRING) is
    require
      id >= 0
      passwd /=Void
    do
      number := id;
      password := passwd.twin
    ensure
      number >= 0
      password = passwd
    end
```

its main task is initialize staff's id and password.

Method match is used to identify a certain staff and is defined as following:

```
feature {ANY}
  STAFF...match (id: INTEGER): BOOLEAN is
    do
      Result := id = number
    end
```

Any staff has an unique employee number, which is generated from 0, and a password, which must not be Void:

```
invariant
  STAFF...number_positive: number >= 0
  password_not_void: password /= Void
end
```

4.1.9 Automatic Code Listing

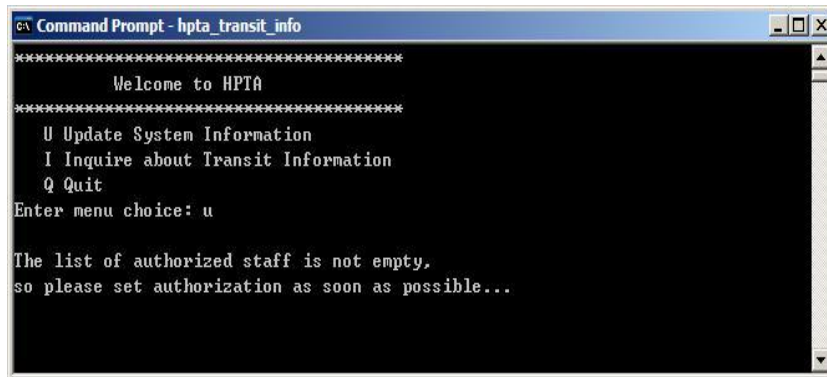
So far, we have implement the system. In order to give an integrated view for ones who are used to read code, Spark inserts all program code here automatically.

Automatical code is listed in Appendix C.

4.1.10 Testing

Updating system

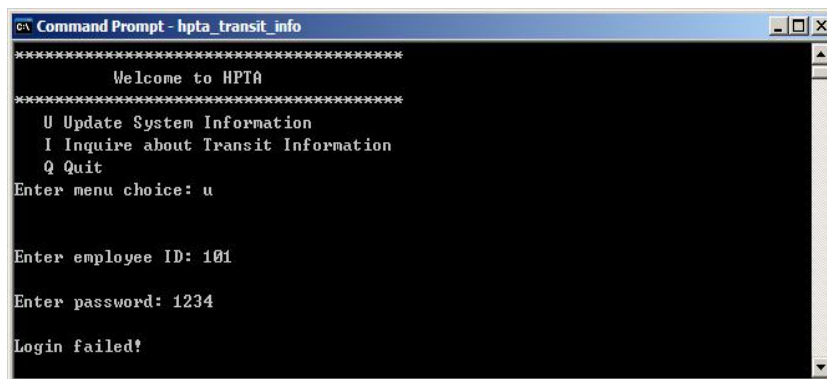
When no staff is authorized, we try to update system information. The result is



```
c:\ Command Prompt - hpta_transit_info
*****
      Welcome to HPTA
*****
  U Update System Information
  I Inquire about Transit Information
  Q Quit
Enter menu choice: u

The list of authorized staff is not empty,
so please set authorization as soon as possible...
```

Otherwise, we try to update system information. The system requires ID and password for logging in as following:



```
c:\ Command Prompt - hpta_transit_info
*****
      Welcome to HPTA
*****
  U Update System Information
  I Inquire about Transit Information
  Q Quit
Enter menu choice: u

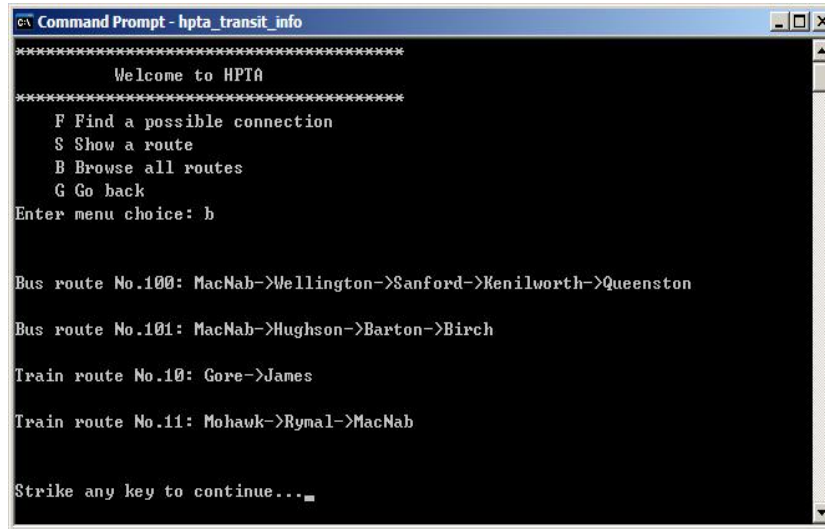
Enter employee ID: 101
Enter password: 1234
Login failed!
```

These results satisfy the design requirements.

Browsing all routes

We try to browse the information of all routes as following:

These result satisfies the design requirements.



```
c:\ Command Prompt - hpta_transit_info
*****
Welcome to HPTA
*****
F Find a possible connection
S Show a route
B Browse all routes
G Go back
Enter menu choice: b

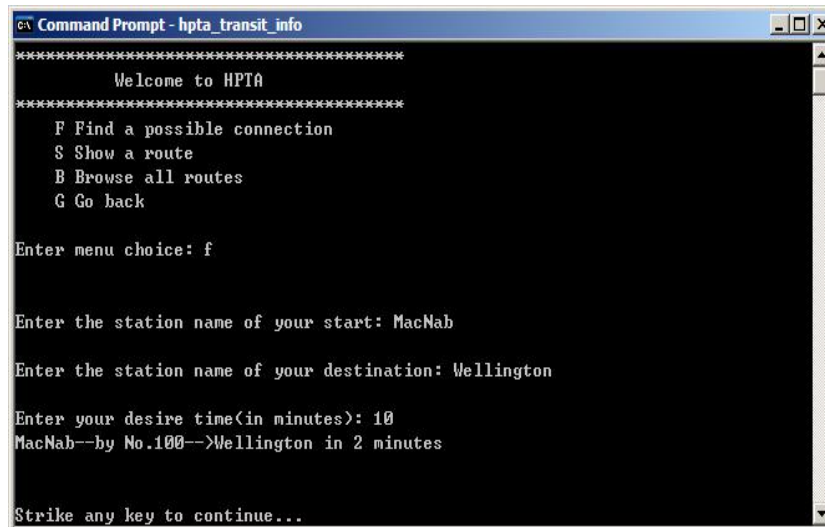
Bus route No.100: MacNab->Wellington->Sanford->Kenilworth->Queenston
Bus route No.101: MacNab->Hughson->Barton->Birch

Train route No.10: Gore->James
Train route No.11: Mohawk->Rymal->MacNab

Strike any key to continue...
```

Finding connection

We try to find a connection between two stations as following:



```
c:\ Command Prompt - hpta_transit_info
*****
Welcome to HPTA
*****
F Find a possible connection
S Show a route
B Browse all routes
G Go back
Enter menu choice: f

Enter the station name of your start: MacNab
Enter the station name of your destination: Wellington

Enter your desire time(in minutes): 10
MacNab--by No.100-->Wellington in 2 minutes

Strike any key to continue...
```

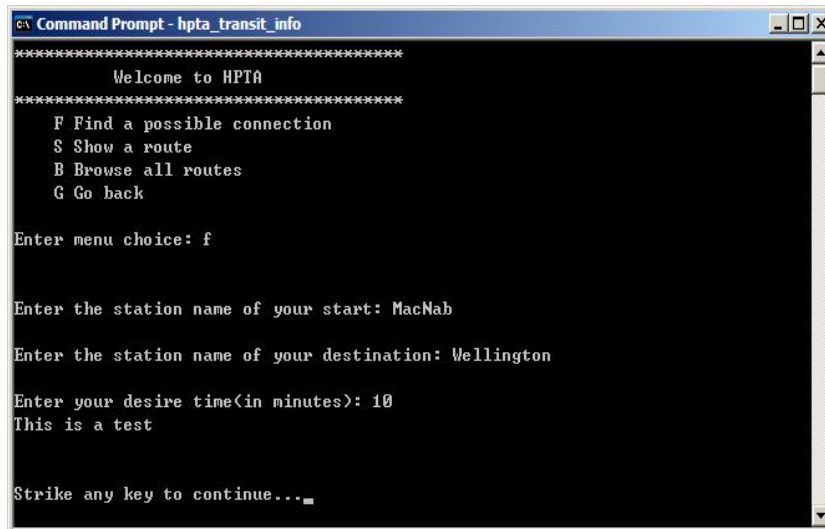
These result satisfies the design requirements.

Strategy pattern

We construct a sample connection finder class and change the algorithm at run-time.

```
class FAKE.FINDER
inherit
    CONNECTION.FINDER
creation {ANY}
make
feature {HPTA.TRANSIT.INFO}
make is
do
end
get_connection(start, destination: STRING; time: INTEGER): STRING is
do
    Result := "This_is_a_test%N";
end
end
```

The result is



```
c:\ Command Prompt - hpta_transit_info
*****
Welcome to HPTA
*****
F Find a possible connection
S Show a route
B Browse all routes
G Go back

Enter menu choice: f

Enter the station name of your start: MacNab

Enter the station name of your destination: Wellington

Enter your desire time(in minutes): 10
This is a test

Strike any key to continue..._
```

These result satisfies the design requirements.

Chapter 5

Implementation

Chapter 3 has showed us the key features of OOLP and some supporting tools. In this chapter, we present the implementation of Spark in a literate way using Spark itself. The rest of this chapter that follow is the actual output of Spark source file.

5.1 Introduction

The two reasons why we present the implementation of Spark in a literate way are that we want to show the universality of Spark, i.e. it can work with not only object-oriented programming languages, but also structured programming languages, and that since the main task of Spark is to parse the syntax of a certain programming language, which is a fairly stable structure, structured programming makes the program clear, simple and efficient. This choice is also followed by one drawback that we have to draw the diagrams by hand.

Spark is implemented entirely in perl. We choose perl mainly because it is good at text manipulation, which is the main task of Spark, and perl is a stable, cross platform programming language, which leads to Spark being inherently platform-independent.

In order to gain more flexibility, we separate Spark into two parts, i.e., front end and back end (see Figure 5.1). The front end is responsible for explaining graphical notation settings and parsing program code chunks; the back end takes care of producing graphical notation files. So far, we have developed three front ends, which are used for Eiffel, Lime, and perl itself respectively.

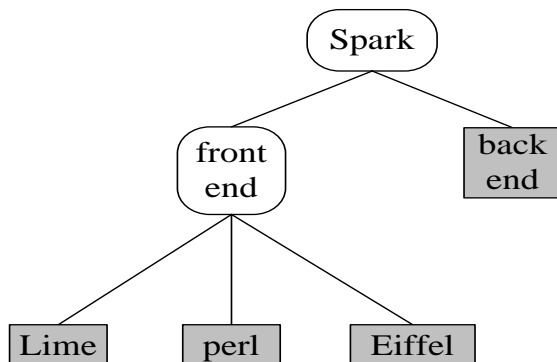


Figure 5.1: Top-level structure for Spark

5.2 Graphic Notation describer

The graphic notation describer is an interim file used to describe all the graphic notations included in the software documentation and the only input file of the back end of Spark. So the changes coming from either document formatting language or programming language do not impact the back end. In addition, except several important setting tags, programmers do not know anything about it, because this file is going to be deleted by the back end before it finishes its work.

In the following table of graphic notation describer structure, terminals are shown in bold font and nonterminals in italics, parentheses ‘(’ and ‘)’ indicate grouping when needed, square brackets ‘[’ and ‘]’ enclose optional items, curly braces ‘{’ and ‘}’ show the (zero or more) repeatable items, and vertical bars ‘|’ separate alternatives.

```

modFile ::= { classDiagram } { class }
classDiagram ::= @CLASSDIAGRAM nameSequence [ @VERTICAL ] [format]
format ::= @BRIEF | @CONCISE | @METHOD | @ATTRIBUTE | @ACTION
class ::= @CLASS name { statementSequence }
           | @INTERFACE name { statementSequence }
statementSequence ::= statement statementSequence
statement ::= extendStmt
              | inheritStmt
              | implementStmt
              | fieldStmt
              | contrStmt
  
```

<i>methodStmt</i>
<i>actionStmt</i>
<i>dependentStmt</i>
<i>extendStmt</i> ::= @EXTEND <i>nameSequence</i>
<i>inheritStmt</i> ::= @INHERIT <i>nameSequence</i>
<i>implementStmt</i> ::= @IMPLEMENT <i>nameSequence</i>
<i>fieldStmt</i> ::= @VAR <i>name</i> { <i>modifier</i> }
<i>cotrStmt</i> ::= @INIT <i>name</i> { <i>modifier</i> }
<i>methodStmt</i> ::= @METHOD <i>name</i> { <i>modifier</i> }
<i>actionStmt</i> ::= @ACTION <i>name</i> { <i>modifier</i> }
<i>nameSequence</i> ::= <i>name nameSequence</i>
<i>modifier</i> ::= <i>visibility</i> <i>type</i> '(' <i>paraSequence</i> ')'
<i>visibility</i> ::= @PRIVATE @PUBLIC @PROTECTED
<i>paraSequence</i> ::= <i>parameter; paraSequence</i>
<i>parameter</i> ::= <i>name</i> ':' <i>type</i>
<i>type</i> ::= <i>name</i>
<i>name</i> ::= <i>letter</i> { <i>letter</i> <i>digital</i> }
<i>letter</i> ::= a b c d e f g h i j k l m n o p q r s t
u v w x y z A B C D E F G H I J D L
M N O P Q R S T U V W X Y Z
<i>digital</i> ::= 1 2 3 4 5 6 7 8 9 0

Table 5.1: The block structure of graphic notation describer.

A graphic notation describer consists of multiple *classDiagrams* followed by multiple *classes*. *classDiagram* begin with the key word “@CLASSDIAGRAM”. *nameSequence* is a list of class names included in this diagram. “@VERTICAL” is used to set the direction of the specified graphic notation, i.e. if “@VERTICAL” is set, the diagram will be drawn vertically, otherwise horizontally. “@HEAD”, “@BRIEF” and “@CONCISE” are used to control the format of the specified graphic notation. If “@CONCISE” is set, the class diagram will hide all the information about the method’s parameters of the involved class. If “@BRIEF” is set, the class diagram will be shown without parameters and types. If “@HEAD” is set, the class diagram will be shown with class name only. Ones also can use “@METHOD”, “@ATTRIBUTE”, and “@ACTION” to control the display areas. For example, if “@METHOD” is set, all class methods will be shown in the diagram and if “@ATTRIBUTE” is set, only attributes of class can be saw in the diagram. All the setting tags are included in

source file as specific comments, so they are transparent for everything except the front end. *class* begin with the key words “@CLASS” or “@INTERFACE”. *class* name followed by a list of statements in the form of one statement each line. “*inheritStmt*”, “*extendStmt*” and “*implementStmt*” begin respectively with the key words “@INHERIT”, “@EXTEND”, and “@IMPLEMENT”, which are followed by a list of superclass names. “*fieldStmt*”, “*cotrStmt*”, “*methodStmt*”, and “*actionStmt*” are the member statements of this class and begin with their key word respectively.

5.3 Front End

The front end of Spark takes an OOLP program file as input and produces machine-readable program code files and one graphic notation script file as output. In addition, it can insert the copy of machine-readable program code list back into the OOLP program file upon the request. The only entrance of front end is *main*, which depends on three modules, i.e. *parseSource*, *doOutput*, and *callBackEnd* (see Figure 5.2). The rest of this section discusses the implementation of Spark front end for Lime in details.

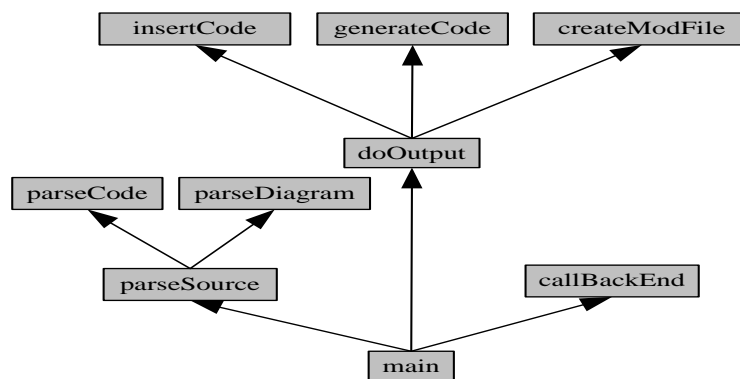


Figure 5.2: Top-level structure for Spark front end

The entrance of Spark front end for lime,

```
Sparkf-lime...& main;
```

is defined as following:

```
Sparkf-lime...sub main{
  $#ARGV == 0 || die "Usage: perl sparkf.pl filename\n";
```

```

open(SOURCE, "< $old") || die "Cannot open $old!\n";
close SOURCE || die "Can't close $old: $!";
&parseSource;
&doOutput;
&callBackEnd;
}

```

The first process in main is parsing the OOLP program, whose name is hold by a local variable \$old.

```
Sparkf-lime...my $old = $ARGV[0];
```

Module *parseSource* keeps reading the content of the program line by line:

```

Sparkf-lime...sub parseSource{
while(1){
&newLine;
last if($done);
if($currentLine =~ /image::/){
&parseDiagram;
} elsif($currentLine =~ /^-{5,}/){
&parseCode;
}
}
}
}

```

using the following function:

```

Sparkf-lime...sub newLine{
my @array = @currentWords;
push(@lastWords, \@array);
if($#lineBuffer < 0){
unless($done){
$done = 1;
while(<>){
chop;
$currentLine = $_;
@currentWords = split;
$pointer = 0;
$done = 0;
last if($#currentWords >= 0);
$done = 1;
}
}
} else{
my $refArray = pop(@lineBuffer);
@currentWords = @$refArray;
$pointer = 0;
}
}
}

```

where local variable @currentWords defined as

```
Sparkf-lime...my @currentWords = ();
```

always keeps the words of the current line in array form and local variable @lastWords defined as

```
Sparkf-lime...my @lastWords = ();
```

keeps all the addresses of old @currentWords in order and local variable \$currentLine defined as

```
Sparkf-lime...my $currentLine = "";
```

always keeps the words of the current line in string form and local variable \$done defined as

```
Sparkf-lime...my $done = 0;
```

will be set as 1 after the last line is read and local variable @lineBuffer defined as

```
Sparkf-lime...my @lineBuffer = ();
```

is used to hold the current line temporarily in the case that front end need read again the last word, which is in the last line, and the current line is still needed. The local variable \$pointer defined as

```
Sparkf-lime...my $pointer = 0;
```

is used to point out the current word the front end is reading and increases by one after the execution of function *nextWord* defined as

```
Sparkf-lime...sub nextWord{
  if($pointer >= $#currentWords){
    &newline;
  }else{
    $pointer++;
  }
  $currentWords[$pointer];
}
```

and decreases by one after the execution of function *lastWord* defined as

```
Sparkf-lime...sub lastWord{
  if($pointer == 0){
    if($#lastWords >= 0){
      my @array = @currentWords;
      push(@lineBuffer, \@array);
      my $refArray = pop(@lastWords);
      @currentWords = @$refArray;
      $pointer = $#currentWords;
    }
  }else{
    $pointer--;
  }
  $currentWords[$pointer];
}
```

If a diagram tag such as “image:” is encountered, the front end will enter the status of parsing diagrams:

```
Sparkf-lime...sub parseDiagram{
  /\w+::(\w+)\.(\w+)/;
  my $picture = $1.".".$2;
```

```

my $token = &nextWord;
if($token =~ /\s\/\s/){
  for(my $i = 1; $i <= $#currentWords; $i++){
    $picture .= " ". $currentWords[$i];
  }
} else{
  &lastWord;
}
push(@diagramList, $picture);
}

```

Front end keeps all the information concerning the current diagram in a local variable, \$picture, and then pushes it into the diagram list:

```
Sparkf-lime...my @diagramList = ();
```

If a code tag such as a serial of “-” is encountered, the front end will enter the status of parsing code, which is the main difference between different front ends:

```

Sparkf-lime...sub parseCode{
my $token = &nextWord;
until($token =~ /^-{5,}/){
  if($token eq "class"){
    parseClass(0);
  } elsif($token eq "final"){
    &nextWord;
    parseClass(1);
  } elsif($memberList{$token}){
    parseMembers($token);
  }
  $token = &nextWord;
}
}

```

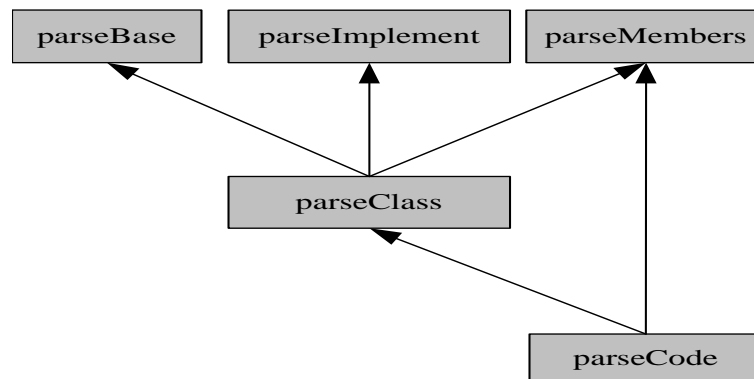


Figure 5.3: The structure of Module parseCode

The front end for lime parses various code units such as a class, a class member, and a statement, according to the syntax one by one as following:

If the first token of a code block is “final” or “class”, then this block is a class block. For class block, front end parse it according to the syntax that class ::= [“final”] “class” identifier base implement members “end” as following:

```
Sparkf-lime...sub parseClass{
  $currentClass = &nextWord;
  my $classBody = hashAdd(\%classList, $currentClass);
  $classBody->{"final"} = $_[0];
  my $token = &nextWord;
  if($token eq "inherit"){
    parseBase("inherit");
    $token = &nextWord;
  }
  if($token eq "extend"){
    parseBase("extend");
    $token = &nextWord;
  }
  while($token eq "implement"){
    $token = &nextWord;
    parseImplement($token);
    $token = &nextWord;
  }
  while($memberList{$token}){
    parseMembers($token);
    $token = &nextWord;
  }
  die "parsing class failed![$token]\n" if($token ne "end");
}
```

where the local variable `$currentClass` keeps the name of current class and is defined as:

```
Sparkf-lime...my $currentClass = "";
```

and supporting function `hashAdd` is used to add item to a hash table without duplicate and defined as:

```
Sparkf-lime...sub hashAdd{
  my($hash, $item) = @_;
  if(not $hash->{$item}){
    my %newHash = ();
    $hash->{$item} = \%newHash;
  }
  $hash->{$item};
}
```

If the first token of a code block is one of the members of

```
Sparkf-lime...my %memberList = ("public", 1, "action", 1, "template", 1, "const", 1,
                               "initialization", 1, "method", 1, "var", 1);
```

then the block is a class member block. There are six kinds of legal members, i.e. constant, variable, method, action, initialization, and label:

```
Sparkf-lime...sub parseMembers{
  my($token, $public) = ($_[0], 0);
```

```

if($token eq "public"){
    $public = 1;
    $token = &nextWord;
}
if($token eq "const"){
    parseConst(&nextWord, $public);
} elseif($token eq "var"){
    parseVariable($public, &nextWord);
} elseif(($token eq "method") || ($token eq "template")){
    parseMethod($token, $public);
} elseif($token eq "action"){
    parseAction(&nextWord);
} elseif($token eq "initialization"){
    parseInitialization(&nextWord);
} elseif($token =~ /\w+\/){
    parseLabel($token);
}
}
}

```

For constants, the syntax is
constant ::= “const” identifier [“.” type] “=” expression:

```

Sparkf-lime...sub parseConst{
    my ($token, $public) = @_;
    my $identifier = "";
    my $const = "$public.";
    if($token =~ /\w+\/){
        $identifier = $token;
        $token = &nextWord;
    } elseif($token =~ /\w+(\S+)/){
        $identifier = $1;
        $token = $2;
    }
    if($token eq "."){
        $const .= " : ".parseType(&nextWord);
        $token = $currentWords[$pointer];
    } elseif($token =~ /\w+\S*/){
        $const .= " : ".parseType($1);
        $token = $currentWords[$pointer];
    }
    if($token =~ /\S+\/){
        $const .= " = ".parseExpression($1);
    } elseif($token =~ /=$/){
        $const .= " = ".parseExpression(&nextWord);
    } else{
        $token = &nextWord;
        if($token =~ /\S+\/){
            $const .= " = ".parseExpression($1);
        } elseif($token =~ /=$/){
            $const .= " = ".parseExpression(&nextWord);
        } else{
            die "parse const statement failed!\n";
        }
    }
}
my $classBody = hashAdd(\%classList, $currentClass);
my $constFields = hashAdd($classBody, "const");
$constFields->{$identifier} = $const;
}

```

The supporting module *parseExpression* is defined according to its syntax
expression ::= conjunction “or” conjunction
as following:

```

Sparkf-lime...sub parseExpression{
  my $token = $_[0];
  my $expression = "";
  $token = $1 if($token =~ /(.)+);
  if($token =~ /\WorW/){
    my @wordBuffer = split(/\s+/, $token);
    $token = shift(@wordBuffer);
  }else{
    $expression = parseConjunction($token);
    while($expression ne ""){
      $token = &nextWord;
      if($token eq "or"){
        $expression .= " or ".parseConjunction(&nextWord);
      }else{
        &lastWord;
        last;
      }
    }
  }
  $expression;
}

```

The syntax of conjunction is
 conjunction ::= relational “and” relational :

```

Sparkf-lime...sub parseConjunction{
  my $conjunction = parseRelational($_[0]);
  while($conjunction ne ""){
    my $token = &nextWord;
    if($token eq "and"){
      $conjunction .= " and ".parseRelational(&nextWord);
    }else{
      &lastWord;
      last;
    }
  }
  $conjunction;
}

```

The syntax of relational is
 relational ::= additive [(“<” | “>” | “≤” | “≥” | “=” | “!=”) additive]:

```

Sparkf-lime...sub parseRelational{
  my $relational = "";
  my $token = $_[0];
  if($token =~ /(.)+(<|>|!=|<=|>=| | )(.+)/){
    $relational = parseAdditive($1)." ".$2." ".parseAdditive($3);
  }elseif($token =~ /(.)+(<|>|!=|<=|>=| | )$/){
    $relational = parseAdditive($1)." ".$2." ".parseAdditive(&nextWord);
  }else{
    $relational = parseAdditive($token);
    $token = &nextWord;
    if($token =~ /^(<|>|!=|<=|>=| | )$/){
      $relational .= " ".$1." ".parseAdditive(&nextWord);
    }elseif($token =~ /^(<|>|<=|>=|!=| | )(.+)/){
      $relational .= " ".$1." ".parseAdditive($2);
    }else{
      &lastWord;
    }
  }
  $relational;
}

```

The syntax of additive is
 additive ::= multiplicative (“+” | “-”) multiplicative

```
Sparkf-lime...sub parseAdditive{
my $additive = parseMultiplicative($_[0]);
my $token = $currentWords[$pointer];
while($additive ne ""){
  if($token =~ /(\+|-)$/){
    $additive .= " $1 ".parseMultiplicative(&nextWord);
  }elseif($token =~ /(\+|-)(\S+)/){
    $additive .= " $1 ".parseMultiplicative($2);
  }else{
    $token = &nextWord;
    if($token =~ /(\+|-)$/){
      $additive .= " $1 ".parseMultiplicative(&nextWord);
    }elseif($token =~ /(\+|-)(\S+)/){
      $additive .= " $1 ".parseMultiplicative($2);
    }else{
      $token = $currentWords[$pointer];
      &lastWord;
      $token = $currentWords[$pointer];
      last;
    }
  }
}
$additive;
}
```

The syntax of multiplicative is
 multiplicative ::= unary (“*” | “/” | “div” | “mod”) unary

```
Sparkf-lime...sub parseMultiplicative{
my $multiplicative = parseUnary($_[0]);
my $token = $currentWords[$pointer];
while($multiplicative ne ""){
  if($token =~ /(\*|\/)(\S+)/){
    $token = $2;
    $multiplicative .= " $1 ".parseUnary($token);
  }elseif($token =~ /(\+|-)$/){
    $token = &nextWord;
    $multiplicative .= " $1 ".parseUnary($token);
  }else{
    $token = &nextWord;
    if($token =~ /(\*|\/)(\S+)/){
      $token = $2;
      $multiplicative .= " $1 ".parseUnary($token);
    }elseif($token =~ /(\+|-)$/){
      $token = &nextWord;
      $multiplicative .= " $1 ".parseUnary($token);
    }elseif($token eq "div"){
      $token = &nextWord;
      $multiplicative .= " div ".parseUnary($token);
    }elseif($token eq "mod"){
      $token = &nextWord;
      $multiplicative .= " mod ".parseUnary($token);
    }else{
      &lastWord;
      last;
    }
  }
}
$multiplicative;
}
```

The syntax of unary is

unary ::= (“-” | “not” | “+”)unary | primary

and in turn the syntax of primary is

primary ::= integer | “nil” | “true” | “false” | designator | “new” name[actuals]:

```
Sparkf-lime...sub parseUnary{
  my ($unary, $token) = ("", $_[0]);
  if($token =~ /^(-|\+)(\S+)/){
    $unary = "$1".parseUnary($2);
  }elseif($token =~ /^(-|\+)$/){
    $unary = "$1".parseUnary(&nextWord);
  }elseif($token eq "not"){
    $unary = "not ".parseUnary(&nextWord);
  }elseif($token =~ /not\s+(\.+)/){
    $unary = "not ".parseUnary($1);
  }elseif($token =~ /(\d+)/){
    $unary = "$1";
  }elseif(($token =~ /^(nil)\W*/ ||
    ($token =~ /^(true)\W*/ ||
    ($token =~ /^(false)\W*/)){
    $unary = $1;
  }elseif($token eq "new"){
    $unary = "new " .&nextWord;
    $token = &nextWord;
    if($token =~ /^\(/{
      $unary .= " ".parseActuals($token);
    }else{
      &lastWord;
    }
  }elseif($token =~ /new\s+(\w+)\s*(\(.+\))/){
    $unary = "new " . $1 . parseActuals($2);
  }elseif($token =~ /(\w+)/){
    $unary = parseDesignator($token) unless($keyWordsList{$1});
  }
  $unary;
}
```

The syntax of designator is

designator ::= identifier “.” identifier | actuals :

```
Sparkf-lime...sub parseDesignator{
  my ($designator, $token) = ("", $_[0]);
  if($token =~ /^(\w+)$/){
    $designator = $token;
    $token = &nextWord;
  }elseif(($token =~ /^(\w+)(\.\S+)/)||($token =~ /^(\w+)(\(\S+)/)){
    $designator = $1;
    $token = $2;
  }else{
    die "parse designator failed![$token]\n";
  }
  while(1){
    if($keyWordsList{$designator}){
      $designator = "";
      &lastWord;
      last;
    }
    if($token =~ /^\.\/){
      $token = "." .&nextWord if($token =~ /^\.$/);
      if($token =~ /^\.(\w+)$/){
        $designator .= "." . $1;
        $token = &nextWord;
      }elseif($token =~ /^\.(\w+)(\S+)/){

```

```

    $designator .= ".".$1;
    $token = $2;
  }else{
    die "parse designator failed![3]\n";
  }
}elsif($token =~ /\^\(\){
  $designator .= parseActuals($token);
  $token = $currentWords[$pointer];
  if(($token =~ /\)(\.\S*)/)||($token =~ /\)(\(\S*)/)){
    $token = $2;
  }elsif($token =~ /\)$/){
    $token = &nextWord;
  }else{last;}
}else{
  &lastWord;
  last;
}
}
$designator;
}

```

For variables, the syntax is variable ::= “var” idList “:” type:

```

Sparkf-lime...sub parseVariable{
  my @vars = parseIdList($_[1]);
  my ($type, $public) = ("", $_[0]);
  my $token = $currentWords[$pointer];
  if($token =~ /\w:(\w+\S*)/){
    $type = parseType($1);
  }elsif($token =~ /\w+$/){
    $type = parseType(&nextWord);
  }else{
    $token = &nextWord;
    if($token =~ /\^:(\w+\S*)/){
      $type = parseType($1);
    }elsif($token eq ":{"){
      $type = parseType(&nextWord);
    }else{
      die "parse variable failed!\n";
    }
  }
}
foreach my $v(@vars){
  if($v =~ /\(\w+)\.\(\w+)/){
    $currentClass = $1;
    $v = $2;
  }
  my $classBody = hashAdd(\%classList, $currentClass);
  my $varFields = hashAdd($classBody, "var");
  $varFields->{$v} = $type;
}
}

```

For methods, the syntax method ::= “method” identifier formals[“:” type] statement:

```

Sparkf-lime...sub parseMethod{
  my @vars = ();
  my ($token, $public) = @_;
  my ($identifier, $type) = ("", "");
  my ($formals, $statement) = ("", "");
  if($token eq "method"){
    $token = &nextWord;
    if($token =~ /\^(\w+\.\w+)/){
      $identifier = $1;
      $token = &nextWord;
    }
  }
}

```

```

} elsif($token =~ /\^(\\w+\\.\\w+)(\\S+)/){
    $identifier = $1;
    $token = $2;
} elsif($token =~ /\^(\\w+)$/){
    $identifier = $1;
    $token = &nextWord;;
} elsif($token =~ /\^(\\w+)(\\S+)/){
    $identifier = $1;
    $token = $2;
} else{
    die "parse method failed!\n";
}
}
$formals = parseFormal($token);
$token = $currentWords[$pointer];
if(($token =~ /\):$/)||
((($formals eq "")&&($token =~ /:$/))){
    $token = &nextWord;
    $type = parseType($token);
} elsif(($token =~ /\):(\\w+\\.*)/)||
((($token =~ /:(\\w+\\.*)/)&&($formals eq ""))){
    $type = parseType($1);
} else{
    $token = &nextWord;
    if($token =~ /:$/){
        $token = &nextWord;
        $type = parseType($token);
    } elsif($token =~ /:(\\w+\\.*)/){
        $type = parseType($1);
    } else{
        &lastWord;
    }
}
}
my $newWord = 0;
$token = $currentWords[$pointer];
if(($type ne "")||(($type eq "")&&($formals eq ""))){
    if($token =~ /(\'.*)/){
        $token = $1;
    } else{
        $newWord = 1;
        $token = &nextWord;
    }
} elsif($formals ne ""){
    if($token =~ /\)(.+)/){
        $token = $1;
    } else{
        $newWord = 1;
        $token = &nextWord;
    }
}
}
$statement = parseStatementList($token);
&lastWord if(($statement eq "") && $newWord);
if($identifier =~ /(\\w+\\.\\w+)/){
    $currentClass = $1;
    $identifier = $2;
}
}
$identifier .= ".".$formals;
my $classBody = hashAdd(\%classList, $currentClass);
my $varFields = hashAdd($classBody, "method");
$varFields->{$identifier}={TYPE=>$type,STATEMENT=>$statement}
unless($varFields->{$identifier});
$varFields->{$identifier}->{STATEMENT} = $statement
if(($varFields->{$identifier}->{STATEMENT} eq "")&&($statement ne ""));
$varFields->{$identifier}->{TYPE} = $type
if(($varFields->{$identifier}->{TYPE} eq "")&&($type ne ""));
}
}

```

where the supporting module *parseFormal* is defined as following according to the syntax (formals ::= [“(” idList “:” type “,” idList “:” type “)”]):

```
Sparkf-lime...sub parseFormal{
  my ($formals, $token) = ("", $-[0]);
  if($token =~ /\(/){
    $formals = "(";
    if($token =~ /\($/){
      $token = &nextWord;
    }elseif($token =~ /\((\S+)/){
      $token = $1;
    }
  }
  while(1){
    if($token =~ /(\.\*\)\S*/){
      $formals .= $1;
      last;
    }else{
      $formals .= $token;
      $token = &nextWord;
    }
  }
  $formals;
}
```

where the supporting module *parseType* is defined as following according to the syntax (type ::= [“shared”] “array”[expression“,” expression]“of”(name | “integer” | “boolean”)):

```
Sparkf-lime...sub parseType{
  my $type = "";
  my $token = $-[0];
  if($token eq "shared"){
    $type = "shared ";
    $token = &nextWord;
  }
  while(1){
    if($token eq "array"){
      $type .= "array ";
      while($token ne "of"){
        $token = &nextWord;
        $type .= "$token ";
      }
      $type .= "of ";
      $token = &nextWord;
    }elseif($token eq "set"){
      $type .= "set ";
      $type .= &nextWord." ";
      $token = &nextWord;
    }else{
      last;
    }
  }
  if(($token =~ /^(integer)\W*/)||($token =~ /^(boolean)\W*/)){
    $type .= $1;
  }elseif($token =~ /\w+){
    $type .= parseName($token);
  }else{
    die "parse type failed![2]\n";
  }
  $type;
}
```

the supporting module *parseStatementList* is defined as following according to the syntax (statementList ::= statement “,” statement):

```
Sparkf-lime...sub parseStatementList{
my $statementList = parseStatement($-[0]);
$statementList = " if($statementList =~ /\s+$/);
while($statementList ne ""){
my $token = $currentWords[$pointer];
my @array = split(/;/, $token);
for(my $i = 1; $i <= $#array; $i++){
$statementList .=";\n".parseStatement($array[$i]);
}
if($token =~ /;$/){
$statementList .=";\n".parseStatement(&nextWord);
}else{
$token = &nextWord;
if($token =~ /\s+$/){
$statementList .=";\n".parseStatement(&nextWord);
}elseif($token =~ /\s+(\S+)/){
$statementList .=";\n".parseStatement($1);
}else{
&lastWord;
$statementList .="\n";
last;
}
}
}
$statementList;
}
```

Statement is the most complex one in Lime. Its syntax is

```
statement ::= designator [“:=” expression] |
designatorList “:=” expressionList |
“begin” statement “,” statement “end” |
“when” expression “do” statement |
“if” expression “then” statement [“else” statement] |
“while” expression “do” statement |
“repeat” statement “,” statement “until” expression |
variable statement |
constant statement |
“return” [expression] |
label
```

```
Sparkf-lime...sub parseStatement{
my $space = " ";
$layers++;
for(my $i = 0; $i < $layers; $i++){
$space .= " ";
}
my ($statement, $token) = ($space, $-[0]);
if($token eq "begin"){
```



```

    $statement .= " : ".parseType($1);
  }else{
    &lastWord;
  }
  $token = $currentWords[$pointer];
}
if($token =~ /=$/){
  $statement .= " = ".parseExpression(&nextWord);
}elseif($token =~ /=(\S+)/){
  $statement .= " = ".parseExpression($1);
}else{
  $token = &nextWord;
  if($token =~ /=$/){
    $statement .= " = ".parseExpression(&nextWord);
  }elseif($token =~ /=(\S+)/){
    $statement .= " = ".parseExpression($1);
  }else{
    die "parse const statement failed!\n";
  }
}
}elseif($token eq "return"){
  my $t = parseExpression(&nextWord);
  &lastWord if($t eq "");
  $statement .= "return ".$t;
}elseif($token =~ /\^\w+\/){
  $statement .= $token;
}elseif($token =~ /\w+\/){
  $statement .= parseAssignment($token) unless($keyWordsList{$token});
}
$layers--;
$statement;
}

```

where local variable `%keyWordsList` is used to identify the key words of Lime and defined as:

```

Sparkf-lime...my %keyWordsList = ("abort", 1, "action", 1, "and", 1,
    "array", 1, "char", 1, "do", 1,
    "begin", 1, "boolean", 1, "case", 1,
    "class", 1, "const", 1, "div", 1,
    "downto", 1, "else", 1, "end", 1,
    "export", 1, "integer", 1, "map", 1,
    "false", 1, "final", 1, "if", 1,
    "import", 1, "initialization", 1,
    "method", 1, "mod", 1, "new", 1,
    "nil", 1, "package", 1, "real", 1,
    "not", 1, "of", 1, "or", 1,
    "private", 1, "program", 1, "procedure", 1,
    "repeat", 1, "return", 1, "set", 1,
    "this", 1, "sequence", 1, "until", 1,
    "skip", 1, "super", 1, "then", 1,
    "to", 1, "true", 1, "type", 1,
    "var", 1, "when", 1, "while", 1);

```

and local variable `$layer` is used to count the nest of statement and defined as:

```

Sparkf-lime...my $layers = -1;

```

For actions, the syntax is `action ::= "action" identifier statement`:

```

Sparkf-lime...sub parseAction{
  my $identifier = $_[0];
  die "parse action failed!\n" if($identifier !~ /\w+$/);
  my $statement = parseStatementList(&nextWord);
}

```

```

my $classBody = hashAdd(\%classList, $currentClass);
my $actionFields = hashAdd($classBody, "action");
$actionFields->{$identifier} = $statement
unless($actionFields->{$identifier});
}

```

For initializations, the syntax is `initialization ::= "initialization" formals statement`:

```

Sparkf-lime...sub parseInitialization{
my $token = $_[0];
my $statement = "";
my $formals = parseFormal($token);
if($formals ne ""){
$token = $currentWords[$pointer];
if($token =~ /\)(\S+)/){
$statement = parseStatementList($1);
}else{
$statement = parseStatementList(&nextWord);
}
}else{
$formals = "init";
$statement = parseStatementList($token);
}
my $classBody = hashAdd(\%classList, $currentClass);
my $actionFields = hashAdd($classBody, "initialization");
$actionFields->{$formals} = $statement
unless($actionFields->{$formals});
}

```

For labels, the syntax is `label ::= 'char'`:

```

Sparkf-lime...sub parseLabel{
my @array = ($_[0]);
my $classBody = hashAdd(\%classList, $currentClass);
if($classBody->{"label"}){
my $old = $classBody->{"label"};
my $find = 0;
foreach my $v(@$old){
if($v eq $_[0]){
$find = 1;
last;
}
}
push(@$old, $_[0]) unless($find);
}else{
$classBody->{"label"} = ($_[0]);
}
}

```

Lime also support multiple assignment as following:

```

Sparkf-lime...sub parseAssignment{
my (@designatorList, @expressionList) = ((), ());
my ($finish, $token, $assignment) = (0, $_[0], "");
my @statementList = split(/;/, $token);
$token = $statementList[0];
$finish = 1 if((($token =~ /\)/) || (($token =~ /\)/) else \'/));
until($finish){
last if($#statementList > 0);
my $temp = &nextWord;
if($temp =~ /\w+){
if(($1 eq "end") || ($1 eq "else") ||
($1 eq "until") || ($memberList{$1})){

```

```

    &lastWord;
    $finish = 1;
  }else{
    if($token =~ /\w+$/){
      $token.=" ". $temp;
    }else{
      $token.= $temp;
    }
  }
}else{
  if($token !~ /\w+$/){
    $token.=" ". $temp;
  }else{
    $token.= $temp;
  }
}
}
@statementList = split(/;/, $token);
}
$token = $statementList[0];
if($token =~ /(.)=(.+)/){
  @designatorList = split(/,/ , $1);
  @expressionList = split(/,/ , $2);
  for(my $i=0; $i< $#designatorList; $i++){
    $assignment .= parseDesignator($designatorList[$i])." , ";
  }
  $assignment.=parseDesignator($designatorList[$#designatorList])." := ";
  for(my $i=0; $i< $#expressionList; $i++){
    $assignment .= parseExpression($expressionList[$i])." , ";
  }
  $assignment .= parseExpression($expressionList[$#expressionList]);
}else{
  $assignment .= parseDesignator($token);
}
}
$assignment;
}
}

```

For actuals, the syntax is actuals ::= “(” expression “,” expression “)”:

```

Sparkf-lime ... sub parseActuals{
  my $token = $_[0];
  my $actuals = "(";
  if($token =~ /\((.+)/){
    $token = $1;
  }else{
    die "parse actuals statement failed![1]\n";
  }
  while(1){
    if($token =~ /(.)+(\S+)){
      $actuals .= parseExpression($1)." , ";
      $token = $2;
    }elseif($token =~ /(.)$/){
      $actuals .= parseExpression($1)." , ";
      $token = &nextWord;
    }elseif($token =~ /(.)\)/){
      $actuals .= parseExpression($1);
      last;
    }else{
      die "parse actuals statement failed![2]\n";
    }
  }
  $actuals .= ")";
  $actuals;
}

```

The syntax of name is name ::= identifier “.” identifier :

```

Sparkf-lime...sub parseName{
my ($name, $token) = ("", $_[0]);
if($token =~ /\w+$/){
$name = $token;
$token = &nextWord;
}elsif($token =~ /^\w+(\S+)/){
$name = $1;
$token = $2;
}else{
die "parse name failed!\n";
}
while(1){
if($token =~ /\./){
if($token =~ /\.(\\w+$/){
$token = &nextWord;
$name .= ".".$1;
}elsif($token =~ /\.(\\w+(\S+))/){
$token = $2;
$name .= ".".$1;
}else{
die "parse name failed!\n";
}
}
}
}
$name;
}

```

The syntax of idList is idList ::= identifier “,” identifier:

```

Sparkf-lime...sub parseIdList{
my @idList = ();
my $token = $_[0];
my $finish = 0;
if($token =~ /\w+$/){
push(@idList, $1);
$token = &nextWord;
}elsif($token =~ /\w+(\, \S*)/){
push(@idList, $1);
$token = $2;
}elsif($token =~ /\w+//){
push(@idList, $1);
$finish = 1;
}else{
die "parse idList failed![1]\n";
}
until($finish){
if($token =~ /\,(\\w+$/){
$token = &nextWord;
push(@idList, $1);
}elsif($token =~ /\,$/){
$token = ",".&nextWord;
}elsif($token =~ /\,(\\w+(\, \S*)/){
$token = $2;
push(@idList, $1);
}elsif($token =~ /\,(\\w+(\, +))/){
$finish = 1;
push(@idList, $1);
}else{
&lastWord;
$finish = 1;
}
}
}
@idList;

```

```
}

```

For base, the syntax is `base ::= [“inherit” name | “extend” name]:`

```
Sparkf-lime...sub parseBase{
  my $classBody = hashAdd(\%classList, $currentClass);
  $classBody->{$_[0]} = &nextWord;
}
```

For implement, the syntax is `implement ::= “implement” name:`

```
Sparkf-lime...sub parseImplement{
  my @array = ();
  my $token = $_[0];
  while(1){
    while($token =~ /(\w+)(\S+)/){
      push(@array, $1);
      $token = $2;
    }
    if($token =~ /(\w+)/){
      push(@array, $1);
      $token = &nextWord;
    } else{
      push(@array, $token);
      $token = &nextWord;
      if($token =~ /\^(\S+)/){
        $token = $1;
      } elsif($token =~ /\^./){
        $token = &nextWord;
      } else{
        &lastWord;
        last;
      }
    }
  }
  die "parse implement failed!\n[2]" if($#array < 0);
  my $classBody = hashAdd(\%classList, $currentClass);
  if($classBody->{"implement"}){
    my $old = $classBody->{"implement"};
    my $find = 0;
    foreach my $v(@array){
      foreach my $w(@$old){
        if($v eq $w){
          $find = 1;
          last;
        }
      }
      $find ? $find = 0 : push(@$old, $v);
    }
  } else{
    $classBody->{"implement"} = \@array;
  }
}
```

All the parsing result of code blocks will be inserted into the class list:

```
Sparkf-lime...my %classList = ();
```

The second process in main is outputting the parsing result.

```
Sparkf-lime...sub doOutput{
  generateCode("");
  &insertCode;
  &createModFile;
}
```

```
}

```

Module *generateCode* is responsible for generating the actual program code files for compiler using the information in the variable `%classList` and defined as:

```
Sparkf-lime...sub generateCode{
my $outFile;
my @keyList = keys(%classList);
foreach my $v(@keyList){
my $classBody = $classList{$v};
if($_[0] eq ""){
open $outFile, "> $v.lime" || die "Create file failed!\n";
}else{
$outFile = $_[0];
}
print $outFile "final " if($classBody->{"final"});
print $outFile "class $v ";
if($classBody->{extend}){
print $outFile "extend $classBody->{extend} ";
}elseif($classBody->{inherit}){
print $outFile "inherit $classBody->{inherit} ";
}
if($classBody->{implement}){
print $outFile "implement ";
my $first = 1;
my $memberBody = $classBody->{implement};
foreach my $w(@$memberBody){
if($first){
print $outFile "$w";
$first = 0;
}else{
print $outFile ", $w";
}
}
}
print $outFile "\n";
my @memberList = keys(%$classBody);
foreach my $u(@memberList){
my $memberBody = $classBody->{$u};
unless(($u eq "implement")||($u eq "inherit")||
($u eq "extend")||($u eq "final")){
my @fieldList = keys(%$memberBody);
foreach my $f(@fieldList){
if($u eq "var"){
print $outFile " var $f : $memberBody->{$f}\n\n";
}elseif($u eq "const"){
if($memberBody->{$f} =~ /\(d+\)\.(\.+)/){
print $outFile " ";
print $outFile "public " if($1);
print $outFile "const $$f2\n";
}
}elseif($u eq "initialization"){
print $outFile " initialization ";
print $outFile "$f" if($f ne "init");
print $outFile "\n";
print $outFile "$memberBody->{$f}\n";
if($memberBody->{$f} ne " ");
}elseif($u eq "action"){
print $outFile " action $f\n";
print $outFile "$memberBody->{$f}\n";
if($memberBody->{$f} ne " ");
}elseif($u eq "method"){
if($f =~ /\(w+\)\.(\.+)/){
print $outFile " method $1";
print $outFile "$1" if($f =~ /\w+\.(\.+)/);
print $outFile " : $memberBody->{$f}->{TYPE}"

```

```

        if($memberBody->{$f}->{TYPE} ne "");
        print $outFile "\n";
        print $outFile "$memberBody->{$f}->{STATEMENT}"
        if($memberBody->{$f}->{STATEMENT} ne "");
        print $outFile "\n";
    }
}
}#end of foreach
}#end of unless
}
print $outFile "end\n\n";
close $outFile if($_[0] eq "");
}#end of foreach
}

```

After generating code files, front end will insert a copy of code to the origin file, if it find the special tag-pair, “//CODE LIST BEGIN” and “//CODE LIST END”, by module *insertCode*, which is defined as:

```

Sparkf-lime...sub insertCode{
my $switch = 0;
open my $oldFile, "< $old" || die "Can't open $old: $!";
open my $newFile, "> $new" || die "Can't open $new: $!";
while(<$oldFile>){
if($_ =~ /\^\/\CODE LIST BEGIN/){
print $newFile $_;
print $newFile "----\n";
$switch = 1;
generateCode($newFile);
}elseif($_ =~ /\^\/\CODE LIST END/){
print $newFile "----\n";
print $newFile $_;
$switch = 0;
}else{
print $newFile $_ unless($switch);
}
}
close $oldFile || die "Can't close $old: $!";
close $newFile || die "Can't close $old: $!";
unlink($old);
rename($new, $old) || die "Can't rename $old: $!";
}

```

In fact, this module creates a file named by the variable `$new`, which is defined as:

```

Sparkf-lime...my $new = "temp";

```

and then copy the content other than the part between the code list tags into this new file, and after that insert the content of `%classList` into this new file, and finally deletes the old file and renames the new file with the old name.

Module *createModFile* is responsible for creating a scripts file according to the content of variable `@diagramList`. If it is empty, nothing will happen; otherwise, front end will create a file with the name defined by variable `$filename`, which is defined as:

```
Sparkf-lime...my $filename = "oolp.mod";
```

to describe the diagrams used in this program.

```
Sparkf-lime...sub createModFile{
  open FILE, "> $filename" || die "Open file failed!";
  foreach my $v(@diagramList){
    print FILE "\@CLASSDIAGRAM $v\n";
  }
  print FILE "\n";
  my @keyList = keys(%classList);
  foreach my $v(@keyList){
    my $classBody = $classList{$v};
    print FILE "\@CLASS $v\n";
    print FILE "\@FINAL\n" if($classBody->{"final"});
    if($classBody->{extend}){
      print FILE "\@EXTEND $classBody->{extend}\n";
    }elseif($classBody->{inherit}){
      print FILE "\@INHERIT $classBody->{inherit}\n";
    }
    if($classBody->{implement}){
      print FILE "\@IMPLEMENT ";
      my $sep = " ";
      my $memberBody = $classBody->{implement};
      foreach my $w(@$memberBody){
        print FILE "$sep$w";
        $sep = ", " if($sep eq " ");
      }
      print FILE "\n";
    }
    my @memberList = keys(%$classBody);
    foreach my $u(@memberList){
      my $memberBody = $classBody->{$u};
      unless(($u eq "implement")||($u eq "inherit")||
        ($u eq "extend")||($u eq "final")){
        my @fieldList = keys(%$memberBody);
        foreach my $f(@fieldList){
          if($u eq "var"){
            print FILE "\@VAR $f $memberBody->{$f}\n";
          }elseif($u eq "const"){
            if($memberBody->{$f} =~ /\(d+\)\.(.+)\){
              print FILE "\@CONST ";
              print FILE "\@PUBLIC " if($1);
              print FILE "$f$2\n";
            }
          }elseif($u eq "initialization"){
            print FILE "\@INIT ";
            print FILE "$f" if($f ne "init");
            print FILE "\n";
          }elseif($u eq "action"){
            print FILE "\@ACTION $f\n";
          }elseif($u eq "method"){
            if($f =~ /\(w+\)\.\/){
              print FILE "\@METHOD $1 ";
              print FILE "$1 " if($f =~ /\w+\.(.+)\);
              print FILE "$memberBody->{$f}->{TYPE}"
                if($memberBody->{$f}->{TYPE} ne " ");
              print FILE "\n";
            }
          }
        }
      }
    }
  }#end of foreach
}#end of unless
}#end of foreach
print FILE "\n";
}
close FILE || die "Close $filename failed!";
```

```
}

```

The third process in main is to call the back end of Spark. It is simply defined as:

```
Sparkf-lime...sub callBackEnd{
  system "perl sparke.pl";
}
```

5.4 Back End

The back end of Spark takes the graphic notation describer mentioned above as input and produce all the graphic notation files as output. However, layout algorithm is out of the range of this thesis, so we choose an automatic diagram layout tool, Graphviz, to fulfill this task. Graphviz is a package of open source tools initiated by AT&T Research Labs for drawing graphs specified in dot language scripts. Now, let us look at how the back end works.

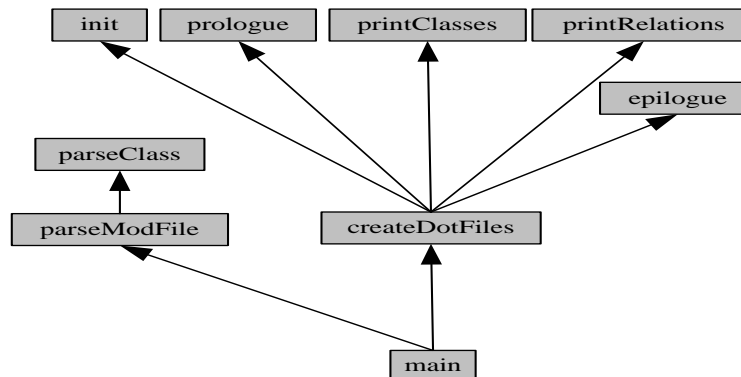


Figure 5.4: Top-level structure for Spark back end

The entrance of Spark back end for lime,

```
Sparkb...&main;
```

is defined as following:

```
Sparkb...sub main{
  &parseModFile;
  &createDotFiles;
}
```

The first process in main is to parse the script file created by front end and defined as:

```
Sparkb...sub parseModFile{
open FILE, "< $filename" || die "Open file failed!\n";
while(<FILE>){
  if($_- = ~ /\@CLASSDIAGRAM/){
    push(@diagramList, $_);
  } elsif($_- = ~ /\@CLASS/){
    parseClass($_);
  }
}
close FILE || die "Close file failed!\n";
unlink($filename);
}
```

where local variable \$filename hold the name of script file and defined as:

```
Sparkb...my $filename = "oolp.mod";
```

and variable @diagramList keeps all the information about diagrams to be created and defined as:

```
Sparkb...my @diagramList = ();
```

Module *parseClass* is used to retrieve all the class information by reconstructing the class list %classList, which is defined as:

```
Sparkb...my %classList = ();
```

and *parseClass* is defined as:

```
Sparkb...sub parseClass{
chop($_[0]);
my @words = split(/\s+/, $_[0]);
my $name = $words[1];
my %newHash = ();
while(<FILE>){
  last if($_ eq "\n");
  chop;
  @words = split;
  if($words[0] eq "\@VAR"){
    unless($newHash{var}){
      my %hash = ();
      $newHash{var} = \%hash;
    }
    my $type = $words[2];
    for(my $i=3; $i<=#words; $i++){
      $type .= " $words[$i]";
    }
    $newHash{var}->{$words[1]} = $type;
  } elsif($words[0] eq "\@ACTION"){
    unless($newHash{action}){
      my @array = ();
      $newHash{action} = \@array;
    }
    my $body = $newHash{action};
    push(@$body, $words[1]);
  } elsif($words[0] eq "\@METHOD"){
    unless($newHash{method}){
      my %hash = ();
    }
  }
}
```

```

    $newHash{method} = \%hash;
  }
  if($#words == 3){
    $newHash{method}->{$words[1].".". $words[2]} = $words[3];
  }elseif($#words == 1){
    $newHash{method}->{$words[1]} = "";
  }elseif($#words == 2){
    if($words[2] =~ /\(.+\)/){
      $newHash{method}->{$words[1].".". $words[2]} = "";
    }else{
      $newHash{method}->{$words[1]} = $words[2];
    }
  }
} elseif($words[0] eq "\@INIT"){
  unless($newHash{init}){
    my @array = ();
    $newHash{init} = \@array;
  }
  my $body = $newHash{init};
  if($#words == 1){
    push(@$body, $name.".". $words[1]);
  }elseif($#words == 0){
    push(@$body, $name);
  }
} elseif($words[0] eq "\@IMPLEMENT"){
  unless($newHash{implement}){
    my @array = ();
    $newHash{implement} = \@array;
  }
  my @temp = split(/./, $words[1]);
  my $body = $newHash{implement};
  while($#temp >= 0){
    push(@$body, shift(@temp));
  }
} elseif($words[0] eq "\@EXTEND"){
  $newHash{extend} = $words[1];
} elseif($words[0] eq "\@INHERIT"){
  unless($newHash{inherit}){
    my @array = ();
    $newHash{inherit} = \@array;
  }
  my @temp = split(/./, $words[1]);
  my $body = $newHash{inherit};
  while($#temp >= 0){
    push(@$body, shift(@temp));
  }
} elseif($words[0] eq "\@CONST"){
}
}
$classList{$name} = \%newHash;
}

```

The second process in main is to *createDotFiles*, which is defined as:

```

Sparkb...sub createDotFiles{
  foreach my $v(@diagramList){
    my @diagram = split(/\s+/, $v);
    if($diagram[1] =~ /(\w+)\.(\w+)/){
      &init;
      for(my $i=2; $i<=#diagram; $i++){
        if($diagram[$i] =~ /\@VERTICAL/){
          $isHorizontal = 0;
        }elseif($diagram[$i] =~ /\@BRIEF/){
          $isBrief = 1;
        }elseif($diagram[$i] =~ /\@CONCISE/){
          $isConcise = 1;
        }elseif($diagram[$i] =~ /\@HEAD/){
          $isHead = 1;
        }
      }
    }
  }
}

```

```

    }elseif($diagram[$i] =~ /\@METHOD/){
        $isMethodOnly = 1;
    }elseif($diagram[$i] =~ /\@ATTRIBUTE/){
        $isAttributeOnly = 1;
    }elseif($diagram[$i] =~ /(\w+)/){
        push(@classesAdded, $diagram[$i]) unless(&findClass($diagram[$i]));
    }
}
next if($#classesAdded < 0);
open my $out, "> $1.dot" || die "Create file failed!\n";
prologue($out);
printClasses($out);
printRelations($out);
epilogue($out);
system "dot -T$2 -o$1.$2 $1.dot";
unlink("$1.dot");
}
}
}

```

For each item in @diagramList, it will create one dot file for GraphViz and one dot file is corresponding to one picture. The supporting function *findClass* is used to determine if the class will appear in this diagram and defined as:

```

Sparkb...sub findClass{
    my $found = 0;
    foreach my $v(@classesAdded){
        if($v eq $-[0]){
            $found = 1;
            last;
        }
    }
    $found;
}

```

Every diagram is new, so the back end clears the environment first every time. The supporting function *init* is defined as:

```

Sparkb...sub init{
    (@classesAdded, @relationsAdded) = ((), ());
    ($isHorizontal, $isBrief, $isConcise, $isHead) = (1, 0, 0, 0);
    ($isMethodOnly, $isAttributeOnly, $isActionOnly) = (0, 0, 0);
}

```

where the variable @classesAdded is used to hold all the classes, which are going to appear in the current diagram and defined as:

```

Sparkb...my @classesAdded = ();

```

and the variable @relationsAdded is used to describe all the relations between these classes and defined as:

```

Sparkb...my @relationsAdded = ();

```

and the variable \$isHorizontal is used to describe the direction of this diagram and its default value is 1, which means that the diagram will be drawn horizontally.


```

    push(@relationsAdded, "$v.extend.$classBody->{extend}\n")
    if(&findClass($classBody->{extend}));
  }
  if($classBody->{var}){
    my $memberBody = $classBody->{var};
    my @varList = keys(%$memberBody);
    foreach my $w(@varList){
      my $stpe = $memberBody->{$w};
      if($stpe =~ /^(.+)\s+\@LIST\s+(\w+)/){
        push(@relationsAdded, "$v.haslist.$2.$w\n") if(&findClass($2));
      }else{
        push(@relationsAdded, "$v.has.$stpe.$w\n") if(&findClass($stpe));
      }
    }
  }
}
unless($isHead){
  #show Attributes
  my $displayed = 0;
  innerTableStart($out);
  unless($isMethodOnly || $isActionOnly){
    if($classBody->{var}){
      my $memberBody = $classBody->{var};
      my @varList = keys(%$memberBody);
      foreach my $w(@varList){
        $displayed |= attribute($out, $w, $memberBody->{$w});
      }
    }
  }
  tableLine($out, "LEFT", " ") unless($displayed);
  innerTableEnd($out);
  #show operation
  innerTableStart($out);
  $displayed = 0;
  unless($isAttributeOnly || $isActionOnly){
    if($classBody->{init}){
      $displayed = 1;
      my $memberBody = $classBody->{init};
      foreach my $w(@$memberBody){
        if($w =~ /^(.+)\.(.+){
          operation($out, $1, $2, "");
        }else{
          operation($out, $w, "()", "");
        }
      }
    }
  }
  if($classBody->{method}){
    $displayed = 1;
    my $memberBody = $classBody->{method};
    my @methodList = keys(%$memberBody);
    foreach my $w(@methodList){
      my $stpe = $memberBody->{$w};
      my ($para, $name) = ("", "");
      if($w =~ /^(.+)\.(.+){
        ($name, $para) = ($1, $2);
      }else{
        $name = $w;
      }
      operation($out, $name, $para, $stpe);
    }
  }
  tableLine($out, "LEFT", " ") unless($displayed);
  innerTableEnd($out);
  #show action
  innerTableStart($out);
  $displayed = 0;
  unless($isAttributeOnly || $isMethodOnly){
    if($classBody->{action}){

```

```

        $displayed = 1;
        my $memberBody = $classBody->{action};
        foreach my $w(@$memberBody){
            operation($out, $w, "()", "");
        }
    }
}
tableLine($out, "LEFT", " ") unless($displayed);
innerTableEnd($out);
}#end of unless($isHead)
externalTableEnd($out);
nodeProperties($out);
}
}

```

where the variable `$isInterface` is defined as:

```
Sparkb...my $isInterface = 0;
```

The supporting function `externalTableStart` is used to draw the start border of class diagram and defined as:

```

Sparkb...sub externalTableStart {
    my ($out, $bgcolor) = ($_[0], "");
    $bgcolor = " bgcolor=\"".$nodeFillColor.\" \" if($nodeFillColor ne "");
    print $out "<<table border=\"0\" cellborder=\"1\" cellspacing=\"0\" \"
        \" cellpadding=\"2\" port=\"p\".\"$bgcolor.\">\".$linePostfix;
}

```

where the variable `$nodeFillColor` is defined as:

```
Sparkb...my $nodeFillColor = "";
```

and function `innerTableStart` is used to draw the inner border of class diagram and defined as:

```

Sparkb...sub innerTableStart {
    my $out = $_[0];
    print $out $linePrefix.$linePrefix."<tr><td><table border=\"0\" \"
        \" cellspacing=\"0\" cellpadding=\"1\">\".$linePostfix;
}

```

where the variable `$linePrefix` and `$linePostfix` are defined as:

```
Sparkb...my ($linePrefix, $linePostfix) = ("\t", "\n");
```

and function `tableLine` is used to draw a common line of class diagram and defined as:

```

Sparkb...sub tableLine {
    my $out = $_[0];
    my ($stopen, $stclose) = ("", "</td></tr>");
    my $prefix = $linePrefix.$linePrefix.$linePrefix;
    if($_[1] eq "CENTER"){
        $stopen = $prefix."<tr><td> ";
    } elsif($_[1] eq "LEFT"){
        $stopen = $prefix."<tr><td align=\"left\"> ";
    } elsif($_[1] eq "RIGHT"){

```

```

    $stopen = $prefix.<tr><td align=\"right\"> ";
  }
  print $out $stopen.$_[2].$tclose.$linePostfix;
}

```

where the variable \$align is defined as:

```
Sparkb...my $align = "CENTER";
```

and function *innerTableEnd* is used to draw inner end of class diagram and defined as:

```
Sparkb...sub innerTableEnd{
  my $out = $_[0];
  print $out $linePrefix.$linePrefix.</table></td></tr>". $linePostfix;
}

```

and function *externalTableEnd* is used to close class diagram and defined as:

```
Sparkb...sub externalTableEnd{
  my $out = $_[0];
  print $out $linePrefix.$linePrefix.</table>>>";
}

```

and function *attribute* is used to display the attributes of class in diagram and defined as:

```
Sparkb...sub attribute{
  my $display = 1;
  if(&findClass($_[2])){
    $display = 0;
  }else{
    if($_[2] =~ /^(.+)\s+\@LIST\s+(\w+)/){
      if(&findClass($2)){
        $display = 0;
      }else{
        my $att = $_[1];
        $att .= " : ".$1 unless($isConcise || $isBrief);
        tableLine($_[0], "LEFT", $att);
      }
    }else{
      my $att = $_[1];
      $att .= " : ".$_[2] unless($isConcise || $isBrief);
      tableLine($_[0], "LEFT", $att);
    }
  }
  $display;
}

```

and function *operation* is used to print operations of class in diagram and defined as:

```
Sparkb...sub operation{
  my ($out, $name, $para, $type) = @_;
  my $cs = $name;
  if($isBrief){
    $cs .= "()";
  }else{
    $cs .= $para;
    $type = "" if($isConcise);
    $cs .= " : ".$type if($type ne "");
  }
}

```

```

}
tableLine($out, "LEFT", $cs);
}

```

and function *nodeProperties* is used to print the common properties of the nodes of diagram and defined as:

```

Sparkb...sub nodeProperties{
my $out = $_[0];
print $out " , fontname=\"".$nodeFontName."\"".
" , fontcolor=\"".$nodeFontColor."\"".
" , fontsize=\"".$nodeFontSize."\";\n";
}

```

where the variable *\$nodeFontName*, *\$nodeFontColor*, and *\$nodeFontSize* are defined as:

```

Sparkb...my ($nodeFontName, $nodeFontColor, $nodeFontSize) = ("arial", "black", 10);

```

After that, back end begins to print the relations listed in *@relationsAdded* as following:

```

Sparkb...sub printRelations{
my $out = $_[0];
foreach my $r(@relationsAdded){
my ($tailLabel, $headLabel) = ("", "");
if($r =~ /(\w+)\.extend\.(\\w+)/){
print $out "\t//".$1." extend ".$2."\\n"."\\t".$2.:p -> ".$1.
":p [dir=back, arrowtail=empty, color=\"".$edgeColor."\"];\\n";
}elsif($r =~ /(\w+)\.implement\.(\\w+)/){
print $out "\t//".$1." implement ".$2."\\n"."\\t".$2.:p -> ".$1.
":p [dir=back, arrowtail=empty, style=dashed, color=\"".$edgeColor."\"];\\n";
}elsif($r =~ /(\w+)\.inherit\.(\\w+)/){
print $out "\t//".$1." inherit ".$2."\\n"."\\t".$2.:p -> ".$1.
":p [dir=back, arrowtail=empty, color=\"".$edgeColor."\"];\\n";
}elsif($r =~ /(\w+)\.haslist\.(\\w+)\.(\\w+)/){
print $out "\t// ".$1." has a list of ".$2."\\n"."\\t".$1.:p -> ".$2.:p ["
" taillabel=\"".$tailLabel."\\", ". label=\"".$3."\\", ". headlabel=\"".$headLabel.
"\\", ". fontname=\"".$edgeFontName."\\", ". fontcolor=\"".$edgeFontColor."\\", ".
" fontsize=\"".$edgeFontSize." , ". color=\"".$edgeColor."\\", ".$associationMap{"list"}."];\\n";
}
}
}
}

```

where the variable *\$edgeFontName*, *\$edgeColor*, *\$edgeFontSize*, and *\$edgeFontColor* are defined as:

```

Sparkb...my ($edgeFontName, $edgeColor) = ("arial", "red");
my ($edgeFontSize, $edgeFontColor) = (10, "black");

```

and the variable *%associationMap* is defined as:

```
Sparkb...my %associationMap = ( "assoc",      "arrowhead=none",
                               "navassoc",   "arrowhead=open",
                               "has",        "arrowhead=none, arrowtail=diamond",
                               "composed",   "arrowhead=none, arrowtail=diamond",
                               "list",       "arrowhead=dot, arrowtail=diamond",
                               "depend",     "arrowhead=open, style=dashed");
```

The last job is to print epilogue as following:

```
Sparkb...sub epilogue{
  my $out = $_[0];
  print $out "}\n";
  close $out;
}
```

5.5 Testing

To verify the design and implementation of Spark, we performed testing following the strategy mentioned in Section 3.6.4.

5.5.1 Usability Testing

Correct Usage

We tried to use Spark with an actual OOLP program file name, leaf.txt, as following:

```
perl sparkf-lime.pl leaf.txt
```

The result is that there are three Lime files and four pictures generated by Spark.

No Parameters

We tried to use Spark without parameters as following:

```
perl sparkf-lime.pl
```

The result is

```
Usage: perl sparkf.pl filename
```

Wrong Parameters

We tried to use Spark with fake file name as following:

```
perl sparkf-lime.pl aaa
```

The result is

```
Cannot open aaa!
```

These three results show that Spark can handle both legal and illegal usages and satisfies the design requirements.

5.5.2 Unit Testing

Syntax Coverage

We composed a sample code program file that coverage all the syntax of Eiffel.

```
Testing of a declaration of a class .
-----
class STUDENT
end

Testing of the inheritance relation of a class .
-----
inherit STUDENT...PEOPLE

Testing of two features of a class .
-----
feature {NONE} PEOPLE...name: STRING
      age: INTEGER

Testing of a deferred class .
-----
deferred class PEOPLE
end

Testing of a operation with formal specification and various statements of a class
-----
feature {NONE}
  STUDENT...set_name(new_name: STRING) is
    local
      a: INTEGER
    do
      name := new_name
      create employees.make
      getup
      if text_file_read.is_connected then
        split := input_string.split
      end
      if text_file_read.is_connected then
        split := input_string.split
      else
        split := input_string.split
      end
      inspect text_file_read.last_string.first.to_upper
```

```

        when 'S' then
            text_file_read.read_line
        when 'B', 'T' then
            text_file_read.read_line
        else
            end
        from text_file_read.read_line
        until text_file_read.end_of_input
        loop
            text_file_read.read_line
        end
    end
end

-----
Testing of a deferred operation with formal specification of a class
-----
feature
    PEOPLE...set_name(n: STRING) is
        require
            n /= Void
        deferred
        end
end

-----
Testing of a constant of a class
-----
feature {NONE} STUDENT...min_age: INTEGER is 5
-----
Testing of an invariant of a class
-----
invariant
    PEOPLE...age > 0
           age < 200
end

-----
Testing of an operation with result of a class.
-----
feature {ANY}
    STUDENT...match (n: STRING): BOOLEAN is
        do
            Result := name = n
        end
end

-----
//CODE LIST BEGIN
-----
//CODE LIST END

```

After running Spark on this sample,

[perl sparkf-eiffel.pl coverage.txt](#)

we got:

```

Testing of a declaration of a class.
-----
class STUDENT
end
-----
Testing of the inheritance relation of a class.
-----
inherit STUDENT...PEOPLE
-----
Testing of two features of a class.
-----

```

```

feature {NONE} PEOPLE...name: STRING
    age: INTEGER
-----
Testing of a deferred class.
-----
deferred class PEOPLE
end
-----
Testing of a operation with formal specification and various statements of a class
-----
feature {NONE}
    STUDENT...set_name(new_name: STRING) is
        local
            a: INTEGER
        do
            name := new_name
            create employees.make
            getup
            if text_file_read.is_connected then
                split := input_string.split
            end
            if text_file_read.is_connected then
                split := input_string.split
            else
                split := input_string.split
            end
            inspect text_file_read.last_string.first.to_upper
            when 'S' then
                text_file_read.read_line
            when 'B', 'T' then
                text_file_read.read_line
            else
            end
            from text_file_read.read_line
            until text_file_read.end_of_input
            loop
                text_file_read.read_line
            end
        end
    end
-----
Testing of a deferred operation with formal specification of a class
-----
feature
    PEOPLE...set_name(n: STRING) is
        require
            n /= Void
        deferred
        end
-----
Testing of a constant of a class
-----
feature {NONE} STUDENT...min_age: INTEGER is 5
-----
Testing of an invariant of a class
-----
invariant
    PEOPLE...age > 0
            age < 200
end
-----
Testing of an operation with result of a class.
-----
feature {ANY}
    STUDENT...match (n: STRING): BOOLEAN is
        do
            Result := name = n
        end
    end
-----

```

```
//CODE LIST BEGIN
-----
class STUDENT
inherit
  PEOPLE
feature {NONE}
  set_name(new_name: STRING) is
    local
      a : INTEGER
    do
      name := new_name
      create employees.make
      getup
      if text_file_read.is_connected then
        split := input_string.split
      end
      if text_file_read.is_connected then
        split := input_string.split
      else
        split := input_string.split
      end
      inspect
        text_file_read.last_string.first.to_upper
      when 'S' then
        text_file_read.read_line
      when 'B', 'T' then
        text_file_read.read_line
      else
      end
      from
        text_file_read.read_line
      until text_file_read.end_of_input
      loop
        text_file_read.read_line
      end
    end
    min_age : INTEGER is 5
feature {ANY}
  match(n: STRING) : BOOLEAN is
    do
      Result := name = n
    end
end

deferred class PEOPLE
feature
  set_name(n: STRING) is
    require
      n /= Void
    deferred
    end
feature {NONE}
  name : STRING
  age : INTEGER
invariant
  age > 0
  age < 200
end

-----
//CODE LIST END
```

The result shows that Spark can parse syntax of Eiffel and generates files correctly.

Diagram Files Generating

We composed a mod file for testing of diagram generation.

```
@CLASSDIAGRAM student1.ps2 PEOPLE STUDENT ATHLETE @VERTICAL
@CLASSDIAGRAM student2.ps2 PEOPLE STUDENT @HEAD
@CLASSDIAGRAM student3.ps2 STUDENT @METHOD
@CLASSDIAGRAM student4.ps2 STUDENT @ATTRIBUTE
@CLASSDIAGRAM student5.ps2 STUDENT @CONCISE
@CLASSDIAGRAM student6.ps2 STUDENT @BRIEF

@CLASS PEOPLE
@METHOD set_name (n:STRING)
@VAR name STRING
@VAR num STRING

@CLASS STUDENT
@INHERIT PEOPLE
@IMPLEMENT ATHLETE
@METHOD match (n:STRING) BOOLEAN
@VAR num INTEGER
@METHOD run

@INTERFACE ATHLETE
@METHOD run
```

According to this mod file, Spark should generate 6 diagrams. Figure 5.5 is drawn vertically and includes all these three classes or interface.

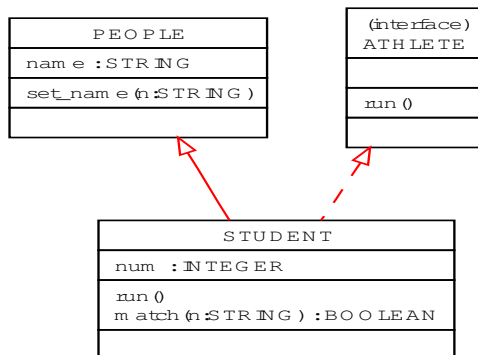


Figure 5.5: Vertical testing



Figure 5.6: Horizontal testing

Figure 5.6 is drawn horizontally and shows the class head only.

Figure 5.7 shows the methods of a class only and figure 5.8 shows the attributes of a class only.

Figure 5.9 and figure 5.10 show classes in concise form and brief form respectively.

The results shows that Spark satisfies the design requirements.

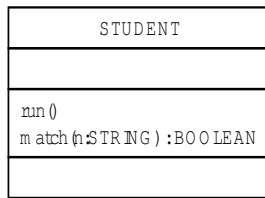


Figure 5.7: Method only testing

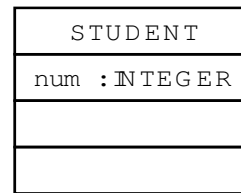


Figure 5.8: Attribute only testing

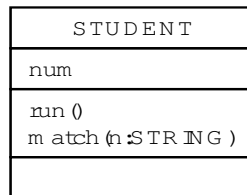


Figure 5.9: Concise form testing

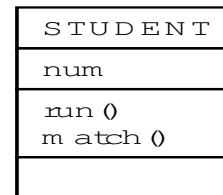


Figure 5.10: Brief form testing

5.5.3 Integration Testing

Our case study itself is perfect integration testing for Spark. The result that program files can be compiled successfully and the graphic notation files is successfully included in the documentation (see Chapter 4) shows that Spark satisfies the design requirements.

5.5.4 System Testing

We conducted all testing mentioned above on MS Windows and Macintosh. The same results show that Spark is platform-independent and Spark satisfies the design requirements.

Chapter 6

Conclusion and Future Work

In this thesis, we presented a new programming paradigm, object-oriented literate programming, which combines several existed significant ideas and is used to construct object-oriented programs in literature style. A set of software tools, Spark, is implemented to support this technique. So far the implementation of Spark altogether contains about hand-written 5000 lines and consists of the following four parts:

- The front end of Spark for Lime (1200 lines).
- The front end of Spark for Eiffel (2500 lines).
- The front end of Spark for perl (800 lines).
- The back end of Spark (500 lines).

Chapter 3 introduced OOLP and all features of Spark, from which we can see how programmers can enjoy the freedom of choosing the combination of languages to develop their software. Chapter 4 gave a case study, Transit Information System, implemented with this technique. It turns out that object-oriented software applications can be expressed in literate style well. In addition, programmers do not need to worry about the graphical notations as well as any extra cost spent on training and tools. Of course, more complex examples are still needed to validate this technique.

Spark itself needs further development to make it more complete and useful. One of the most important things is to build more front ends for Spark in order to adapt more object-oriented programming languages, such as Java and C#.

More graphical notations are need to be supported. So far, Spark can only generate class diagrams. In fact, further development could let Spark have the ability to produce dynamic diagrams, such as sequence diagram and statechart, which can make software documentation more expressive and more complete.

Appendix A

Installation

In order to run the study case presented in this thesis it is necessary to get hold of the following five separate tools. All of them have setup programs as well as installation instructions that can be found on the Web.

A.1 Perl

Perl is a dynamic programming language created by Larry Wall. As an open source software, every body can download its latest version for free from www.perl.com/download.csp.

Perl is necessary, because Spark is developed entirely in this language. The version 5.8.8.820 is employed in the testing of Spark.

A.2 Graphviz

Graphviz is a package of open source tools initiated by AT&T Research Labs for drawing graphs specified in DOT language scripts. Since it is free software licensed under the Common Public License, every one can download it for free from www.ryandesign.com/graphviz (for Mac OS) and from www.graphviz.org/Download_windows.php (for MS Windows)

Graphviz is necessary, because the back end of Spark depends on it. The version 2.12 is employed in the testing of Spark.

A.3 AsciiDoc

AsciiDoc is a text document format for writing short documents, articles, books and UNIX man pages. Its files can be translated to HTML and DocBook markups easily. Free use of AsciiDoc is granted under the terms of the GNU General Public License, so every one can download the latest version for free from www.methods.co.nz/asciidoc/downloads.html.

AsciiDoc is necessary, because it is used as the document formatting language in the study case. The version 8.2.1 is employed in the testing of Spark.

A.4 Python

Python is a dynamic object-oriented programming language. As an OSI certified open source software, every body can download its latest version for free from www.python.org/download.

Python is necessary, because AsciiDoc depends on it. The version 2.5.1 is employed in the testing of Spark.

A.5 SmartEiffel

SmartEiffel is a free Eiffel compiler. It has been developed at the Lorraine Laboratory of Research in Information Technology and its Applications, an institute affiliated to the French National Institute for Research in Computer Science and Control. SmartEiffel can be downloaded for free from smarteiffel.loria.fr.

SmartEiffel is necessary, because Eiffel is chose as the programming language in the study case. The version 2.2 is employed in the testing of Spark.

Appendix B

Source Code of Case Study

```
== The Requirement of Transit Information System
In this project, we are asked to develop an information system for a
local train and bus service. Our customer, HPTA (Happy Passenger
Transit Authority), has no clear picture what it should do, except
to increase customer satisfaction and make traveling more
attractive. All the information we have goes as follows:

.....
- It will be used by passengers as well as by HPTA staff.
- Selected staff members would be allowed to update the information.
- Passengers should be able to enter their start and destination, a desired time, and get a bunch
of possible connections.
- Connections can be direct or with changing busses or trains.
- For each bus and train station, the information like opening hours and accessibility is maintained.
- Users can browse a list of all bus and train routes or check the details of a certain route.
- Some bus stops and train stops are conjoint, but some not.
- Trains have two-digit numbers and busses have three-digit numbers.
- Connections between trains and busses must have at least five minutes for the change.
.....

For simplicity, we assume that detours and delays do not occur, stops are never skipped.

== An Overview
The following picture is an overview of this transit information system. As the root class,
HPTA_TRANSIT_INFO is a subclass of SYSTEM, which is a predefined class in Eiffel and allow its
subclasses to execute system command. Class DATABASE is a deferred class, whose subclasses, such as
class FILE_DATABASE, are responsible for maintaining system data. Class CONNECTION_FINDER is also a
deferred class, whose subclasses, such as class PRIME_FINDER, are responsible for finding the
possible connections.

image:: hpta_transit_info.jpg [Class Family]
//$ HPTA_TRANSIT_INFO DATABASE FILE_DATABASE CONNECTION_FINDER PRIME_FINDER ROUTE STAFF STATION KNOT @VERTICAL

The purpose of the application is to maintain the system information, including local train or bus
service and the status of staffs, and provide users current public transit service information,
including possible connections, and routes.

== Dictionary
To understand the main terms used in the requirement, we create a dictionary.
.....
- passenger: a person, who want to get his or her destination by bus or train.
- staff: a person, who works for HPTA.
- start: a station, where a passenger begin his or her journey.
- destination: a station to which a passenger is going or directed.
```

```

- desired time: an interval, within which one want get to the destination from the start.
- connection: a sequence of stations.
- bus: a long motor vehicle for carrying passengers, usually along a fixed route.
- train: a series of connected railroad cars pulled or pushed by one or more locomotives.
- route: a course for buses or trains to travel from one station to another.
- opening hour: a time, at which the first vehicle departs.
- accessibility: a description of the running status of a station.
- update: a change of system information.
- browse: a display of the information of all routes.
- check: a detail show of a certain route information.
.....

== Identifying Classes
Basing on the requirements, we defined the classes as follows:

-----
class HPTA_TRANSIT_INFO
end
-----
HPTA_TRANSIT_INFO is identified as a class of the entire system.

-----
class STAFF
feature {NONE}
  number: INTEGER
  password: STRING
end
-----
STAFF is a class with attributes employee number and password. The requirements state that selected
staff members would be allowed update the system.

-----
class STATION
feature {NONE}
  name: STRING
  open: STRING
  accessibility: STRING
end
-----
STATION is a class with attributes name, opening hour, and accessibility.

-----
class ROUTE
feature {NONE}
  number: INTEGER
  stops: LINKED_LIST [STATION]
end
-----
ROUTE is a class with attributes station list and route number.

== Identifying Operations
All three operations listed in the directory belong naturally in the class HPTA_TRANSIT_INFO,
because they are dependent on the interface of the system.
.....
- login should belong in class STAFF, because it keep the secret of a certain staff.
.....

== Consulting The Library of Model
There is no suitable business model in our existing library, so we have to build this system from
the beginning.

== Applying Design Patterns
According to the requirements, our application needs to keep all system information and to calculate
possible connections. There exist so many different methods for these two tasks. Hence, we apply the
strategy design pattern. We declare two deferred classes

-----
deferred class DATABASE
end

```

and

```
deferred class CONNECTION_FINDER
end
```

Then, we define two private members for class HPTA_TRANSIT_INFO denoted by the class name followed by three dots as following:

```
feature {NONE}
  HPTA_TRANSIT_INFO...db: DATABASE
  finder: CONNECTION_FINDER
```

i.e.

```
image::hptal.jpg[attributes of class HPTA_TRANSIT_INFO]
//$ HPTA_TRANSIT_INFO @ATTRIBUTE
```

In this way, we can add new algorithms easily and even change mechanisms at runtime with the following private methods:

```
feature {NONE}
  HPTA_TRANSIT_INFO...set_finder(new_finder: CONNECTION_FINDER) is
    require
      new_finder /= Void
    do
      finder := new_finder
    ensure
      finder = new_finder
    end
```

and

```
feature {NONE}
  HPTA_TRANSIT_INFO...set_database(new_database: DATABASE) is
    require
      new_database /= Void
    do
      db := new_database
    ensure
      db = new_database
    end
```

Their preconditions require that the new comers are not invalid and their postconditions ensure that the private member db and finder are set correctly.

Class CONNECTION_FINDER describes the interface that is common to all concrete mechanisms as following:

```
image::connection.jpg[Class connection_finder]
//$ CONNECTION_FINDER @METHOD
```

```
feature {HPTA_TRANSIT_INFO}
  CONNECTION_FINDER...
  get_connection(dbase: DATABASE; start, destination: STRING; time: INTEGER): STRING is
    require
      start /= Void
      destination /= Void
      time >= 0
      dbase /= Void
    deferred
    end
```

Class DATABASE describes the interface that is common to all concrete data maintain mechanisms as following:

```
image::database.jpg[Class database]
//$ DATABASE @METHOD
```

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_first_bus_route: ROUTE is
    deferred
  end
```

This method can return the first bus route object and is used by class HPTA_TRANSIT_INFO and class CONNECTION_FINDER. Together with the following method, its clients can browse all bus routes one by one.

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_next_bus_route(route: ROUTE): ROUTE is
    deferred
  end
```

Similarly, we can browse all train routes by the following two methods:

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_first_train_route: ROUTE is
    deferred
  end
```

and

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... get_next_train_route(route: ROUTE): ROUTE is
    deferred
  end
```

Browsing all staff information is not necessary, but we need to find given staff object by the following method.

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... find_staff (num: INTEGER): STAFF is
    require
      num > 0
    deferred
  end
```

This method can return an STAFF object, whose employee number equals to the parameter num. It is because all employee number start from 1 that the precondition is added.

For convenience, we also provide a route finding method as follows:

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
  DATABASE... find_route (num: INTEGER): ROUTE is
    require
      num >= min_train_route_num
      num <= max_bus_route_num
    deferred
  end
```

The following method is the creation of class DATABASE and invoked by class HPTA_TRANSIT_INFO only.

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... make is
    deferred
  end
```

In order to update system information, class DATABASE also requires the interfaces of adding and deleting as following:

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... add_route(new_route: ROUTE) is
    require
      new_route /= Void
    deferred
  end
```

and

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... delete_route(route: ROUTE) is
    require
      route /= Void
    deferred
  end
```

These two methods can add or delete a certain route to or from this system respectively and is called by class HPTA_TRANSIT_INFO only.

Similarly, class HPTA_TRANSIT_INFO also can add or delete a certain staff by the following two methods:

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... add_staff(new_staff: STAFF) is
    require
      new_staff /= Void
    deferred
  end
```

and

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... delete_staff(staff: STAFF) is
    require
      staff /= Void
    deferred
  end
```

As long as some system information is updated, DATABASE object must be informed to save the change by the following method.

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... do_save is
    deferred
  end
```

According to the requirements, only selected members can update the system. We define that when the database is locked, only the user, who knows both employee number and password, can conduct an update.

```
feature {HPTA_TRANSIT_INFO}
  DATABASE... is_locked: BOOLEAN is
    deferred
  end
```

The subclasses of these two deferred classes implement each concrete behavior mentioned above.

The following four constants are used to point out the bound of route number

```
feature {NONE}
  DATABASE... max_bus_route_num: INTEGER is 999
  min_bus_route_num: INTEGER is 100
```

```

max_train_route_num: INTEGER is 99
min_train_route_num: INTEGER is 10
-----

== Algorithms Design
=== File Database
For simplicity, we save the system information in a file named "sys_info.txt". So we define a
subclass of class DATABASE, FILE_DATABASE as following:
-----

class FILE_DATABASE
inherit DATABASE
feature {NONE}
    file_name: STRING is "sys_info.txt"
end
-----

i.e.

image::databases.jpg[the hirarchy of databases]
//$ DATABASE FILE_DATABASE @HEAD @VERTICAL

class FILE_DATABASE keep bus routes, train routes and staffs with LINKED_LIST as following:
-----

feature {NONE}
    FILE_DATABASE...train_routes: LINKED_LIST [ROUTE]
    bus_routes: LINKED_LIST [ROUTE]
    employees: LINKED_LIST [STAFF]
-----

now, class FILE_DATABASE becomes:

image::file_database.jpg[attributes of file database]
//$ FILE_DATABASE @ATTRIBUTE

The creation of FILE_DATABASE is method make
-----

create FILE_DATABASE...make
-----

The main task of make is to initialize this three list
-----

feature {HPTA_TRANSIT_INFO}
    FILE_DATABASE...make is
        do
            create employees.make
            create bus_routes.make
            create train_routes.make
            load
        ensure
            employees /= Void
            bus_routes /= Void
            train_routes /= Void
        end
-----

and to load the system information for that file:
-----

feature {NONE}
    FILE_DATABASE...load is
        local
            input_string : STRING
            text_file_read: TEXT_FILE_READ
            text_file_write: TEXT_FILE_WRITE
            split: ARRAY[STRING]
            new_staff: STAFF
            route: ROUTE
        do
            create text_file_read.connect_to(file_name)
            if text_file_read.is_connected then
                from text_file_read.read_line

```

```

until text_file_read.end_of_input
loop
  if text_file_read.last_string.upper = 1 then
    inspect text_file_read.last_string.first.to_upper
    when 'S' then
      text_file_read.read_line
      input_string := text_file_read.last_string.twin
      split := input_string.split
      create new_staff.make (split.first.to_integer, split.last)
      employees.add_last(new_staff)
    when 'B', 'T' then
      text_file_read.read_line
      input_string := text_file_read.last_string.twin
      split := input_string.split
      route := find_route(split.item(4).to_integer)
      if route = Void then
        create route.make(split.item(4).to_integer)
        route.add_station(split.first, split.item(2), split.item(3), split.last)
        if split.item(4).to_integer > max_train_route_num then
          bus_routes.add_last(route)
        else
          train_routes.add_last(route)
        end
      else
        route.add_station(split.first, split.item(2), split.item(3), split.last)
      end
    end
  else
    end
  end
  text_file_read.read_line
end
text_file_read.disconnect
else
  create text_file_write.connect_to(file_name)
  if text_file_write.is_connected then
    text_file_write.disconnect
  end
end
end
end

```

By the following method, one can get the specific route object.

```

feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
FILE_DATABASE...find_route (num: INTEGER): ROUTE is
  local
    i: INTEGER
    route: ROUTE
  do
    if num > max_train_route_num then
      from i := bus_routes.lower
      until i > bus_routes.upper or else bus_routes.item(i).match(num)
      loop
        i := i+1
      end
      if i <= bus_routes.upper then
        route := bus_routes.item(i)
      end
    else
      from i := train_routes.lower
      until i > train_routes.upper or else train_routes.item(i).match(num)
      loop
        i := i+1
      end
      if i <= train_routes.upper then
        route := train_routes.item(i)
      end
    end
  end
  Result := route

```

end

Similarly, using the following method, one can get the staff with such employee number:

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
FILE_DATABASE... find_staff (num: INTEGER): STAFF is
  local
    i: INTEGER
    staff: STAFF
  do
    from i := employees.lower
    until i > employees.upper or else employees.item(i) .match(num)
    loop
      i := i+1
    end
    if i <= employees.upper then
      staff := employees.item(i)
    end
    Result := staff
  end
end
```

By the following four methods, one can browse all train routes and bus routes:

```
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
FILE_DATABASE... get_first_bus_route: ROUTE is
  local
    route: ROUTE
  do
    if not bus_routes.is_empty then
      route := bus_routes.first
    end
    Result := route
  end
end
get_next_bus_route(route: ROUTE): ROUTE is
  require
    bus_routes.index_of(route) > 0
  local
    next_route: ROUTE
  do
    if bus_routes.index_of(route) < bus_routes.upper then
      next_route := bus_routes.item(bus_routes.index_of(route)+1)
    end
    Result := next_route
  end
end
get_first_train_route: ROUTE is
  local
    route: ROUTE
  do
    if not train_routes.is_empty then
      route := train_routes.first
    end
    Result := route
  end
end
get_next_train_route(route: ROUTE): ROUTE is
  require
    train_routes.index_of(route) > 0
  local
    next_route: ROUTE
  do
    if train_routes.index_of(route) < train_routes.upper then
      next_route := train_routes.item(train_routes.index_of(route)+1)
    end
    Result := next_route
  end
end
```

By the following method, HPTA_TRANSIT_INFO object can add an arbitrary route to this database

```

feature {HPTA_TRANSIT_INFO}
  FILE_DATABASE... add_route(new_route: ROUTE) is
    do
      if is_bus(new_route.get_number) then
        bus_routes.add_last(new_route)
      elseif is_train(new_route.get_number) then
        train_routes.add_last(new_route)
      end
    end
end

```

By the following method, HPTA_TRANSIT_INFO object can add a staff to this database

```

feature {HPTA_TRANSIT_INFO}
  FILE_DATABASE... add_staff(new_staff: STAFF) is
    do
      employees.add_last(new_staff)
    end
end

```

By the following method, HPTA_TRANSIT_INFO object can remove an arbitrary route from this database

```

feature {HPTA_TRANSIT_INFO}
  FILE_DATABASE... delete_route (route: ROUTE) is
    do
      if is_bus(route.get_number) then
        bus_routes.remove(bus_routes.index_of(route))
      elseif is_train(route.get_number) then
        train_routes.remove(train_routes.index_of(route))
      end
    end
end

```

By the following method, HPTA_TRANSIT_INFO object can remove a staff from this database

```

feature {HPTA_TRANSIT_INFO}
  FILE_DATABASE... delete_staff (staff: STAFF) is
    do
      employees.remove(employees.index_of(staff))
    end
end

```

In FILE_DATABASE, as long as employees is not empty, this database is locked, which means the user has to log in before updating.

```

feature {HPTA_TRANSIT_INFO}
  FILE_DATABASE... is_locked: BOOLEAN is
    do
      Result := not employees.is_empty
    end
end

```

Whenever the database is changed, it have to save the new data to the specific file by the following method:

```

feature {HPTA_TRANSIT_INFO}
  FILE_DATABASE... do_save is
    local
      file_2_write: TEXT_FILE_WRITE
      i: INTEGER
    do
      create file_2_write.connect_to(file_name)
      if file_2_write.is_connected then
        from i := employees.lower
        until i > employees.upper
        loop
          employees.item(i) .do_save(file_2_write)
          i := i + 1
        end
      end
    end
end

```

```

        from i := bus_routes.lower
        until i > bus_routes.upper
        loop
            bus_routes.item(i) .do_save(file_2_write)
            i := i + 1
        end
        from i := train_routes.lower
        until i > train_routes.upper
        loop
            train_routes.item(i) .do_save(file_2_write)
            i := i + 1
        end
        file_2_write.disconnect
    else
        io.put_string("Update database failed!%N")
    end
end
end

```

For convenience, we define the following two methods to tell if the current route is bus or train route:

```

feature {NONE}
    FILE_DATABASE...is_train(num: INTEGER): BOOLEAN is
    do
        Result := num >= min_train_route_num and num <= max_train_route_num
    end
    is_bus(num: INTEGER): BOOLEAN is
    do
        Result := num >= min_bus_route_num and num <= max_bus_route_num
    end

```

==== Conditional Shortest Path

According to the requirement that connections between trains and busses must have at least five minutes for the change, we have to consider bus station and train station as two different stations even they share the same name. In addition, we define a constant change_time in class CONNECTION_FINDER, whose subclasses need it.

```

feature {NONE}
    CONNECTION_FINDER...change_time: INTEGER is 5

```

For convenience, we assume that a bus needs 2 minutes to get to the second stop and a train needs only 1 minute. So we also define the following two members in class CONNECTION_FINDER.

```

feature {NONE}
    CONNECTION_FINDER...train_time: INTEGER is 1
    bus_time: INTEGER is 2

```

PRIME_FINDER is one of the subclasses of CONNECTION_FINDER

```

inherit
    PRIME_FINDER...CONNECTION_FINDER

```

i.e.

```

image::connection_finder.jpg[hierarchy of class HPTA_TRANSIT_INFO]
//$ CONNECTION_FINDER PRIME_FINDER @HEAD @VERTICAL

```

Our first algorithm, PRIME_FINDER, is that starting from the start stations, including both bus station and train station, we search for all direct neighbors one after another and calculate their time respectively. In this way, as long as we found the destination as the next neighbor or no more new neighbors before get to the destination, our searching work is done.

To implement this algorithm, we declare list in class PRIME_FINDER

```

feature {NONE}

```

```
PRIME_FINDER...stop_list: LINKED_LIST [KNOT]
```

Every node of this list record the following information:

```
image::knot.jpg[attribute of class KNOT]
//$ KNOT @ATTRIBUTE
```

```
feature {NONE}
  KNOT...station: STATION
```

Form the start down to the destination, as long as the station is found as a valid neighbor, it will be set in a KNOT object by the following method.

```
feature {PRIME_FINDER}
  KNOT...set_station(value: STATION) is
  do
    station := value
  end
```

Of course, class KNOT requires PRIME_FINDER object give a non Void value.

```
feature {PRIME_FINDER}
  KNOT...get_station: STATION is
  do
    Result := station
  end
```

After searching, PRIME_FINDER object can get the record of station by the above method.

```
feature {NONE}
  KNOT...number: INTEGER
```

The number of KNOT object keeps the route number of the station and is set by the following method:

```
feature {PRIME_FINDER}
  KNOT...set_number(value: INTEGER) is
  require
    value >= 0
    value <= 999
  do
    number := value
  end
```

According to the requirement that train route number is a two-digit number and bus route number is a three-digit number, we set a precondition like that for this method.

```
feature {PRIME_FINDER}
  KNOT...get_number: INTEGER is
  do
    Result := number
  end
```

The above method can tell PRIME_FINDER object the route, to which this station belongs.

```
feature {NONE}
  KNOT...time: INTEGER
```

Member time records the total time needed from start and is set by the following method

```
feature {PRIME_FINDER}
  KNOT...set_time(value: INTEGER) is
  require
    value >= 0
  do
```

```

    time := value
end

```

The time of start node is 0 and the time of destination is desire time plus one, so here KNOT object requires a nonnegative number.

```

feature {PRIME_FINDER}
  KNOT...get_time: INTEGER is
  do
    Result := time
  end
end

```

The above method is used to provide time for PRIME_FINDER object.

```

feature {NONE}
  KNOT...pred: INTEGER

```

This member is used to record the index of last stop in this list. The pred of start is -1. That the pred of two destination are all -1 means that there is no possible connection between the start and the destination.

PRIME_FINDER object set this member by the following method:

```

feature {PRIME_FINDER}
  KNOT...set_pred(value: INTEGER) is
  do
    pred := value
  end
end

```

and get the value of this member by the following method:

```

feature {PRIME_FINDER}
  KNOT...get_pred: INTEGER is
  do
    Result := pred
  end
end

```

Then, how can we judge if this node should be check for new neighbors? we define the member status in class KNOT.

```

feature {NONE}
  KNOT...status: INTEGER

```

If there is no more new neighbors can be found for the current station, this member should be set as permanent, which is a constant of class KNOT;

```

feature {PRIME_FINDER}
  KNOT...permanent: INTEGER is 1

```

otherwise, member status should be set as tentative, which is another constant of class KNOT.

```

feature {PRIME_FINDER}
  KNOT...tentative: INTEGER is 0

```

This member can be set by the following method

```

feature {PRIME_FINDER}
  KNOT...set_status(value: INTEGER) is
  require
    value >= tentative
    value <= permanent
  do
    status := value
  end
end

```

and get by the following method

```
feature {PRIME_FINDER}
  KNOT...get_status: INTEGER is
  do
    Result := status
  end
```

Method make is the creation of class KNOT

```
creation {PRIME_FINDER}
  KNOT...make
```

and its main task is to initialize this object with the given parameters as following:

```
feature {PRIME_FINDER}
  KNOT...make(sn: STATION; num, t, ss, pr: INTEGER) is
  do
    set_station(sn)
    set_number(num)
    set_time(t)
    set_status(ss)
    set_pred(pr)
  end
```

Every node is added into the list by the following method:

```
feature {NONE}
  PRIME_FINDER...add_node (pr: INTEGER; s: STATION; t, num: INTEGER) is
  require
    t >= 0
  local
    node: KNOT
  do
    create node.make(s, num, t, node.tentative, pr)
    if s = Void then
      node.set_status(node.permanent)
    end
    stop_list.add_last(node)
  end
```

If the station is Void, then the new node will be considered as dead.

The logic of possible connection finding is implemented mainly in the following method.

```
feature {HPTA_TRANSIT_INFO}
  PRIME_FINDER...
  get_connection(dbase: DATABASE; start, destination: STRING; time: INTEGER): STRING is
  require else
    stop_list.upper = 0
  local
    connection, cur_station: STRING
    node: KNOT
    i, monitor: INTEGER
    is_end, break: BOOLEAN
  do
    connection := ""
    desire_time := time
    add_bus_train_station(dbase, destination, desire_time+1)
    add_bus_train_station(dbase, start, 0)

    i := 3
    cur_station := start.twin

  from
    until is_end or else cur_station = Void
```

```

loop
  monitor := stop_list.upper
  find_neighbor(dbase, cur_station, i)
  if monitor = stop_list.upper then
    if stop_list.item(i) /= Void then
      stop_list.item(i) .set_status(node.permanent)
    end
  end
  is_end := True

  from
  until break or else i > stop_list.upper
  loop
    if stop_list.item(i) /= Void then
      node := stop_list.item(i)
      if node.get_status = node.tentative and node.get_station /= Void then
        cur_station := node.get_station.get_name
        is_end := False
        break := True
      end
    end
    if not break then
      i := i + 1
    end
  end

  if break then
    break := False
  end
end

connection := get_connection_mes(1)
connection := connection + get_connection_mes(2)

if connection.same_as("") then
  connection := "There is no connection from your start"
  + " to your destination in such time."
end

Result := connection
ensure
  Result /= Void
end

```

The first parameter provides the source of data; the second and third parameters are the names of start station and destination station respectively; the last parameter is the desire time, which will be used to set the private member `desire_time`:

```

feature {NONE}
  PRIME_FINDER...desire_time: INTEGER

```

At the beginning of searching, we initialize the `stop_list` of a `PRIME_FINDER` object with four nodes, i.e. bus and train stations of destination followed by bus and train stations of start, using the following method:

```

feature {NONE}
  PRIME_FINDER...add_bus_train_station(dbase: DATABASE; name: STRING; time: INTEGER) is
  require
    name /= Void
    time >= 0
  local
    route: ROUTE
    station: STATION
    is_end: BOOLEAN
    num: INTEGER
  do
    route := dbase.get_first_bus_route

```

```

from
until is_end or route = Void
loop
  station := route.get_first_station
  from
  until is_end or station = Void
  loop
    if name.same_as(station.get_name) then
      is_end := True
    end
    if not is_end then
      station := route.get_next_station(station)
    end
  end
  if not is_end then
    route := dbase.get_next_bus_route(route)
  end
end
if not is_end then
  station := Void
end
if route /= Void then
  num := route.get_number
else
  num := 0
end
add_node(-1, station, time, num)

station := Void
is_end := False
route := dbase.get_first_train_route
from
until is_end or route = Void
loop
  station := route.get_first_station
  from
  until is_end or station = Void
  loop
    if name.same_as(station.get_name) then
      is_end := True
    end
    if not is_end then
      station := route.get_next_station(station)
    end
  end
  if not is_end then
    route := dbase.get_next_train_route(route)
  end
end
if not is_end then
  station := Void
end
if route /= Void then
  num := route.get_number
else
  num := 0
end
add_node(-1, station, time, num)
end

```

Then from the bus station of start, we try to find its direct neighbor by the following method:

```

feature {NONE}
  PRIME.FINDER... find_neighbor(dbase: DATABASE; sn: STRING; pr: INTEGER) is
  require
    sn /= Void
  local
    cost, index, switch: INTEGER

```

```

p_node, node: KNOT
route: ROUTE
station, last: STATION
name: STRING
break: BOOLEAN
do
  from switch := 0
  until switch > 1
  loop
    if switch = 0 then
      cost := bus_time
    else
      cost := train_time
    end
    if pr >= stop_list.lower and pr <= stop_list.upper then
      p_node := stop_list.item(pr)
    end
    if p_node /= Void then
      if p_node.get_station /= Void then
        if switch = 0 then
          if is_train(p_node.get_number) then
            cost := change_time + cost
          end
          route := dbase.get_first_bus_route
        else
          if is_bus(p_node.get_number) then
            cost := change_time + cost
          end
          route := dbase.get_first_train_route
        end
      from
      until route = Void
      loop
        station := route.get_first_station
        last := station
        from
        until station = Void or break
        loop
          name := station.get_name.twin
          if name /= Void and name.is_equal(sn) then
            if not last.get_name.is_equal(name) then
              index := get_index(last, route.get_number)
              if index >= 0 then
                node := stop_list.item(index)
                if node.get_station /= Void then
                  if is_train(node.get_number) then
                    if node.get_time > p_node.get_time + cost then
                      node.set_pred(pr)
                      node.set_time(p_node.get_time + cost)
                      node.set_number(route.get_number)
                    end
                  end
                end
              end
            else
              add_node(pr, last, p_node.get_time+cost, route.get_number)
            end
          end
          last := route.get_next_station(station)
          if last /= Void then
            index := get_index(last, route.get_number)
            if index >= 0 then
              node := stop_list.item(index)
              if node.get_station /= Void then
                if is_train(node.get_number) then
                  if node.get_time > p_node.get_time + cost then
                    node.set_pred(pr)
                    node.set_time(p_node.get_time + cost)
                    node.set_number(route.get_number)
                  end
                end
              end
            end
          end
        end
      end
    end
  end
end

```

```

                                end
                                end
                                end
                                else
                                add_node(pr, last, p_node.get_time+cost, route.get_number)
                                end
                                end
                                break := True
                                else
                                last := station;
                                station := route.get_next_station(station)
                                end
                                end
                                break := False
                                if switch = 0 then
                                route := dbase.get_next_bus_route(route)
                                else
                                route := dbase.get_next_train_route(route)
                                end
                                end
                                end
                                end
                                switch := switch + 1
                                end
                                end
                                end

```

For convenience, we define the following two methods to tell if the current route is train or bus:

```

feature {NONE}
  PRIME.FINDER...is_train(num: INTEGER): BOOLEAN is
  do
    Result := num >= 10 and num <= 99
  end
end

```

and

```

feature {NONE}
  PRIME.FINDER...is_bus(num: INTEGER): BOOLEAN is
  do
    Result := num >= 100 and num <= 999
  end
end

```

The following method is used to get the index of a certain station in the list; if the target station is not in the list, -1 will be return.

```

feature {NONE}
  PRIME.FINDER...get_index(s: STATION; num: INTEGER): INTEGER is
  require
    s /= Void

  local
    ind, i: INTEGER
    node: KNOT
    name: STRING
  do
    ind := -1
    from i := stop_list.lower
    until i > stop_list.upper
    loop
      node := stop_list.item(i)
      if node.get_station /= Void then
        name := node.get_station.get_name
        if name.is_equal(s.get_name) then
          if is_bus(num) and is_bus(node.get_number) then
            ind := i
          elseif is_train(num) and is_train(node.get_number) then
            ind := i
          end
        end
      end
    end
  end
end

```

```

        end
      end
    end
    i := i + 1
  end
  Result := ind
end

```

When the searching is done, we can get the information of possible connections by the following method:

```

feature {NONE}
  PRIME_FINDER... get_connection_mes(index: INTEGER): STRING is
  require
    index >= 0
  local
    node: KNOT
    mes: STRING
  do
    mes := ""
    node := stop_list.item(index)
    if node /= Void then
      if node.get_station /= Void then
        if node.get_pred /= -1 and node.get_time <= desire_time then
          mes := "-No." + node.get_number.to_string + "->"
            + node.get_station.get_name + " in "
            + node.get_time.to_string + " minutes%N"
          node := stop_list.item(node.get_pred)
          from
            until node = Void or else node.get_station = Void or else node.get_pred = -1
            loop
              mes := "-No." + node.get_number.to_string + "->"
                + node.get_station.get_name + mes
              node := stop_list.item(node.get_pred)
            end
          if node /= Void then
            if node.get_station /= Void then
              mes := "%N" + node.get_station.get_name + mes
            end
          else
            mes := ""
          end
        end
      end
    end
    Result := mes
  end
end

```

The creation of PRIME_FINDER is method make

```

creation {ANY} PRIME_FINDER... make

```

it is defined as following:

```

feature {HPTA_TRANSIT_INFO}
  PRIME_FINDER... make is
  do
    create stop_list.make
  ensure
    stop_list /= Void
  end

```

Now, let us talk about the root class HPTA_TRANSIT_INFO.

```

image::hpta2.jpg [methods of class HPTA_TRANSIT_INFO]
//$ HPTA_TRANSIT_INFO @METHOD

```

The creation of class HPTA_TRANSIT_INFO is make

```
create HPTA_TRANSIT_INFO...make
```

Its main task is to initialize the database and connection finder, and then run the whole system:

```
feature {ANY}
  HPTA_TRANSIT_INFO...make is
    local
      prime_finder: PRIME_FINDER
      file_database: FILE_DATABASE
    do
      create file_database.make
      set_database(file_database)
      create prime_finder.make
      set_finder(prime_finder)
      run
    end
```

In order to increase customer satisfaction, we run the system by a series of menus

```
feature {NONE}
  HPTA_TRANSIT_INFO...run is
    do
      from
        until io.last_character.to_upper = 'Q'
        loop
          menu
            io.read_character
            io.put_new_line
            inspect io.last_character.to_upper
            when 'U' then do_update
            when 'I' then do_inquire
            else
              end
          end
        end
      end
```

In order to use OS command, we let class HPTA_TRANSIT_INFO be a subclass of class SYSTEM, which is a predefined class in Eiffel.

```
inherit
  HPTA_TRANSIT_INFO...SYSTEM
```

Method menu is the main menu of the interface of this system and

```
feature {NONE}
  HPTA_TRANSIT_INFO...menu is
    do
      execute_command_line("cls")
      io.put_string("[
        *****
        Welcome to HPTA
        *****
        U Update System Information
        I Inquire about Transit Information
        Q Quit

        Enter menu choice:
      ]")
    end
```

This is the main menu and there are two items in it, through which users can either update or

inquire system information. The first line of the method body is used to clear the screen.

If users chose the first menu item, they are going to enter the following menu, i.e. update_menu:

```
feature {NONE}
  HPTA_TRANSIT_INFO...update_menu is
  do
    execute_command_line(" cls ")
    io.put_string("[
    *****
    Welcome to HPTA
    *****
    A Add
    D Delete
    G Go back

    Enter menu choice:
    ]")
  end
```

In this menu, users can add new information, such as staffs and stations, as follow:

```
feature {NONE}
  HPTA_TRANSIT_INFO...add_menu is
  do
    execute_command_line(" cls ")
    io.put_string("[
    *****
    Welcome to HPTA
    *****
    S Add a station
    E Add a staff
    G Go back

    Enter menu choice:
    ]")
  end
```

Follows the logic of method do_add:

```
feature {NONE}
  HPTA_TRANSIT_INFO...do_add is
  local
    employee: STAFF
    id: INTEGER
    input, name, password, open, access, last: STRING
    is_end: BOOLEAN
    route: ROUTE
  do
    from
    until is_end
    loop
      add_menu
      io.read_line
      input := io.last_string.twin
      io.put_new_line
      if not input.is_empty then
        inspect input.first.to_upper
        when 'G' then is_end := True
        when 'S' then
          io.put_string("%NEnter station name: ")
          io.read_line
          name := io.last_string.twin
          io.put_string("%NEnter open hour: ")
          io.read_line
          open := io.last_string.twin
          io.put_string("%NEnter its accessibility: ")
```

```

        io.read_line
        access := io.last_string.twin
        io.put_string("%NEnter route number: ")
        io.read_line
        id := io.last_string.to_integer
        io.put_string("%NEnter the name of its last station: ")
        io.read_line
        last := io.last_string.twin
        route := db.find_route(id)
        if route = Void then
            create route.make(id)
            route.add_station(name, access, open, last)
            db.add_route(route)
        else
            route.add_station(name, access, open, last)
        end
    when 'E' then
        io.put_string("%NEnter your ID: ")
        io.read_line
        id := io.last_string.to_integer
        io.put_string("%NEnter your password: ")
        io.read_line
        password := io.last_string.twin
        create employee.make(id, password)
        db.add_staff(employee)
    else
        end
    end
end
end
end

```

they can also delete those information as follow:

```

feature {NONE}
  HPTA_TRANSIT_INFO...delete_menu is
  do
    execute_command_line(" cls ")
    io.put_string("[
      *****
      Welcome to HPTA
      *****
      S Delete a station
      E Delete a staff
      R Delete a route
      G Go back

      Enter menu choice:
    ]")
  end
end

```

Follows the logic of method do_delete:

```

feature {NONE}
  HPTA_TRANSIT_INFO...do_delete is
  local
    is_end: BOOLEAN
    num: INTEGER
    staff: STAFF
    route: ROUTE
    input, name: STRING
  do
    from
    until is_end
    loop
      delete_menu
      io.read_line
      input := io.last_string.twin
    end
  end
end

```

```

io.put_new_line
if not input.is_empty then
  inspect input.first.to_upper
  when 'G' then is_end := True
  when 'S' then
    io.put_string("%NEnter route number: ")
    io.read_line
    num := io.last_string.to_integer
    io.put_string("%NEnter station name: ")
    io.read_line
    name := io.last_string.twin
    route := db.find_route(num)
    if route /= Void then
      route.remove_station(name)
    else
      io.put_string ("%NNNo such a station%N")
      io.read_line
    end
  when 'R' then
    io.put_string("%NEnter route number: ")
    io.read_line
    num := io.last_string.to_integer
    route := db.find_route(num)
    if route /= Void then
      db.delete_route(route)
    else
      io.put_string ("%NNNo such a station%N")
      io.read_line
    end
  when 'E' then
    io.put_string("%NEnter ID: ")
    io.read_line
    num := io.last_string.to_integer
    staff := db.find_staff(num)
    if staff /= Void then
      db.delete_staff(staff)
    else
      io.put_string ("%NNNo such a staff%N")
      io.read_line
    end
  else
    end
end
end
end
end
end

```

According to the requirement, only authorized staffs can do such things, so this system will ask the user to log in the system before he or she enter the update menu. The following method `do.update` has the logic to require the user to enter his or her employee number and password first.

```

feature {NONE}
  HPTA_TRANSIT_INFO...do.update is
    local
      id: INTEGER
      passed, is_end: BOOLEAN
      password, input: STRING
      staff: STAFF
    do
      io.read_line
      if db.is_locked then
        io.put_string ("%NEnter employee ID: ")
        io.read_line
        id := io.last_string.to_integer
        staff := db.find_staff(id)
        if staff /= Void then
          io.put_string ("%NEnter password: ")

```

```

        io.read_line
        password := io.last_string.twin
        passed := staff.login(password)
    end
else
    io.put_string (" [
        The list of authorized staff is not empty,
        so please set authorization as soon as possible...
    ]")
    passed := True
    io.read_line
end
if passed then
    from
    until is_end
    loop
        update_menu
        io.read_line
        input := io.last_string.twin
        io.put_new_line
        if not input.is_empty then
            inspect input.first.to_upper
            when 'A' then do_add
            when 'D' then do_delete
            when 'G' then is_end := True
            else
                end
            end
        end
    end
    db.do_save
else
    io.put_string ("%NLogin failed!%N")
    io.read_line
end
end
end

```

The actual logging responsibility is assigned to class STAFF as public feature to class HPTA_TRANSIT_INFO:

```

feature {HPTA_TRANSIT_INFO}
    STAFF...login(passwd: STRING): BOOLEAN is
        require
            passwd /= Void
        do
            Result := password.is_equal(passwd)
        end
end

```

If the result is True, the user can continue his or her update, otherwise, this system will remain on the main menu.

If users chose the second menu item of the main menu, they will enter the following query menu without any bother, because the requirement says that any one can have access to the transit information.

```

feature {NONE}
    HPTA_TRANSIT_INFO...inquire_menu is
        do
            execute_command_line(" cls ")
            io.put_string (" [
                *****
                Welcome to HPTA
                *****
                F Find a possible connection
                S Show a route
                B Browse all routes
                G Go back

                Enter menu choice:
            ]")
        end
end

```

```

end
-----
The first item of this menu is used for users to find a possible connection. Following the
logic of method do_inquire, users are required to enter their start, destination, as well as
their desire time.
-----
feature {NONE}
  HPTA_TRANSIT_INFO...do_inquire is
    local
      input, start, dest: STRING
      is_end: BOOLEAN
      num, time: INTEGER
      route: ROUTE
    do
      from
      until is_end
      loop
        inquire_menu
        io.read_line
        input := io.last_string.twin
        io.put_new_line
        if not input.is_empty then
          inspect input.first.to_upper
          when 'B' then
            from route := db.get_first_bus_route
            until route = Void
            loop
              route.show
              route := db.get_next_bus_route(route)
            end
            from route := db.get_first_train_route
            until route = Void
            loop
              route.show
              route := db.get_next_train_route(route)
            end
            io.put_string ("%N%NStrike any key to continue...")
            io.read_line
          when 'F' then
            io.put_string ("%NEnter the station name of your start: ")
            io.read_line
            start := io.last_string.twin
            io.put_string ("%NEnter the station name of your destination: ")
            io.read_line
            dest := io.last_string.twin
            io.put_string ("%NEnter your desire time(in minutes): ")
            io.read_line
            time := io.last_string.to_integer
            io.put_string (finder.get_connection(db, start, dest, time))
            io.put_string ("%N%NStrike any key to continue...")
            io.read_line
          when 'S' then
            io.put_string ("Input the route number (10 - 999): ")
            io.read_line
            num := io.last_string.to_integer
            route := db.find_route(num)
            if route /= Void then
              route.show
            else
              io.put_string ("Sorry there is no such a route")
            end
            io.put_string ("%N%NStrike any key to continue...")
            io.read_line
          when 'G' then is_end := True
          else
            end
        end
      end
    end
end

```

```

        end
    end

```

Now, it is time to implement the methods of class ROUTE

```

image::Route.jpg[class ROUTE]
//$ ROUTE @METHOD

```

The creation of ROUTE is make, which can be invoke by class HPTA_TRANSIT_INFO

```

creation ROUTE...make

```

The main task of make is initialize the route number and station list

```

feature {HPTA_TRANSIT_INFO}
    ROUTE...make (num: INTEGER) is
        require
            num > 9
            num < 1000
        do
            number := num
            create stops.make
        ensure
            number = num
            stops /= Void
        end
end

```

According the requirement, route number must be two- or three-digit number, so we define the following invariant for class ROUTE.

```

invariant
    ROUTE...number > 9
    number < 1000
end

```

At any time, its client get route number by the following method:

```

feature {ANY}
    ROUTE...get_number : INTEGER is
        do
            Result := number
        end
end

```

also, by the following method to tell if the current route is which we want:

```

feature {ANY}
    ROUTE...match (num: INTEGER): BOOLEAN is
        do
            Result := num = number
        end
end

```

By the following method, its client adds new stations for this ROUTE object and at the same time set the name, the accessibility, the opening hour, and last station for this new station.

```

feature {HPTA_TRANSIT_INFO, DATABASE}
    ROUTE...add_station(new_name, access, open_hour, last_stop: STRING) is
        local
            new_station: STATION
            i: INTEGER
            last: STRING
        do
            from i := stops.lower
            until i > stops.upper or else stops.item(i) .match(new_name)
            loop
                i := i + 1
            end

```

```

        if i > stops.upper then
            create new_station.make(new_name, access, open_hour)
            last := last_stop.twin
            last.to_upper
            if last.same_as("NONE") then
                stops.add_first(new_station)
            else
                from i := stops.lower
                until i > stops.upper or else stops.item(i).match(last_stop)
                loop
                    i := i+1
                end
                if i <= stops.upper then
                    stops.add(new_station, i+1)
                else
                    create new_station.make(last_stop, access, open_hour)
                    stops.add_last(new_station)
                    stops.add_last(new_station)
                end
            end
        end
    end
end
end

```

HPTA_TRANSIT_INFO object removes a certain station by the following method, whose only parameter is the name of the target station.

```

feature {HPTA_TRANSIT_INFO}
    ROUTE...remove_station(name: STRING) is
        local
            i: INTEGER
        do
            from i := stops.lower
            until i > stops.upper or else stops.item(i).match(name)
            loop
                i := i + 1
            end
            if i <= stops.upper then
                stops.remove(i)
            end
        end
    end
end

```

The subclasses of CONNECTION_FINDER use the following two methods to visit all stations in this route

```

feature {CONNECTION_FINDER}
    ROUTE...get_first_station: STATION is
        local
            station: STATION
        do
            if stops.upper > 0 then
                station := stops.first
            end
            Result := station
        end
    get_next_station(station1: STATION): STATION is
        require
            station1 /= Void
        local
            station: STATION
        do
            if stops.index_of(station1) < stops.upper then
                station := stops.item(stops.index_of(station1)+1)
            end
            Result := station
        end
    end
end

```

Class ROUTE keep the secret of saving itself, so DATABASE object can call this method to

fulfill the task. Actually, such assignment is worth to discuss. Maybe should move to the subclasses of DATABASE, because only they know exactly how to save those data.

```
feature {DATABASE}
  ROUTE...do_save(file: TEXT_FILE_WRITE) is
    require
      file.is_connected
    local
      i: INTEGER
      tag, last: STRING
    do
      if number > 99 then
        tag := "b"
      else
        tag := "t"
      end
      last := "None"
      from i := stops.lower
      until i > stops.upper
      loop
        file.put_string(tag+"%N")
        stops.item(i).do_save(file)
        file.put_string(" " + number.to_string + " " + last + "%N")
        last := stops.item(i).get_name.twin
        i := i + 1
      end
    end
end
```

Similarly, the following method is responsible for showing the details of this route, but only class HPTA_TRANSIT_INFO know exactly how to display with interface, so this method should be move to class HPTA_TRANSIT_INFO.

```
feature {HPTA_TRANSIT_INFO}
  ROUTE...show is
    local i: INTEGER
    do
      if number > 99 then
        io.put_string ("%NBus route No.")
      else
        io.put_string ("%NTrain route No.")
      end
      io.put_integer (number)
      io.put_string (": ")
      from i := stops.lower
      until i > stops.upper
      loop
        stops.item(i).show
        if i < stops.upper then
          io.put_string (">")
        end
        i := i+1;
      end
      io.put_new_line
    end
end
```

Same problem can be found on the method show of class STATION

```
feature {ROUTE}
  STATION...show is
    do
      io.put_string (name)
    end
end
```

Now, let us look at the class STATION, whose creation is method make too,

```
create STATION...make
```

and defined as following:

```
feature {ROUTE}
  STATION...make (new_name, new_open, new_acc: STRING) is
    require
      new_name /= Void
      new_open /= Void
      new_acc /= Void
    do
      name := new_name.twin
      open := new_open.twin
      accessibility := new_acc.twin
    end
end
```

The main task of it is to initial these three features of class STATION. At any time, its client can visit these three features by the following methods:

```
feature {ROUTE, CONNECTION_FINDER}
  STATION...get_name: STRING is
    do
      Result := name.twin
    end
  get_acc: STRING is
    do
      Result := accessibility.twin
    end
  get_open: STRING is
    do
      Result := open.twin
    end
end
```

Similar with the method do_save of class ROUTE, this method should be moved into the subclasses of DATABASE.

```
feature {ROUTE}
  STATION...do_save(file: TEXT_FILE_WRITE) is
    require
      file.is_connected
    do
      file.put_string(name + " " + accessibility + " " + open)
    end
end
```

The same problem can be found on class STAFF

```
feature {DATABASE}
  STAFF...do_save(file: TEXT_FILE_WRITE) is
    require
      file.is_connected
    do
      file.put_string("s%N" + number.to_string + " " + password + "%N")
    end
end
```

We identify station with name only, i.e. if two stations share the same name, we assume they are the same station. Here case is insensitive.

```
feature {ROUTE}
  STATION...match (targetname: STRING): BOOLEAN is
    require
      targetname /= Void
    do
      Result := name.same_as (targetname)
    end
end
```

Now, let us talk about the implementation of class STAFF.

```
image::staff.jpg[class STAFF]
```

```

// $ STAFF @METHOD

The creation of class STAFF is make

-----
creation{ANY} STAFF...make
-----
it is defined as following:
-----
feature {ANY}
  STAFF...make (id: INTEGER; passwd: STRING) is
  require
    id >= 0
    passwd /=Void
  do
    number := id;
    password := passwd.twin
  ensure
    number >= 0
    password = passwd
  end
-----
its main task is initialize staff's id and password.

Method match is used to identify a certain staff and is defined as following:
-----
feature {ANY}
  STAFF...match (id: INTEGER): BOOLEAN is
  do
    Result := id = number
  end
-----
Any staff has an unique employee number, which is generated from 0, and a password,
which must not be Void:
-----
invariant
  STAFF...number_positive: number >= 0
  password_not_void: password /= Void
end
-----

So far, we have implement the system.

== Testing
== Updating system
When no staff is authorized, we try to update system information.
The result is

image::empty.jpg[empty]

Otherwise, we try to update system information. The system requires
ID and password for logging in as following:

image::updating.jpg[updating]

These results satisfy the design requirements.

== Browsing all routes
We try to browse the information of all routes as following:

image::browse.jpg[browse]

These result satisfies the design requirements.

== Finding connection
We try to find a connection between two stations as following:

image::connection.jpg[connection]

```

These result satisfies the design requirements.

== Strategy pattern

We construct a sample connection finder class and change the algorithm a run-time.

```
class FAKE_FINDER
inherit
    CONNECTION_FINDER
creation {ANY}
    make
feature {HPTA_TRANSIT_INFO}
    make is
        do
        end
    get_connection(start, destination: STRING; time: INTEGER): STRING is
        do
            Result := "This is a test%N";
        end
end
```

image::strategy.jpg[strategy]

These result satisfies the design requirements.

In order to give an integrated view for ones who are used to read code, we list all program code here.

Appendix C

Generated Code of Case Study

```
class KNOT
creation {PRIME_FINDER}
make
feature {PRIME_FINDER}
get_number : INTEGER is
do
    Result := number
end
make(sn: STATION; num, t, ss, pr: INTEGER) is
do
    set_station(sn)
    set_number(num)
    set_time(t)
    set_status(ss)
    set_pred(pr)
end
set_pred(value: INTEGER) is
do
    pred := value
end
get_time : INTEGER is
do
    Result := time
end
set_time(value: INTEGER) is
require
    value >= 0
do
    time := value
end
get_pred : INTEGER is
do
    Result := pred
end
permanent : INTEGER is 1
get_station : STATION is
do
    Result := station
end
tentative : INTEGER is 0
set_station(value: STATION) is
do
```

```

        station := value
    end
get_status : INTEGER is
do
    Result := status
end
set_status(value: INTEGER) is
    require
        value >= tentative
        value <= permanent
    do
        status := value
    end
set_number(value: INTEGER) is
    require
        value >= 10
        value <= 999
    do
        number := value
    end
feature {NONE}
    station : STATION
    number : INTEGER
    status : INTEGER
    time : INTEGER
    pred : INTEGER
end

deferred class DATABASE
feature {HPTA_TRANSIT_INFO}
do_save is
    deferred
    end
add_route(new_route: ROUTE) is
    require
        new_route /= Void
    deferred
    end
make is
    deferred
    end
delete_staff(staff: STAFF) is
    require
        staff /= Void
    deferred
    end
delete_route(route: ROUTE) is
    require
        route /= Void
    deferred
    end
add_staff(new_staff: STAFF) is
    require
        new_staff /= Void
    deferred
    end
is_locked : BOOLEAN is
    deferred
    end
feature {NONE}
    min_train_route_num : INTEGER is 10
    max_bus_route_num : INTEGER is 999
    max_train_route_num : INTEGER is 99
    min_bus_route_num : INTEGER is 100
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
    find_route(num: INTEGER) : ROUTE is
        require
            num >= min_train_route_num

```

```

        num <= max_bus_route_num
    deferred
    end
    find_staff(num: INTEGER) : STAFF is
    require
        num > 0
    deferred
    end
    get_next_bus_route(route: ROUTE) : ROUTE is
    deferred
    end
    get_first_bus_route : ROUTE is
    deferred
    end
    get_next_train_route(route: ROUTE) : ROUTE is
    deferred
    end
    get_first_train_route : ROUTE is
    deferred
    end
end

deferred class CONNECTION_FINDER
feature {HPTA_TRANSIT_INFO}
    get_connection(dbase: DATABASE; start, destination: STRING; time: INTEGER) : STRING is
    require
        start /= Void
        destination /= Void
        time >= 0
        dbase /= Void
    deferred
    end
feature {NONE}
    train_time : INTEGER is 1
    bus_time : INTEGER is 2
    change_time : INTEGER is 5
end

class STAFF
creation {ANY}
    make
feature {HPTA_TRANSIT_INFO}
    login(passwd: STRING) : BOOLEAN is
    require
        passwd /= Void
    do
        Result := password.is_equal(passwd)
    end
feature {DATABASE}
    do_save(file: TEXT_FILE_WRITE) is
    require
        file.is_connected
    do
        file.put_string("s%N" + number.to_string + "_" + password + "%N")
    end
feature {NONE}
    password : STRING
    number : INTEGER
feature {ANY}
    make(id: INTEGER; passwd: STRING) is
    require
        id >= 0
        passwd /= Void
    do
        number := id
        password := passwd.twin
    ensure
        number >= 0

```

```

        password = passwd
    end
    match(id: INTEGER) : BOOLEAN is
    do
        Result := id = number
    end
invariant
    number_positive : number >= 0
    password_not_void : password /= Void
end

class PRIME.FINDER
inherit
    CONNECTION.FINDER
creation {ANY}
make
feature {HPTA.TRANSIT.INFO}
make is
do
    create stop_list.make
ensure
    stop_list /= Void
end
get_connection(dbase: DATABASE; start, destination: STRING; time: INTEGER) : STRING is
require else
    stop_list.upper = 0
local
    connection : STRING
    cur_station : STRING
    node : KNOT
    i : INTEGER
    monitor : INTEGER
    is_end : BOOLEAN
    break : BOOLEAN
do
    connection := ""
    desire_time := time
    add_bus_train_station(dbase, destination, desire_time+1)
    add_bus_train_station(dbase, start, 0)
    i := 3
    cur_station := start.twin
from
until is_end or else cur_station = Void
loop
    monitor := stop_list.upper
    find_neighbor(dbase, cur_station, i)
    if monitor = stop_list.upper then
        if stop_list.item(i) /= Void then
            stop_list.item(i).set_status(node.permanent)
        end
    end
    is_end := True
from
until break or else i > stop_list.upper
loop
    if stop_list.item(i) /= Void then
        node := stop_list.item(i)
        if node.get_status = node.tentative and node.get_station /= Void then
            cur_station := node.get_station.get_name
            is_end := False
            break := True
        end
    end
    if not break then
        i := i + 1
    end
end
if break then

```

```

        break := False
    end
end
connection := get_connection_mes(1)
connection := connection + get_connection_mes(2)
if connection.same_as("") then
    connection := "_There_is_no_connection_from_your_start" + "_to_your_destination_in_such_time."
end
Result := connection
ensure
    Result /= Void
end
feature {NONE}
    get_index(s : STATION; num: INTEGER) : INTEGER is
    require
        s /= Void
    local
        ind : INTEGER
        i : INTEGER
        node : KNOT
        name : STRING
    do
        ind := - 1
        from
            i := stop_list.lower
        until i > stop_list.upper
        loop
            node := stop_list.item(i)
            if node.get_station /= Void then
                name := node.get_station.get_name
                if name.is_equal(s.get_name) then
                    if is_bus(num) and is_bus(node.get_number) then
                        ind := i
                    elseif is_train(num) and is_train(node.get_number) then
                        ind := i
                    end
                end
            end
            i := i + 1
        end
        Result := ind
    end
end
add_bus_train_station(dbase: DATABASE; name: STRING; time: INTEGER) is
    require
        name /= Void
        time >= 0
    local
        route : ROUTE
        station : STATION
        is_end : BOOLEAN
        num : INTEGER
    do
        route := dbase.get_first_bus_route
        from
            until is_end or route = Void
            loop
                station := route.get_first_station
                from
                    until is_end or station = Void
                    loop
                        if name.same_as(station.get_name) then
                            is_end := True
                        end
                        if not is_end then
                            station := route.get_next_station(station)
                        end
                    end
                end
            end
            if not is_end then

```

```

        route := dbase.get_next_bus_route(route)
    end
end
if not is_end then
    station := Void
end
if route /= Void then
    num := route.get_number
else
    num := 0
end
add_node(-1, station, time, num)
station := Void
is_end := False
route := dbase.get_first_train_route
from
until is_end or route = Void
loop
    station := route.get_first_station
    from
    until is_end or station = Void
    loop
        if name.same_as(station.get_name) then
            is_end := True
        end
        if not is_end then
            station := route.get_next_station(station)
        end
    end
    if not is_end then
        route := dbase.get_next_train_route(route)
    end
end
if not is_end then
    station := Void
end
if route /= Void then
    num := route.get_number
else
    num := 0
end
add_node(-1, station, time, num)
end
get_connection_mes(index: INTEGER) : STRING is
require
    index >= 0
local
    node : KNOT
    mes : STRING
do
    mes := "_"
    node := stop_list.item(index)
    if node /= Void then
        if node.get_station /= Void then
            if node.get_pred /= - 1 and node.get_time <= desire_time then
                mes := "_No." + node.get_number.to_string + "_>" + node.get_station.get_name + "_in_" + node
                node := stop_list.item(node.get_pred)
            from
            until node = Void or else node.get_station = Void or else
node.get_pred = - 1
        loop
            mes := "_No." + node.get_number.to_string + "_>" + node.get_station.get_name + mes
            node := stop_list.item(node.get_pred)
        end
        if node /= Void then
            if node.get_station /= Void then
                mes := "_%N" + node.get_station.get_name + mes
            end
        end
    end
end

```



```

        route := dbase.get_first_train_route
    end
    from
    until route = Void
    loop
        station := route.get_first_station
        last := station
        from
        until station = Void or break
        loop
            name := station.get_name.twin
            if name /= Void and name.is_equal(sn) then
                if not last.get_name.is_equal(name) then
                    index := get_index(last, route.get_number)
                    if index >= 0 then
                        node := stop_list.item(index)
                        if node.get_station /= Void then
                            if is_train(node.get_number) then
                                if node.get_time > p_node.get_time + cost then
                                    node.set_pred(pr)
                                    node.set_time(p_node.get_time + cost)
                                    node.set_number(route.get_number)
                                end
                            end
                        end
                    else
                        add_node(pr, last, p_node.get_time+cost, route.get_number)
                    end
                end
                last := route.get_next_station(station)
                if last /= Void then
                    index := get_index(last, route.get_number)
                    if index >= 0 then
                        node := stop_list.item(index)
                        if node.get_station /= Void then
                            if is_train(node.get_number) then
                                if node.get_time > p_node.get_time + cost then
                                    node.set_pred(pr)
                                    node.set_time(p_node.get_time + cost)
                                    node.set_number(route.get_number)
                                end
                            end
                        end
                    else
                        add_node(pr, last, p_node.get_time+cost, route.get_number)
                    end
                end
                break := True
            else
                last := station
                station := route.get_next_station(station)
            end
        end
        break := False
        if switch = 0 then
            route := dbase.get_next_bus_route(route)
        else
            route := dbase.get_next_train_route(route)
        end
    end
end
end
switch := switch + 1
end
end
class STATION

```

```

creation
  make
feature {ROUTE, CONNECTION_FINDER}
  get_open : STRING is
    do
      Result := open.twin
    end
  get_acc : STRING is
    do
      Result := accessibility.twin
    end
  get_name : STRING is
    do
      Result := name.twin
    end
feature {NONE}
  accessibility : STRING
  open : STRING
  name : STRING
feature {ROUTE}
  do_save(file : TEXT_FILE_WRITE) is
    require
      file.is_connected
    do
      file.put_string(name + "_" + accessibility + "_" + open)
    end
  make(new_name, new_open, new_acc : STRING) is
    require
      new_name /= Void
      new_open /= Void
      new_acc /= Void
    do
      name := new_name.twin
      open := new_open.twin
      accessibility := new_acc.twin
    end
  show is
    do
      io.put_string(name)
    end
  match(targetname : STRING) : BOOLEAN is
    require
      targetname /= Void
    do
      Result := name.same_as(targetname)
    end
end

class FILE_DATABASE
inherit
  DATABASE
creation
  make
feature {HPTA_TRANSIT_INFO}
  do_save is
    local
      file_2_write : TEXT_FILE_WRITE
      i : INTEGER
    do
      create file_2_write.connect_to(file_name)
      if file_2_write.is_connected then
        from
          i := employees.lower
        until i > employees.upper
        loop
          employees.item(i).do_save(file_2_write)
          i := i + 1
        end
      end
    end
  end

```

```

        from
            i := bus_routes.lower
        until i > bus_routes.upper
        loop
            bus_routes.item(i).do_save(file_2_write)
            i := i + 1
        end
        from
            i := train_routes.lower
        until i > train_routes.upper
        loop
            train_routes.item(i).do_save(file_2_write)
            i := i + 1
        end
        file_2_write.disconnect
    else
        io.put_string("Update_database_failed!%N")
    end
end
add_route(new_route: ROUTE) is
do
    if is_bus(new_route.get_number) then
        bus_routes.add_last(new_route)
    elseif is_train(new_route.get_number) then
        train_routes.add_last(new_route)
    end
end
make is
do
    create employees.make
    create bus_routes.make
    create train_routes.make
    load
ensure
    employees /= Void
    bus_routes /= Void
    train_routes /= Void
end
delete_staff(staff: STAFF) is
do
    employees.remove(employees.index_of(staff) )
end
delete_route(route: ROUTE) is
do
    if is_bus(route.get_number) then
        bus_routes.remove(bus_routes.index_of(route) )
    elseif is_train(route.get_number) then
        train_routes.remove(train_routes.index_of(route) )
    end
end
add_staff(new_staff: STAFF) is
do
    employees.add_last(new_staff)
end
is_locked : BOOLEAN is
do
    Result := not employees.is_empty
end
feature {NONE}
    file_name : STRING is "sys.info.txt"
    is_bus(num: INTEGER) : BOOLEAN is
        do
            Result := num >= min_bus_route_num and num <= max_bus_route_num
        end
    employees : LINKED_LIST[STAFF]
    is_train(num: INTEGER) : BOOLEAN is
        do
            Result := num >= min_train_route_num and num <= max_train_route_num

```

```

end
bus_routes : LINKED_LIST [ROUTE]
load is
  local
    input_string : STRING
    text_file_read : TEXT_FILE_READ
    text_file_write : TEXT_FILE_WRITE
    split : ARRAY [STRING]
    new_staff : STAFF
    route : ROUTE
  do
    create text_file_read.connect_to(file_name)
    if text_file_read.is_connected then
      from
        text_file_read.read_line
      until text_file_read.end_of_input
      loop
        if text_file_read.last_string.upper = 1 then
          inspect
            text_file_read.last_string.first.to_upper
          when 'S' then
            text_file_read.read_line
            input_string := text_file_read.last_string.twin
            split := input_string.split
            create new_staff.make(split.first.to_integer, split.last)
            employees.add_last(new_staff)
          when 'B', 'T' then
            text_file_read.read_line
            input_string := text_file_read.last_string.twin
            split := input_string.split
            route := find_route(split.item(4).to_integer)
            if route = Void then
              create route.make(split.item(4).to_integer)
              route.add_station(split.first, split.item(2), split.item(3), split.last)
              if split.item(4).to_integer > max_train_route_num then
                bus_routes.add_last(route)
              else
                train_routes.add_last(route)
              end
            else
              route.add_station(split.first, split.item(2), split.item(3), split.last)
            end
          else
            end
          end
        end
        text_file_read.read_line
      end
      text_file_read.disconnect
    else
      create text_file_write.connect_to(file_name)
      if text_file_write.is_connected then
        text_file_write.disconnect
      end
    end
  end
train_routes : LINKED_LIST [ROUTE]
feature {HPTA_TRANSIT_INFO, CONNECTION_FINDER}
find_route(num: INTEGER) : ROUTE is
  local
    i : INTEGER
    route : ROUTE
  do
    if num > max_train_route_num then
      from
        i := bus_routes.lower
      until i > bus_routes.upper or else bus_routes.item(i).match(num)
      loop
        i := i + 1

```

```

        end
        if i <= bus_routes.upper then
            route := bus_routes.item(i)
        end
    else
        from
            i := train_routes.lower
        until i > train_routes.upper or else train_routes.item(i).match(num)
        loop
            i := i + 1
        end
        if i <= train_routes.upper then
            route := train_routes.item(i)
        end
    end
    Result := route
end

get_next_bus_route(route: ROUTE) : ROUTE is
    require
        bus_routes.index_of(route) > 0
    local
        next_route : ROUTE
    do
        if bus_routes.index_of(route) < bus_routes.upper then
            next_route := bus_routes.item(bus_routes.index_of(route) + 1)
        end
        Result := next_route
    end

get_first_bus_route : ROUTE is
    local
        route : ROUTE
    do
        if not bus_routes.is_empty then
            route := bus_routes.first
        end
        Result := route
    end

find_staff(num: INTEGER) : STAFF is
    local
        i : INTEGER
        staff : STAFF
    do
        from
            i := employees.lower
        until i > employees.upper or else employees.item(i).match(num)
        loop
            i := i + 1
        end
        if i <= employees.upper then
            staff := employees.item(i)
        end
        Result := staff
    end

get_next_train_route(route: ROUTE) : ROUTE is
    require
        train_routes.index_of(route) > 0
    local
        next_route : ROUTE
    do
        if train_routes.index_of(route) < train_routes.upper then
            next_route := train_routes.item(train_routes.index_of(route) + 1)
        end
        Result := next_route
    end

get_first_train_route : ROUTE is
    local
        route : ROUTE
    do

```

```

        if not train_routes.is_empty then
            route := train_routes.first
        end
        Result := route
    end
end

class ROUTE
creation
    make
feature {HPTA_TRANSIT_INFO}
    remove_station(name: STRING) is
        local
            i : INTEGER
        do
            from
                i := stops.lower
            until i > stops.upper or else stops.item(i).match(name)
            loop
                i := i + 1
            end
            if i <= stops.upper then
                stops.remove(i)
            end
        end
    show is
        local
            i : INTEGER
        do
            if number > 99 then
                io.put_string ("%NBus_route_No.")
            else
                io.put_string ("%NTrain_route_No.")
            end
            io.put_integer(number)
            io.put_string (":_")
            from
                i := stops.lower
            until i > stops.upper
            loop
                stops.item(i).show
                if i < stops.upper then
                    io.put_string (">")
                end
                i := i + 1
            end
            io.put_new_line
        end
    make(num: INTEGER) is
        require
            num > 9
            num < 1000
        do
            number := num
            create stops.make
        ensure
            number = num
            stops /= Void
        end
feature {DATABASE}
    do_save(file: TEXT_FILE_WRITE) is
        require
            file.is_connected
        local
            i : INTEGER
            tag : STRING
            last : STRING
        do

```

```

        if number > 99 then
            tag := "_b"
        else
            tag := "_t"
        end
        last := "_None"
        from
            i := stops.lower
        until i > stops.upper
        loop
            file.put_string(tag+"%N")
            stops.item(i).do_save(file)
            file.put_string("_" + number.to_string + "_" + last + "%N")
            last := stops.item(i).get_name.twin
            i := i + 1
        end
    end
feature {NONE}
    number : INTEGER
    stops : LINKED_LIST[STATION]
feature {CONNECTION_FINDER}
    get_first_station : STATION is
        local
            station : STATION
        do
            if stops.upper > 0 then
                station := stops.first
            end
            Result := station
        end
    get_next_station(station1 : STATION) : STATION is
        require
            station1 /= Void
        local
            station : STATION
        do
            if stops.index_of(station1) < stops.upper then
                station := stops.item(stops.index_of(station1) + 1)
            end
            Result := station
        end
feature {HPTA_TRANSIT_INFO, DATABASE}
    add_station(new_name, access, open_hour, last_stop : STRING) is
        local
            new_station : STATION
            i : INTEGER
            last : STRING
        do
            from
                i := stops.lower
            until i > stops.upper or else stops.item(i).match(new_name)
            loop
                i := i + 1
            end
            if i > stops.upper then
                create new_station.make(new_name, access, open_hour)
                last := last_stop.twin
                last.to_upper
                if last.same_as("NONE") then
                    stops.add_first(new_station)
                else
                    from
                        i := stops.lower
                    until i > stops.upper or else stops.item(i).match(last_stop)
                    loop
                        i := i + 1
                    end
                if i <= stops.upper then

```

```

        stops.add(new_station , i+1)
    else
        create new_station.make(last_stop , access , open_hour)
        stops.add_last(new_station)
        stops.add_last(new_station)
    end
end
end
end
feature {ANY}
get_number : INTEGER is
do
    Result := number
end
match(num: INTEGER) : BOOLEAN is
do
    Result := num = number
end
invariant
    number > 9
    number < 1000
end

class HPTA_TRANSIT_INFO
inherit
    SYSTEM
creation
    make
feature {NONE}
delete_menu is
do
    execute_command_line("cls")
    io.put_string(" [
-----
-----Welcome_to_HPTA
-----
-----S_Delete_a_station
-----E_Delete_a_staff
-----R_Delete_a_route
-----G_Go_back
-----Enter_menu_choice:
-----]")
end
db : DATABASE
update_menu is
do
    execute_command_line("cls")
    io.put_string(" [
-----
-----Welcome_to_HPTA
-----
-----A_Add
-----D_Delete
-----G_Go_back
-----Enter_menu_choice:
-----]")
end
menu is
do
    execute_command_line("cls")
    io.put_string(" [
-----
-----Welcome_to_HPTA
-----
-----U_Update_System_Information
-----I_Inquire_about_Transit_Information
-----Q_Quit
-----Enter_menu_choice:

```

```

.....]”)
    end
    set_finder (new_finder: CONNECTION.FINDER) is
    require
        new_finder /= Void
    do
        finder := new_finder
    ensure
        finder = new_finder
    end
    inquire_menu is
    do
        execute_command_line (“cls”)
        io.put_string (“[
.....*****
.....Welcome_to_HPPTA
.....*****
.....F_Find_a_possible_connection
.....S_Show_a_route
.....B_Browse_all_routes
.....G_Go_back
.....Enter_menu_choice:
.....]”)
    end
    do_delete is
    local
        is_end : BOOLEAN
        num : INTEGER
        staff : STAFF
        route : ROUTE
        input : STRING
        name : STRING
    do
        from
        until is_end
        loop
            delete_menu
            io.read_line
            input := io.last_string.twin
            io.put_new_line
            if not input.is_empty then
                inspect
                    input.first.to_upper
                when ‘G’ then
                    is_end := True
                when ‘S’ then
                    io.put_string (“%NEnter_route_number:_”)
                    io.read_line
                    num := io.last_string.to_integer
                    io.put_string (“%NEnter_station_name:_”)
                    io.read_line
                    name := io.last_string.twin
                    route := db.find_route(num)
                    if route /= Void then
                        route.remove_station(name)
                    else
                        io.put_string (“%NNo_such_a_station%N”)
                        io.read_line
                    end
                when ‘R’ then
                    io.put_string (“%NEnter_route_number:_”)
                    io.read_line
                    num := io.last_string.to_integer
                    route := db.find_route(num)
                    if route /= Void then
                        db.delete_route(route)
                    else
                        io.put_string (“%NNo_such_a_station%N”)

```

```

        io.read_line
      end
    when 'E' then
      io.put_string("%NEnter_ID: ")
      io.read_line
      num := io.last_string.to_integer
      staff := db.find_staff(num)
      if staff /= Void then
        db.delete_staff(staff)
      else
        io.put_string ("%NNo_such_a_staff%N")
        io.read_line
      end
    else
      end
    end
  end
end
finder : CONNECTION_FINDER
do-update is
  local
    id : INTEGER
    passed : BOOLEAN
    is_end : BOOLEAN
    password : STRING
    input : STRING
    staff : STAFF
  do
    io.read_line
    if db.is_locked then
      io.put_string ("%NEnter_employee_ID: ")
      io.read_line
      id := io.last_string.to_integer
      staff := db.find_staff(id)
      if staff /= Void then
        io.put_string ("%NEnter_password: ")
        io.read_line
        password := io.last_string.twin
        passed := staff.login(password)
      end
    else
      io.put_string(" [
      ~~~~~The_list_of_authorized_staff_is_not_empty,
      ~~~~~so_please_set_authorization_as_soon_as_possible...
      ~~~~~] ")
      passed := True
      io.read_line
    end
    if passed then
      from
      until is_end
      loop
        update_menu
        io.read_line
        input := io.last_string.twin
        io.put_new_line
        if not input.is_empty then
          inspect
            input.first.to_upper
          when 'A' then
            do_add
          when 'D' then
            do_delete
          when 'G' then
            is_end := True
          else
            end
        end
      end
    end
  end
end

```

```

        end
        db.do_save
    else
        io.put_string ("%NLogin_failed!%N")
        io.read_line
    end
end
run is
do
    from
    until io.last_character.to_upper = 'Q'
    loop
        menu
        io.read_character
        io.put_new_line
        inspect
            io.last_character.to_upper
        when 'U' then
            do_update
        when 'I' then
            do_inquire
        else
            end
        end
    end
end
add_menu is
do
    execute_command_line ("cls")
    io.put_string (" [
-----*****
-----Welcome_to_HPTA
-----*****
-----S_Add_a_station
-----E_Add_a_staff
-----G_Go_back
-----Enter_menu_choice:
-----] ")
    end
do_add is
    local
        employee : STAFF
        id : INTEGER
        input : STRING
        name : STRING
        password : STRING
        open : STRING
        access : STRING
        last : STRING
        is_end : BOOLEAN
        route : ROUTE
    do
        from
        until is_end
        loop
            add_menu
            io.read_line
            input := io.last_string.twin
            io.put_new_line
            if not input.is_empty then
                inspect
                    input.first.to_upper
                when 'G' then
                    is_end := True
                when 'S' then
                    io.put_string ("%NEnter_station_name:_")
                    io.read_line
                    name := io.last_string.twin
                    io.put_string ("%NEnter_open_hour:_")

```



```

        end
      from
        route := db.get_first_train_route
      until route = Void
      loop
        route.show
        route := db.get_next_train_route(route)
      end
      io.put_string ("%N%NStrike_any_key_to_continue...")
      io.read_line
      when 'F' then
        io.put_string ("%NEnter_the_station_name_of_your_start:_")
        io.read_line
        start := io.last_string.twin
        io.put_string ("%NEnter_the_station_name_of_your_destination:_")
        io.read_line
        dest := io.last_string.twin
        io.put_string ("%NEnter_your_desire_time(in_minutes):_")
        io.read_line
        time := io.last_string.to_integer
        io.put_string (finder.get_connection(db, start, dest, time) )
        io.put_string ("%N%NStrike_any_key_to_continue...")
        io.read_line
      when 'S' then
        io.put_string ("Input_the_route_number_(10_-_999):_")
        io.read_line
        num := io.last_string.to_integer
        route := db.find_route(num)
        if route /= Void then
          route.show
        else
          io.put_string ("Sorry_there_is_no_such_a_route")
        end
        io.put_string ("%N%NStrike_any_key_to_continue...")
        io.read_line
      when 'G' then
        is_end := True
      else
        end
      end
    end
  end
feature {ANY}
  make is
    local
      prime_finder : PRIME_FINDER
      file_database : FILE_DATABASE
    do
      create file_database.make
      set_database(file_database)
      create prime_finder.make
      set_finder(prime_finder)
      run
    end
  end
end

```

Appendix D

Reference Manual of Spark

D.1 Code Block Tag

Code Block Tag is used to identify the class, to which this block belongs. So there is nothing need to do for the class block, but for class member block, including invariant block, ones must put code block tag, class name followed by three dots, at the beginning.

D.2 Graphic Notation Setting

All the tags listed as following should be put in setting line, which is right behind graphic including command.

- @VERTICAL: if set, the diagram will be drawn vertically, otherwise horizontally.
- @HEAD: if set, the class diagram will be shown with class name nodes only.
- @BRIEF: if set, the class diagram will be shown without parameters and types.
- @CONCISE: if set, the class diagram will hide all the information about the method's parameters of the involved class.
- @METHOD: if set, all class methods only will be shown in this diagram.

- @ATTRIBUTE: if set, only attributes of class can be saw in the diagram.
- @ACTION” if set, only actions of class can be saw in the diagram.

D.3 Program Code Quotation

“CODE LIST BEGIN” and “CODE LIST END” are the specific tags used to include continuous program code into the source file. This tags can be put anywhere in the source file as comments. Front ends will insert the parsed code between them, if they find them.

Appendix E

Document Structure of AsciiDoc

An AsciiDoc document consists of a series of block elements. Almost any combination of zero or more elements constitutes a valid AsciiDoc document: documents can range from a single sentence to a multi-part book. In the following table of AsciiDoc document structure, parentheses ‘(’ and ‘)’ indicate grouping when needed, square brackets ‘[’ and ‘]’ enclose optional items, curly parentheses ‘{’ and ‘}’ show the (zero or more) repeatable items, and vertical bars ‘|’ separate alternatives.

```
Document ::= [ Header ] [ Preamble ] { Section }
  Header ::= Title [ AuthorLine [ RevisionLine ] ]
  AuthorLine ::= FirstName [ [ MiddleName ] LastName ] [ EmailAddress ]
  RevisionLine ::= [ Revision ] Date
  Preamble ::= SectionBody
  Section ::= Title [ SectionBody ] { Section }
  SectionBody ::= ( ( [ BlockTitle ] Block ) | BlockMacro ) { ( [ BlockTitle ] Block ) | BlockMacro }
  Block ::= Paragraph | DelimitedBlock | List | Table
  List ::= BulletedList | NumberedList | LabeledList | CalloutList
  BulletedList ::= ListItem { ListItem }
  NumberedList ::= ListItem { ListItem }
  CalloutList ::= ListItem { ListItem }
  LabeledList ::= ItemLabel { ItemLabel } ListItem { ItemLabel { ItemLabel } ListItem }
  ListItem ::= ItemText { List | ListParagraph | ListContinuation }
  Table ::= Ruler [ TableHeader ] TableBody [ TableFooter ]
  TableHeader ::= TableRow { TableRow } TableUnderline
  TableFooter ::= TableRow { TableRow } TableUnderline
  TableBody ::= TableRow { TableRow } TableUnderline
```

TableRow ::= TableData { TableData }

Table E.1: The block structure of AsciiDoc.

Appendix F

Syntax of Dot

The following is an abstract grammar for the **dot** language. Terminals are shown in bold font and nonterminals in italics. Literal characters are given in single quotes. Parentheses '(' and ')' indicate grouping when needed. Square brackets '[' and ']' enclose optional items. Curly parentheses '{' and '}' show the (zero or more) repeatable items. Vertical bars '|' separate alternatives.

$$\begin{aligned} \textit{graph} &::= [\textbf{strict}] (\textbf{digraph} | \textbf{graph}) \textit{id} \{ ' ' \textit{stmt-list} ' ' \} \\ \textit{id} &::= \textit{letter} \{ \textit{letter} | \textit{digital} | _ \} \\ \textit{letter} &::= \mathbf{a} | \mathbf{b} | \mathbf{c} | \mathbf{d} | \mathbf{e} | \mathbf{f} | \mathbf{g} | \mathbf{h} | \mathbf{i} | \mathbf{j} | \mathbf{k} | \mathbf{l} | \mathbf{m} | \mathbf{n} | \mathbf{o} | \mathbf{p} | \mathbf{q} | \mathbf{r} | \mathbf{s} | \mathbf{t} | \\ &\quad \mathbf{u} | \mathbf{v} | \mathbf{w} | \mathbf{x} | \mathbf{y} | \mathbf{z} | \mathbf{A} | \mathbf{B} | \mathbf{C} | \mathbf{D} | \mathbf{E} | \mathbf{F} | \mathbf{G} | \mathbf{H} | \mathbf{I} | \mathbf{J} | \mathbf{D} | \mathbf{L} | \\ &\quad \mathbf{M} | \mathbf{N} | \mathbf{O} | \mathbf{P} | \mathbf{Q} | \mathbf{R} | \mathbf{S} | \mathbf{T} | \mathbf{U} | \mathbf{V} | \mathbf{W} | \mathbf{X} | \mathbf{Y} | \mathbf{Z} \\ \textit{digital} &::= \mathbf{1} | \mathbf{2} | \mathbf{3} | \mathbf{4} | \mathbf{5} | \mathbf{6} | \mathbf{7} | \mathbf{8} | \mathbf{9} | \mathbf{0} \\ \textit{stmt-list} &::= [\textit{stmt} [';'] [\textit{stmt-list}] \\ \textit{stmt} &::= \textit{attr-stmt} | \textit{node-stmt} | \textit{edge-stmt} | \textit{subgraph} | \textit{id} '=' \textit{id} \\ \textit{attr-stmt} &::= (\textbf{graph} | \textbf{node} | \textbf{edge}) \textit{attr-list} \\ \textit{attr-list} &::= '[' [\textit{a-list}] ']' [\textit{attr-list}] \\ \textit{a-list} &::= \textit{id} '=' \textit{id} [','] [\textit{a-list}] \\ \textit{node-stmt} &::= \textit{node-id} [\textit{attr-list}] \\ \textit{node-id} &::= \textit{id} [\textit{port}] \\ \textit{port} &::= \textit{port-location} [\textit{port-angle}] | \textit{port-angle} [\textit{port-location}] \\ \textit{port-location} &::= ':' \textit{id} | ':' '(' \textit{id} ',' \textit{id} ')' \\ \textit{port-angle} &::= '@' \textit{id} \end{aligned}$$

```

edge-stmt ::= ( node-id | subgraph ) edgeRHS [ attr-list ]
edgeRHS ::= edgeop ( node-id | subgraph ) [ edgeRHS ]
subgraph ::= [ subgraph id ] '{' stmt-list '}' | subgraph id
edgeop ::= -> | --

```

Table F.1: Abstract grammar for the dot language

The language supports C++ style comments: `/* */` and `//`.

Semicolons aid readability but are not required except in the rare case that a named subgraph with no body immediate precedes an anonymous subgraph, because under precedence rules this sequence is parsed as a subgraph with a heading and a body.

Complex attribute values may contain characters, such as commas and white space, which are used in parsing the dot language. To avoid getting a parsing error, such values need to be enclosed in double quotes.

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