How many double squares can a string contain?

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Outline

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Motivation and background

We are dealing with finite strings over finite alphabets. There is no particular requirement about the order of the alphabet.

What is the maximum number of distinct squares problem?

We are counting types of squares rather than their occurrences.

`aababaaba` has 6 occurrences of squares, but only 4 distinct squares, `aa`, `aabaab`, `abaaba`, and `baabaa`.

How many double squares can a string contain?
A trivial bound: the number of all occurrences of primitively rooted squares in a string of length \( n \) is bounded by \( O(n \log n) \) (Crochemore 1978) and the number of distinct non-primitively rooted squares is \( O(n) \) (Kubica et al. 2013)

Could it be \( O(n) \)? And if so, what would be the constant?

Why this is not simple? In a string of length \( n \), \( O(\log n) \) squares can start at the same position!
It is easy to compute it for short strings, so why induction cannot be used?

\[ \text{a|a|b|a|a|b} + \text{a|a|b|a|a|b} \]

\[ \text{a|a|b|a|a|b|a|a|b|a|a|b|a|a|b} \]

Concatenation does both “destroys” existing types through multiple-occurrences and “creates” new types. Of course, same holds true for the reverse process - partitioning of strings.
Theorem (Fraenkel-Simpson, 1998)

There are at most $2n$ distinct squares in a string of length $n$.

Count only the rightmost occurrences. Fraenkel-Simpson showed that if there are three rightmost squares $uu$, $vv$, and $ww$ starting at the same position so that $|u| < |v| < |w|$, then $ww$ contains a farther copy of $uu$, based on Crochemore-Rytter (1995) Lemma showing that in such a case, $|w| \geq |u| + |v|$.
**Fraenkel-Simpson** hypothesized that the number of distinct squares should be bounded by \( n \), i.e.

\[
\sigma(n) \leq n
\]

where \( \sigma(n) = \max \{ s(x) : x \text{ is a string of length } n \} \).

**Fraenkel-Simpson** gave an infinite sequence of strings \( \{x_n\}_{n=1}^\infty \) so that \( |x_n| \nearrow \infty \) and

\[
\frac{s(x_n)}{|x_n|} \nearrow 1
\]

where \( s(x) = \text{number of distinct squares in } x \).
In 2005 Ilie provided a simpler proof of Fraenkel-Simpson’s Theorem and in 2007 presented an asymptotic upper bound of $2n - \theta(\log n)$.

In 2011 Deza-F. proposed a $d$-step approach to the problem and conjectured that $\sigma_d(n) \leq n - d$, where $\sigma_d(n) = \max \{ s(x) : x \text{ is a string of length } n \text{ with } d \text{ distinct symbols} \}$.
Basic notions and tools

**Definition**

*non-trivial power* of a string $x$ is a concatenation of $m$ copies of $x$ denotes as $x^m$; $x^2$ is a *square*, $x^3$ a *cube*.

A string $x$ is *primitive* if $x \neq y^n$ for any $y$ and any $n \geq 2$.

*primitive root* of $x$ is the smallest primitive $y$ so that $x = y^n$.

$x$ and $y$ are *conjugates* if $x = uv$ and $y = vu$ for some $u, v$. 
Lemma (Synchronization principle)

Given a primitive string $x$, a proper suffix $y$ of $x$, a proper prefix $z$ of $x$, and $m \geq 0$, there are exactly $m$ occurrences of $x$ in $yx^mz$.

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Lemma (Common factor lemma)

For any strings $x$ and $y$, if a non-trivial power of $x$ and a non-trivial power of $y$ have a common factor of length $|x| + |y|$, then the primitive roots of $x$ and $y$ are conjugates. In particular, if $x$ and $y$ are primitive, then $x$ and $y$ are conjugates.
**Double squares**

- **Fraenkel-Simpson**: only two rightmost squares can start at the same position. Thus, only one rightmost square or two rightmost squares may start at any position.

- **Lam (2009 – unpublished)** tried bounding the number of **double squares** and hence bound the number of distinct squares. His approach is based on a taxonomy of all possible configurations of two double squares yielding a bound of $\frac{94}{48} n \approx 1.98n$. 
A configuration of two squares

\[
\begin{array}{|c|c|}
\hline
U & U \\ \hline
u & u \\ \hline
\end{array}
\]

has been investigated in many different contexts:

- *Smyth et. al.*: with intention to find a position for amortization argument for runs conjecture.
- in computational framework by *Deza-F.-Jiang*: such configurations are used in *Liu’s Ph.D. thesis* to speed up computation of $\sigma_d(n)$.
- *Lam*: two rightmost squares have a unique structure.
Lemma

Let $uu$ and $UU$ be two squares in a string $x$ starting at the same position with $|u| < |U|$ such that either

(b) both $uu$ and $UU$ are rightmost occurrences, or

(a) $|U| < |uu|$ and either $uu$ or $UU$ is primitively rooted.

Then $|u| < |U| < |uu| < |UU|$ and there is a unique primitive string $u_1$, a unique proper prefix $u_2$ of $u_1$, and unique integers $e_1$ and $e_2$ satisfying $1 \leq e_2 \leq e_1$ such that $u = u_1^{e_1}u_2$ and $U = u_1^{e_1}u_2u_1^{e_2}$; i.e. $uu$ and $UU$ form a double square.
Thus, only strings of length at least 10 may contain a double square: $|UU| = 2((u(1)+u(2))|u_1|+|u_2|) \geq 2((1+1)2 + 1) = 10$. 
Cyclic shift (rotation) to the right is controlled by

\[ lcp(u_1, \bar{u}_1) \]

while cyclic shift to the left is controlled by

\[ lcs(u_1, \bar{u}_1) \]

\( lcp \) = largest common prefix

\( lcs \) = largest common suffix
Motivation and background  Basic notions and tools  Double squares  Inversion factors  Rightmost double squares  An upper bound for the number of double squares  Conclusion  References

How many double squares can a string contain?  Supeléc, February 2014

\[ u_1 = aaabaa, \ u_2 = aaab, \ \bar{u}_2 = aa, \ u(1) = u(2) = 2 \]
Motivation and background

Basic notions and tools

Double squares

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Conclusion

References

\[ u_1 = aaabaa, \quad u_2 = aaab, \quad \bar{u}_2 = aa, \quad u(1) = 2, \quad \text{and} \quad u(2) = 1. \]

How many double squares can a string contain?

Supeléc, February 2014
Definition

For a double square $\mathcal{U}, \overline{v}vv\overline{v}$ where $|\overline{v}| = |\overline{u}_2|$ and $|v| = |u_2|$ is an inversion factor

\[
\mathcal{U} = u_1 u^{(1)} u_2 u_1 u^{(2)} + u^{(1)} u_2 u_1 u^{(2)} = \\
N_1 \quad \text{and} \quad N_2 \\
\text{natural inversion factors}
\]
A cyclic shift of an inversion factor is an inversion factor, also controlled by $lcp(u_1, \bar{u}_1)$ and $lcs(u_1, \bar{u}_1)$. 

How many double squares can a string contain? Supeléc, February 2014
All inversion factors are cyclic shifts of the natural ones:

**Lemma (Inversion factor lemma)**

*Given a double square $U$, there is an inversion factor of $U$ within the string $UU$ starting at position $i$ $\iff i \in [L_1, R_1] \cup [L_2, R_2]$.***
Inversion factor lemma for distinct squares

**Theorem (Fraenkel-Simpson, Ilie)**

*At most two rightmost squares can start at the same position.*

Let us assume that 3 rightmost squares $uu$, $UU$, $vv$ start at the same position.

By item (c) of Inversion factor lemma, $uu$ and $UU$ form a double square $\mathcal{U}: u = u_1^{\mathcal{U}(1)} u_2$ and $U = u_1^{\mathcal{U}(1)} u_2 u_1^{\mathcal{U}(2)}$.

Since the first $v$ contains an inversion factor, the second $v$ must also contain an inversion factor.

*Cont. on the next slide*
If the inversion factor in the second $v$ were from $[L_2, R_2]$, then $|v| = |U|$, a contradiction. Hence $v$ must not contain an inversion factor from $[L_2, R_2]$ and so $u_1^U(1)u_2u_1^U(1+U(2)-1)u_2$ must be a prefix of $v$. Therefore $vv$ contains another copy of $u_1^U(1)u_2u_1^U(1)u_2 = uu$, a contradiction.
Fundamental Lemma:

**Lemma**

Let $x$ be a string starting with a double square $U$. Let $V$ be a double square with $s(U) < s(V)$, then either

**a)** $s(V) < R_1(U)$, in which case either

(a1) $V$ is an $\alpha$-mate of $U$ (cyclic shift), or
(a2) $V$ is a $\beta$-mate of $U$ (cyclic shift of $U$ to $V$), or
(a3) $V$ is a $\gamma$-mate of $U$ (cyclic shift of $U$ to $v$), or
(a4) $V$ is a $\delta$-mate of $U$ (big tail),

or

**b)** $R_1(U) \leq s(V)$, then

(b1) $V$ is an $\varepsilon$-mate of $U$ (big gap).
**α-mate (cyclic shift):**

\[
\alpha - \text{mate (cyclic shift)}:
\]

\[
\text{[ ] [ ] ( ) ( ) ( )}
\]

\[
\text{aaabaaaaabaaaaabaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaabaaaaabaaaaabaaaaabaaaaabaaaaabaa...}
\]

\[
\text{[ ] [ ] ( ) ( ) ( )}
\]

\[
\text{aaba...aaba...aaba...aaba...aaba...aaba...aaba...}
\]
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How many double squares can a string contain?

**γ-mate (cyclic shift of $U$ to $v$)**

```
\begin{align*}
R_1 & \quad [ \quad ] \quad [ \quad ) \quad ( \quad ] \quad ) \\
\gamma\text{-mate} & \quad [ \quad ] \quad [ \quad ) \quad ( \quad ] \quad ) \\
\text{aabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaab} & \quad [ \quad ] \quad [ \quad ) \quad ( \quad ] \quad ) \\
\text{aabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaabaab} & \quad [ \quad ] \quad [ \quad ) \quad ( \quad ] \quad ) \\
\end{align*}
```
δ-mate (big tail)

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**ε-mate (big gap)**

How many double squares can a string contain?
An upper bound for the number of double squares

We show by induction a bound $\delta(x) \leq \frac{5}{6}|x| - \frac{1}{3}|u|$, where $uu$ is the shorter square of the leftmost double square of $x$.

<table>
<thead>
<tr>
<th>$u$</th>
<th>$u$</th>
<th>$T$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$G$</td>
<td>$v$</td>
<td>$v$</td>
</tr>
</tbody>
</table>

The fundamental lemma basically says that either the gap $G(U, V)$ is "big" or the tail $T(U, V)$ is "big" (for $\delta$-mate and $\varepsilon$-mate), or it is case of $\alpha$-mate, $\beta$-mate, or $\gamma$-mate.
Lemma (Gap-Tail lemma)

\[ \delta(x') \leq \frac{5}{6} |x'| - \frac{1}{3} |v| \] \text{ implies }

\[ \delta(x) \leq \frac{5}{6} |x| - \frac{1}{3} |u| + d - \frac{1}{2} |G(U, V)| - \frac{1}{3} |T(U, V)| \]
We deal with $\alpha$-mates, $\beta$-mates, and $\gamma$-mates separately.

It is possible as they form families, either a pure $\alpha$-family, or $\alpha+\beta$-family, or $\alpha+\beta+\gamma$-family.
\textit{U}-family consists only of $\alpha$-mates

Illustration of $\alpha$-family with $U(1) = U(2)$

\[
\begin{array}{c}
\text{aaaabaaaaaabaaabaabaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaa}
\end{array}
\]
Illustration of $\alpha$-family with $\mathcal{U}(1) > \mathcal{U}(2)$

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It is easy to estimate the size of $\alpha$-family, as it is controlled by $lcp(u_1, \overline{u}_1)$ and $lcp(y, u_2)$ where $y$ is $x$ without $UU$: the size $\leq |u_1|$.

- Either there are no other double squares, and then it can be shown directly that the bound holds, or

- There is a $\mathcal{V}$ underneath, and we can use induction using the Gap-Tail lemma. $\mathcal{V}$ must be either $\gamma$-mate, or $\delta$-mate, or $\varepsilon$-mate, and the Gap-Tail lemma can be applied to propagate the bound.
**U-family consists of α-mates and β-mates**

Illustration of $\alpha+\beta$-family

```
aaabaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaaaabaaa
It is more complicated to estimate the size of a $\alpha+\beta$-family:

$$\left\lfloor \frac{U(1)-U(2)}{2} \right\rfloor |u_1| \quad \text{if } U(2) = 1$$

$$\frac{U(1)-U(2)}{2} |u_1| \quad \text{if } U(2) > 1$$

- Either there are no other double squares, and then it can be shown directly that the bound holds, or
- There is a $V$ underneath, and we can use induction using the Gap-Tail lemma. $V$ must be either $\delta$-mate, or $\epsilon$-mate, and the Gap-Tail lemma can be applied to propagate the bound. *(Special care needed for $\epsilon$-mate case and super-$\epsilon$-mate must be put in play!)*
Illustration of $\alpha+\beta+\gamma$-family

\[ R_1 \]

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It is quite complex to estimate the size of a $\alpha+\beta+\gamma$-family:

$$ \leq \frac{2}{3}(u(1) + 1)|u_1| $$

- Either there are no other double squares, and then it can be shown directly that the bound holds, or

- There is a $\nu$ underneath, and we can use induction using the Gap-Tail lemma. $\nu$ must be either $\delta$-mate, or $\varepsilon$-mate, and the Gap-Tail lemma can be applied to propagate the bound.
Main theorems

Theorem

The number of double squares in a string of length $n$ is bounded by $\lfloor \frac{5n}{6} \rfloor$.

Corollary

The number of distinct squares in a string of length $n$ is bounded by $\lfloor \frac{11n}{6} \rfloor$. 
We presented a universal upper bound of $\frac{11n}{6}$ for the maximum number of distinct squares in a string of length $n$.

A bound of $\frac{5n}{6}$ for the maximum number of double squares.

It improves the universal bound of $2n$ by Fraenkel-Simpson.

It improves the asymptotic bound of $2n - \Theta(n)$ by Ilie.

The combinatorics of double squares is interesting on its own and possibly can be used for some other problems.
THANK YOU
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How many squares can a string contain?  

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