

I/O Exercise 1

In the standard library module you find:

```
data ExitCode = ExitSuccess | ExitFailure Int
  deriving (Eq, Ord, Read, Show)
```

```
getArgs :: IO [String]
system  :: String → IO ExitCode
```

Write a program that accepts *name* as argument and then calls LaTeX and BibTeX on *name.ltx* until the document can be expected to stabilise.

Defining Functions Over Lists by Pattern Matching

Some functions taking lists as arguments can be defined directly via **pattern matching**:

```
null      :: [a] → Bool
null []   = True
null (x : xs) = False
```

```
head      :: [a] → a
head (x : xs) = x
head []   = error "head: empty list!"
```

```
tail      :: [a] → [a]
tail (x : xs) = xs
tail []   = error "tail: empty list!"
```

(`head` and `tail` are **partial functions** — both are undefined on the empty list.)

Defining Functions Over Lists by Structural Induction

Many functions taking lists as arguments can be defined via **structural induction**:

```
length      :: [a] → Int
length []   = 0
length (x : xs) = 1 + length xs

concat      :: [[a]] → [a]
concat []   = []
concat (xs : xss) = xs ++ concat xss
```

```
(+)        :: [a] → [a] → [a]
[]         ++ ys = ys
(x : xs) ++ ys = x : (xs ++ ys)
```

```
x 'elem' [] = False
x 'elem' (y : ys)
  = x == y || x 'elem' ys
```

(All these functions are in the standard prelude.)

Guarded Definitions

```
sign x | x > 0 = 1
       | x == 0 = 0
       | x < 0 = -1

choose :: Ord a ⇒ (a, b) → (a, b) → b
choose (x, v) (y, w)
  | x > y = v
  | x < y = w
  | otherwise = error "I cannot decide!"
```

If no guard succeeds, the next pattern is tried:

```
take_while p (x : xs) | p x = x : take_while p xs
take_while p xs      = []
```

```
take_while (< 5) [1, 2, 3]
= take_while (< 5) (1 : 2 : 3 : [])
= 1 : take_while (< 5) (2 : 3 : [])
= 1 : 2 : take_while (< 5) (3 : [])
= 1 : 2 : 3 : take_while (< 5) []
= 1 : 2 : 3 : []
= [1, 2, 3]
```

Guarded Definitions — Fall-Through

If no guard succeeds, the next pattern is tried:

```
take_while p (x : xs) | p x = x : take_while p xs
take_while p xs           = []
```

```
take_while (< 5) [1, 2, 3, 2, 3, 4, 3, 4, 5, 4, 3, 4, 5, 6]
= take_while (< 5) (1 : 2 : 3 : 2 : 3 : 4 : 3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : take_while (< 5) (2 : 3 : 2 : 3 : 4 : 3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : take_while (< 5) (3 : 2 : 3 : 4 : 3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : take_while (< 5) (2 : 3 : 4 : 3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : 2 : take_while (< 5) (3 : 4 : 3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : 2 : 3 : take_while (< 5) (4 : 3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : 2 : 3 : 4 : take_while (< 5) (3 : 4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : 2 : 3 : 4 : 3 : take_while (< 5) (4 : 5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : 2 : 3 : 4 : 3 : 4 : take_while (< 5) (5 : 4 : 3 : 4 : 5 : 6 : [])
= 1 : 2 : 3 : 2 : 3 : 4 : 3 : 4 : []
= [1, 2, 3, 2, 3, 4, 3, 4]
```

case Expressions

```
sign x = case compare x 0 of
  GT -> 1
  EQ -> 0
  LT -> -1
```

The prelude datatype *Ordering* has three elements and is used mostly as result type of the prelude function *compare*:

```
data Ordering = LT | EQ | GT
```

```
compare :: Ord a => a -> a -> Ordering
```

Another example:

```
choose (x, v) (y, w) = case compare x y of
  GT -> v
  LT -> w
  EQ -> error "I cannot decide!"
```

if ... then ... else ... and case Expressions

The type *Bool* can be considered as a two-element enumeration type:

```
data Bool = False | True
```

Conditional expressions are “syntactic sugar” for **case** expressions over *Bool*:

<pre>if condition then expr1 else expr2</pre>	≡	<pre>case condition of True -> expr1 False -> expr2</pre>
---	---	--

Two ways of defining functions:

Pattern Matching

```
not True = False
not False = True
```

case

```
not b = case b of
  True -> False
  False -> True
```

case Expressions are “Anonymous” Pattern Matching

```
commaWords :: [String] -> String
commaWords [] = []
commaWords (x : xs) = x ++ case xs of
  [] -> []
  _ -> ", " : commaWords xs
```

Every use of a case expression can be transformed into the use of an auxiliary function defined by pattern matching:

```
commaWords :: [String] -> String
commaWords [] = []
commaWords (x : xs) = x ++ commaWordsAux xs
```

```
commaWordsAux [] = []
commaWordsAux xs = ", " : commaWords xs
```

where Clauses

If an auxiliary definition is used only locally, it should be inside a **local definition**, e.g.:

```
commaWords :: [String] → String
commaWords [] = []
commaWords (x : xs) = x ++ commaWordsAux xs
where
  commaWordsAux [] = []
  commaWordsAux xs = ", " : commaWords xs
```

where clauses are visible **only** within their enclosing clause, here “`commaWords (x : xs) = ...`”

where clauses are visible within all guards:

```
f x y | y > z = ...
      | y == z = ...
      | y < z = ...
  where z = x * x
```

let Expressions

Local definitions can also be part of expressions:

```
f k n = let m = k `mod` n
        in if m == 0
           then n
           else f n m

h x y = let x2 = x * x
        y2 = y * y
        in sqrt (x2 + y2)
```

Definitions can use **pattern bindings**:

```
g k n = let (d,m) = divMod k n
          in if d == 0
             then [m]
             else g d n ++ [m]
```

Guards, `let` and `where` bindings, and `case` cases all are **layout sensitive!**

let or where?

- `let bindings` in *expression* is an **expression**
- *fname patterns guardedRHSs* `where bindings` is a clause that is part of a **definition**
- (where clauses can also modify *case* cases)

Frequently, the choice between `let` and `where` is a matter of *style*:

- `where` clauses result in a top-down presentation
- `let` expressions lend themselves also to bottom-up presentations

Some Prelude Functions — Elementary List Access

```
head :: [a] -> a
head (x:_) = x

last :: [a] -> a
last [x] = x
last (_:xs) = last xs

tail :: [a] -> [a]
tail (_:xs) = xs

init :: [a] -> [a]
init [x] = []
init (x:xs) = x : init xs

null :: [a] -> Bool
null [] = True
null (_:_) = False
```

Some Prelude Functions — List Indexing

```
length      :: [a] -> Int
length     = foldl' (\n _ -> n + 1) 0

(!!)       :: [b] -> Int -> b
(x:_ ) !! 0 = x
(_:xs) !! n | n>0 = xs !! (n-1)
(_:_ ) !! _ = error "PreludeList.!!: negative index"
[]        !! _ = error "PreludeList.!!: index too large"
```

Some Prelude Functions — Positional List Splitting

```
take       :: Int -> [a] -> [a]
take 0 _   = []
take _ []  = []
take n (x:xs) | n>0 = x : take (n-1) xs
take _ _   = error "take: negative argument"

drop       :: Int -> [a] -> [a]
drop 0 xs  = xs
drop _ []  = []
drop n (_:xs) | n>0 = drop (n-1) xs
drop _ _   = error "drop: negative argument"

splitAt    :: Int -> [a] -> ([a], [a])
splitAt 0 xs = ([],xs)
splitAt _ [] = ([],[])
splitAt n (x:xs) | n>0 = (x:xs',xs")
              where (xs',xs") = splitAt (n-1) xs
splitAt _ _ = error "splitAt: negative argument"
```

Some Prelude Functions — Concatenation, Iteration

```
(++) :: [a] -> [a] -> [a]
[]    ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)

concat :: [[a]] -> [a]
concat = foldr (++) []

iterate :: (a -> a) -> a -> [a]
iterate f x = x : iterate f (f x)

repeat :: a -> [a]
repeat x = xs where xs = x:xs
{- repeat x = x : repeat x -} -- for understanding

replicate :: Int -> a -> [a]
replicate n x = take n (repeat x)

cycle :: [a] -> [a]
cycle xs = xs' where xs' = xs ++ xs'
```

Separation of Concerns: Generation and Consumption

```
replicate 3 '!'
= take 3 (repeat '!') -- replicate
= take 3 ('!' : repeat '!') -- repeat
= '!' : take (3 - 1) (repeat '!') -- take (iii)
= '!' : take 2 (repeat '!') -- subtraction
= '!' : take 2 ('!' : repeat '!') -- repeat
= '!' : '!' : take (2 - 1) (repeat '!') -- take (iii)
= '!' : '!' : take 1 (repeat '!') -- subtraction
= '!' : '!' : take 1 ('!' : repeat '!') -- repeat
= '!' : '!' : '!' : take (1 - 1) (repeat '!') -- take (iii)
= '!' : '!' : '!' : take 0 (repeat '!') -- subtraction
= '!' : '!' : '!' : [] -- take (i)
= "!!!"
```

Exercise: Splitting with Predicates

- $takeWhile :: (a \rightarrow Bool) \rightarrow [a] \rightarrow [a]$
 $takeWhile$, applied to a predicate p and a list xs , returns the longest prefix (possibly empty) of xs of elements that satisfy p .
- $dropWhile :: (a \rightarrow Bool) \rightarrow [a] \rightarrow [a]$
 $dropWhile p xs$ returns the suffix remaining after $takeWhile p xs$.

Laws:

- $takeWhile p xs ++ dropWhile p xs = xs$
 - $all p (takeWhile p xs) = True$
 - $null (dropWhile p xs) \parallel (not \$ p \$ head \$ dropWhile p xs)$
- if p is total (on xs).

Note: $span p xs = (takeWhile p xs, dropWhile p xs)$

Some Prelude Functions — List Splitting with Predicates

```
takeWhile      :: (a -> Bool) -> [a] -> [a]
takeWhile p [] = []
takeWhile p (x:xs)
  | p x      = x : takeWhile p xs
  | otherwise = []

dropWhile      :: (a -> Bool) -> [a] -> [a]
dropWhile p [] = []
dropWhile p xs@(x:xs')
  | p x      = dropWhile p xs'
  | otherwise = xs

span, break    :: (a -> Bool) -> [a] -> ([a],[a])
span p []     = ([],[a])
span p xs@(x:xs')
  | p x      = let (ys,zs) = span p xs' in (x:ys,zs)
  | otherwise = ([],xs)

break p       = span (not . p)
```

as-Patterns

```
dropWhile      :: (a -> Bool) -> [a] -> [a]
dropWhile p [] = []
dropWhile p xs@(x:xs')
  | p x      = dropWhile p xs'
  | otherwise = xs
```

Consider matching of the third clause against $dropWhile (< 5) [1,2,3]$:

- $p = (< 5)$
- $xs = [1,2,3]$
- $x = 1$
- $xs' = [2,3]$
- $p x = (< 5) 1 = 1 < 5 = True$

Therefore: $dropWhile (< 5) [1,2,3] = dropWhile (< 5) [2,3]$

as-Patterns — 2

```
dropWhile      :: (a -> Bool) -> [a] -> [a]
dropWhile p [] = []
dropWhile p xs@(x:xs')
  | p x      = dropWhile p xs'
  | otherwise = xs
```

Consider matching of the third clause against $dropWhile (< 5) [5,4,3]$:

- $p = (< 5)$
- $xs = [5,4,3]$
- $x = 5$
- $xs' = [4,3]$
- $p x = (< 5) 5 = 5 < 5 = False$

Therefore: $dropWhile (< 5) [5,4,3] = [5,4,3]$

Some Prelude Functions — List Splitting with Predicates

```
takeWhile      :: (a -> Bool) -> [a] -> [a]
takeWhile p [] = []
takeWhile p (x:xs)
  | p x      = x : takeWhile p xs
  | otherwise = []

dropWhile      :: (a -> Bool) -> [a] -> [a]
dropWhile p [] = []
dropWhile p xs@(x:xs')
  | p x      = dropWhile p xs'
  | otherwise = xs

span, break    :: (a -> Bool) -> [a] -> ([a],[a])
span p []     = ([],[a])
span p xs@(x:xs')
  | p x      = let (ys,zs) = span p xs' in (x:ys,zs)
  | otherwise = ([],xs)

break p       = span (not . p)
```