

“Mouldable Code” for Correct-by-Construction SW
Needs Nested Theories —
≈6 years running Agda at the limits of the machine...

WOLFRAM KAHL

McMaster University, Hamilton, Ontario, Canada

17 September 2015

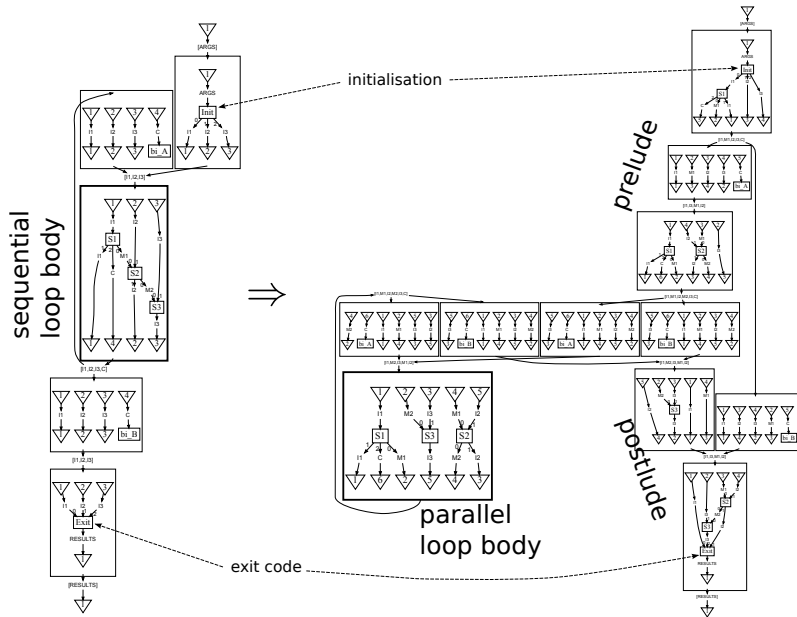
AIM XXII — Leuven, Belgium

This research is supported by the National Science and Engineering Research Council of Canada, NSERC.

“Mouldable Code”? — Background

- 1970s and 1980s at TU Munich: F.L. Bauer’s group: CIP
 - “Computer-aided Intuition-guided Programming”
 - CIP-L: Wide-spectrum language (functional — imperative)
 - CIP-S: (second-order) transformation system
- Gunther Schmidt’s reaction: **Transform Graphs!**
 - Term graph transformation system **HOPS** (several versions)
 - My PhD: second-order term graph transformation
(used relation-algebraic formalisation and proofs)
- My Habilitation: Relation-Algebraic Approach to Graph Transformation
 - relation-algebraically amalgamated DPO and DPB
 - can handle “DPO + graph variables”
- **Coconut** (w. Christopher K. Anand): Software pipelining implemented as code graph transformation
 - generated “vector MASS” library shipped in IBM’s Cell BE SDK
 - implemented in Haskell
 - insufficient support by Haskell type system (no dependent types)

Software Pipelining as Nested Code Graph Transformation



Transformation for Optimisation

- Many transformation patterns
 - are usefully **explained** as graph transformations
 - are normally implemented as AST transformations
- **Implementation** as graph transformations requires:
 - internal representation as graphs (not ASTs)
 - correctness of transformation wrt. graph semantics
 - sufficiently intuitive graph transformation concept

“Mouldable Code” [Gunther Schmidt, 1990s]

- Programs conceptually structured as graphs
- Program development is supported by a graph-based GUI
- Programs are written in a programming language that facilitates correctness proofs
- Program development is supported by a powerful transformation system that allows power-users to “turn the programs inside out”
 - for the purpose of fusion and other efficiency-improving adaptations
 - and also for systematically and without impacting correctness adding what would later become known as “aspects”
- The resulting programs are **correct by construction**

Nested Code Graph Transformation

- Control-flow graphs: Kleene algebra
Kleene categories
- Data-flow graphs: gs-monoidal categories
(tabular allegories)
- **Equations** turn into **transformation rules**
- Matching implemented as graph homomorphisms
- Transformation via variant of DPO approach
 - Correctness wrt. gs-monoidal categories: Zhao Yuhang
 - Correctness wrt. Kleene categories: TBD
- **One-directional rules** can be used for **refinement**
(demonic) Kleene categories

Getting Started — Essential Ingredients

- **RATH-Agda (≈ 500 pages):** Abstract formalisation of semigroupoids, categories, allegories, Kleene categories, collagories, action lattice categories
 - Relatively fine-grained hierarchy of theories
 - Many module splits for performance reasons
 - Allegory and category combinators still slow (>9 GB heap)
- **SUList (≈ 200 pages:)** Sorted unique lists
 - Directly implement sets
 - Key-value-pairs: Finite maps
 - Set-valued maps: Finite relations
 - Invariant-carrying datatype, no irrelevance
 - Many correctness proofs involve large case analyses
 - ≈ 4 GB heap
 - ListSetMap implements Kleene collagory; sub-category of mappings equivalent to FinVecCat
 - ≈ 10 GB heap
- **JSON Parsing and Pretty-printing (≈ 100 pages)**

It Calculates a Pushout! — in 6 seconds...

- A single top-level module brings the three strands together
- Can read and write graphs in JSON format
- Calculates a small (**6 node**) pushout
- MAlonzo:
 - Compilation to Haskell (after typechecking): 40min, >4GB heap
 - GHC call: 40min, >7GB heap
 - Binary size 160MB ; run-time: **≈6s**
 - Probable problem: No compromises:
Invariant-carrying datatypes, no **abstract**, no irrelevance
- UHC (March): Binary size 60MB; segfaults
- UHC whole-program optimisation (-O2,2,2, March):
Binary size 7MB; run-time: **>5min**

Term Graph Decomposition

- Yuhang Zhao implements term graph decomposition into gs-monoidal category expressions [Corradini, Gadducci 1998]
- Concrete model: 2-Category of Term Graphs on top of FinVecCat:
 - Correctness proof involves three levels of categories: Holes unusable
- This is an essential ingredient to proving correctness of DPO term graph rewriting wrt. functorial semantics

Side-Show: AContext

- Abstract formalisation of FCA context categories only needs OCC with powers, residuals, and symmetric quotients
- Agda used as **“just a mechanised mathematical notation”**
that lets me write the mathematics in a natural way
- The **abstract algebraic style** plays to the strengths of Agda
- 189 pages
- Final chapter: Finishing off categoric duality between FCA contexts and complete lower semilattices:
 - Duality proof runs out of 52GB heap
 - One-line definition of the back-and-forth functors takes hours to type-check
 - Issue 1625
 - Andrea Vezzosi supplied experimental patch
 - Will try this week: Does this also help me elsewhere?

Ceterum Censeo ...

... cum grano salis ...

- Agda got many important things right, and has been improving tremendously
 - but even from Agda-dev, I don't get a feeling where Agda is headed
- We need a roadmap towards a trusted kernel
- We need an “Agda report”, perhaps initially limited to the trusted kernel
- We need a roadmap towards self-hosting — Agda in Agda
 - AIM as “Agda hackathon” would profit from the confidence of producing Agda code!
- We need efficient compiled code
 - We need whole-program optimisation
 - We may need *semantics-preserving* pragmas to guide optimisation — not extensions like irrelevance

- First-order sharing is probably not sufficient for efficient type-checking of level-polymorphic code?
- Agda's module system is wonderful to use!
 - Am I the only one using it in certain ways?
 - Documentation of performance implications is needed
 - Nested parameterised modules probably still have problems (Issue 1396)
 - who else besides Ulf understands the implementation of the module system?
 - What would it take for me to understand it?
- Sometimes I look at Agda implementation modules, and lack (pointers to) documentation...
- **I ♥ Agda**