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# Complexity of Self-regular based Infeasible IPM

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# Outline of Presentation

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- Linear Optimization Problem
- Classical primal-Dual IMPs algorithm
- Newton Direction
- Self Regular Function
- Infeasible Self-Regular IPM
- Complexity Of the Algorithm
- Ongoing and Future Work

# The Linear Optimization Problem

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Consider the following LO problem:

$$\begin{array}{ll} \min & c^T x \\ \text{s.t} & Ax = b, \\ & x \geq 0, \end{array}$$

The dual problem is:

$$\begin{array}{ll} \text{(LD)} & \max \quad b^T y \\ & \text{s.t} \quad A^T y + s = c, \\ & \quad \quad s \geq 0, \end{array}$$

where  $y \in \mathcal{R}^m$ ,  $x, s \in \mathcal{R}^n$ .

# The LO Optimality Conditions

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## Optimality Conditions

$$\begin{aligned}Ax &= b, \\A^T y + s &= c, \\Xs &= 0, \\x \geq 0 \quad s \geq 0,\end{aligned}$$

where  $X = \text{diag}(x)$ .

# The Lo Central Path

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The **primal-dual central path** is defined as the set of solutions  $(x(\mu))$  and  $(y(\mu), s(\mu))$  for  $\mu > 0$  of the system

$$\begin{aligned}Ax &= b, \\A^T y + s &= c, \\Xs &= \mu e, \\x > 0 \quad s > 0,\end{aligned}$$

where  $e = [1, 1, \dots, 1]^T$ .

IPM condition:

$$Ax^0 = b, x^0 > 0 \quad A^T y + s = c, s > 0$$

# Classical Primal-Dual Newton Method for LO

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## inputs

A proximity parameter  $\tau$ ;

an accuracy parameter  $\epsilon > 0$ ;

$(x^0, s^0)$  and  $\mu^0 = 1$  such that  $\Phi(x^0, s^0, \mu^0) \leq \tau$ .

$x := x^0$ ;  $s := s^0$ ;  $\mu := \mu^0$ ;

**while**  $n\mu \geq \epsilon$

$\mu := (1 - \theta)\mu$ ;

**while**  $\Phi(x, s, \mu) \geq \tau$

Solve Newton system for  $\Delta x, \Delta y, \Delta s$ ;

determine a step size  $\alpha$ ;

$x := x + \alpha\Delta x$ ;

$s := s + \alpha\Delta s$ ;

$y := y + \alpha\Delta y$ .

**end**

**end**

# Newton Direction

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$$\begin{pmatrix} A & 0 & 0 \\ 0 & A^T & I \\ S & 0 & X \end{pmatrix} \begin{pmatrix} \Delta x \\ \Delta y \\ \Delta s \end{pmatrix} = \begin{pmatrix} -r_b \\ -r_c \\ \mu e - Xs \end{pmatrix},$$

$$r_b = Ax - b, \quad r_c = A^T y + s - c$$

$$v = \sqrt{\frac{xs}{\mu}}, \quad d_x = \frac{v\Delta x}{x}, \quad d_s = \frac{v\Delta s}{s}$$

$$\begin{pmatrix} \bar{A} & 0 & 0 \\ 0 & \bar{A}^T & I \\ I & 0 & I \end{pmatrix} \begin{pmatrix} d_x \\ \Delta y \\ d_s \end{pmatrix} = \begin{pmatrix} -r_b \\ -r_c \\ v^{-1} - v \end{pmatrix}, \quad \bar{A} = \frac{1}{\mu}AV^{-1}X$$

$$\Psi(V) = \frac{e^T v^2 - n}{2} - \sum_{i=1}^n \log(v_i)$$

# Newton Direction

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$$(AD^2A^T)\Delta y = AD^2r_h - r_b,$$

$$\Delta x = D^2(A^T \Delta y - r_h),$$

$$\Delta s = x^{-1}(-\mu v \nabla \Psi(v) - s \Delta x),$$

$$x(\alpha) := x + \alpha \Delta x,$$

$$y(\alpha) := y + \alpha \Delta y,$$

$$s(\alpha) := s + \alpha \Delta s.$$

# Complexity of small-update and large-update

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small-update:

$$O(\sqrt{n} \log(\frac{n}{\epsilon}))$$

large-update:

$$O(n \log(\frac{n}{\epsilon}))$$

# Self-Regular Function

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A function  $\psi(t) \in \mathcal{C}^2 : (0, \infty) \rightarrow \mathbb{R}$  is self-regular if it satisfies the following conditions:

SR.1  $\psi(t)$  is strictly convex with respect to  $t > 0$  and vanishes at its global minimal point  $t = 1$ , i.e.,  $\psi(1) = \psi'(1) = 0$ . Further, there exist positive constants  $\nu_2 \geq \nu_1 > 0$  and  $p \geq 1, q \geq 1$  such that

$$\nu_1(t^{p-1} + t^{-1-q}) \leq \psi''(t) \leq \nu_2(t^{p-1} + t^{-1-q}), \quad \forall t \in (0, \infty);$$

SR.2 For any  $t_1, t_2 > 0$ ,

$$\psi(t_1^r t_2^{1-r}) \leq r\psi(t_1) + (1-r)\psi(t_2), \quad \forall r \in [0, 1].$$

$$\Upsilon_{p,q}(t) = \frac{t^{p+1} - 1}{p(p+1)} + \frac{t^{1-q} - 1}{q(q-1)} + \frac{p-q}{pq}(t-1), \quad p, q \geq 1,$$

# Self-Regular Function

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with  $\nu_1 = \nu_2 = 1$ . The second family is defined as

$$\Gamma_{p,q}(t) = \frac{t^{p+1} - 1}{p + 1} + \frac{t^{1-q} - 1}{q - 1}, \quad p \geq 1, \quad q > 1,$$

with  $\nu_1 = 1$  and  $\nu_2 = q$ .

$$\Psi(v) = \sum_{i=1}^n \psi(v_i), \quad v := \sqrt{\frac{xs}{\mu}}$$

Best complexity for feasible IPMs :  $O(\sqrt{n} \log(n) \log(\frac{n}{\epsilon}))$

$$\begin{pmatrix} \bar{A} & 0 & 0 \\ 0 & \bar{A}^T & I \\ I & 0 & I \end{pmatrix} \begin{pmatrix} dx \\ \Delta y \\ ds \end{pmatrix} = \begin{pmatrix} -r_b \\ -r_c \\ -\nabla \Psi(v) \end{pmatrix},$$

# Properties of Proximity Function

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$$\Phi(x, s, \mu) = \Psi(v) = \frac{e^T v^2 - n}{2} + \frac{e^T v^{-2} - n}{2}$$

$$\mu_g = \frac{x^T s}{n}, \quad \mu_h = \frac{n}{x^{-T} s^{-1}}, \quad \mu^* = \sqrt{\frac{x^T s}{x^{-T} s^{-1}}} = \sqrt{\mu_g \mu_h}$$

$$\mu_t = \frac{2\mu_g}{\tau + 1 + \sqrt{(\tau + 1)^2 - \frac{4\mu_g}{\mu_h}}}, \quad \Phi(x, s, \mu_t) = \frac{(\tau - 1)n}{2}$$

$$\mu_t \leq \mu_h \leq \mu^* \leq \mu_g$$

$$\Phi(x, s, \mu_g) = \Phi(x, s, \mu_h)$$

# Properties of Proximity Function

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$$\Phi(x, s, \mu_g) = \Phi(x, s, \mu^*) + \frac{\Phi(x, s, \mu^*)^2}{2n}$$

$$\frac{\mu_g}{\mu_h} = \frac{x^T s x^{-T} s^{-1}}{n^2} \leq \tau \iff \Phi(x, s, \frac{\mu_g}{\tau}) \leq \frac{(\tau - 1)n}{2}$$

# Infeasible Neighborhood

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$$\mathcal{N}(\tau, \beta) = \left\{ (x, y, s) \mid \Psi(v) \leq \frac{(\tau - 1)n}{2}, \|r_b\| \leq \|r_b^0\| \frac{\mu_g}{\mu^0} \beta, \|r_c\| \leq \|r_c^0\| \frac{\mu_g}{\mu^0} \beta \right\}$$

$$(x^0, y^0, s^0) = (\zeta e, 0, \zeta e)$$

$$\sigma^2 = \|d_x + d_s\|^2, \sigma_1^2 = d_x^2 + d_s^2 = \sigma^2 - 2d_x^T d_s$$

$$C_3 n^2 \sigma^{\frac{2}{3}} \leq \sigma_1^2 \leq C_3 n^2 \sigma^{\frac{2}{3}},$$

# SR-Infeasible IPMs

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## Input:

Proximity parameters  $\tau \geq 10$  and  $\beta \geq 1$ , Neighborhood  $\mathcal{N}(\tau, \beta)$ ;

Initial point  $(x^0, y^0, s^0) \in \mathcal{N}(\tau, \beta)$ ; such that  $(x^0, s^0) > 0$

An accuracy parameter  $\epsilon > 0$ , and damping factor  $\alpha$ .

**begin**

**while**  $\max \left\{ (x^k)^T s^k, \left\| r_b^k \right\|, \left\| r_c^k \right\| \right\} \geq \epsilon$  **do**

**Begin**

If  $\frac{\mu_g^k}{\mu_h^k} \geq \frac{\tau}{2}$  then  $\mu := \mu_h^k$ ; otherwise

$\mu := \mu_t^k$ .

Solve system for  $\Delta x^k, \Delta y^k, \Delta s^k$ .

**Begin**

Determine a step size  $\alpha_k$  such that

$$\Phi(x(\alpha_k), s(\alpha_k), \mu_t^k) \leq \Phi(x^k, s^k, \mu_t^k) - \frac{\alpha^*}{2} \Phi(x, s, \mu_t)$$

and  $(x(\alpha_k), y(\alpha_k), s(\alpha_k)) \in \mathcal{N}(\tau, \beta)$ ;

$$x^{k+1} := x(\alpha_k); \quad y^{k+1} := y(\alpha_k); \quad s^{k+1} := s(\alpha_k);$$

$$k = k + 1.$$

**end**

**end**

# Worst Case Complexity

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$$\Phi(x(\alpha^*), s(\alpha^*), \mu_g(\alpha^*)) \leq \Phi(x^k, s^k, \mu_t^k)$$

$$\Phi(x(\alpha^*), s(\alpha^*), \mu_t^k) \leq \Phi(x^k, s^k, \mu_t^k) - \frac{\alpha^*}{2} \Phi(x^k, s^k, \mu_t^k)$$

$$\mu_t^{k+1} \leq \left(1 - \frac{\alpha^*}{4}\right) \mu_t^k$$

$$O\left(n^2 \log\left(\frac{n}{\epsilon}\right)\right)$$

# Ongoing and future work

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- Dynamic Large-update SR IPMs Developed for a big class of SR functions.
- Implementation of SR Proximity Based IPMs has been Done by X. Zhu and guoqing Zhang.
- Developing new algorithm for pure primal or pure dual problem based on SR proximity.