

# Eigenvalue Problems and Singular Value Decomposition

Sanzheng Qiao

Department of Computing and Software  
McMaster University

August, 2012

# Outline

- 1 Eigenvalue Problems
- 2 Singular Value Decomposition
- 3 Software Packages

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# Problem setting

Eigenvalue problem:

$$A\mathbf{x} = \lambda\mathbf{x},$$

$\lambda$ : eigenvalue

$\mathbf{x}$ : right eigenvector.

$\mathbf{y}^H A = \lambda \mathbf{y}^H$ ,  $\mathbf{y}$  left eigenvector.

# Canonical forms

Decomposition:

$$A = SBS^{-1}$$

where  $B$  is in a canonical (simple) form, whose eigenvalues and eigenvectors can be easily obtained.

- $A$  and  $B$  have the same eigenvalues. (They are similar.)
- If  $\mathbf{x}$  is an eigenvector of  $B$ , then  $S\mathbf{x}$  is the eigenvector of  $A$  corresponding to the same eigenvalue.

# Jordan canonical form

$$A = SJS^{-1}, \quad J = \text{diag}(J_{n_1}(\lambda_1), \dots, J_{n_k}(\lambda_k))$$

$$J_{n_i}(\lambda_i) = \begin{bmatrix} \lambda_i & 1 & & 0 \\ & \ddots & \ddots & \\ & & \ddots & 1 \\ 0 & & & \lambda_i \end{bmatrix}.$$

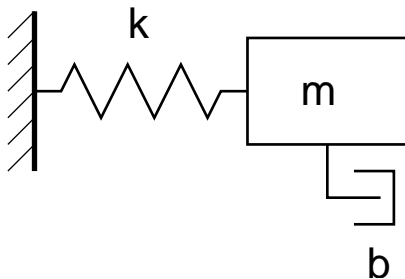
# Jordan canonical form

- The algebraic multiplicity of  $\lambda_i$  is  $n_i$ .
- A Jordan block has one right eigenvector  $[1, 0, \dots, 0]^T$  and one left eigenvector  $[0, \dots, 0, 1]^T$ .
- If all  $n_i = 1$ , then  $J$  is diagonal,  $A$  is called diagonalizable; otherwise,  $A$  is called defective.
- An  $n$ -by- $n$  defective matrix has fewer than  $n$  eigenvectors.

# Example

In practice, confronting defective matrices is a fundamental fact.

Mass-spring problem





# Mass-spring problem

Newton's law  $F = ma$  implies

$$m\ddot{x}(t) = -kx(t) - b\dot{x}(t).$$

Let

$$\mathbf{y}(t) = \begin{bmatrix} \dot{x}(t) \\ x(t) \end{bmatrix},$$

we transform the second order ODE into a system of the first order ODEs

$$\dot{\mathbf{y}}(t) = \begin{bmatrix} -\frac{b}{m} & -\frac{k}{m} \\ 1 & 0 \end{bmatrix} \mathbf{y}(t) =: \mathbf{A}\mathbf{y}(t).$$

# Mass-spring problem

The characteristic polynomial of  $A$  is

$$\lambda^2 + \frac{b}{m}\lambda + \frac{k}{m}$$

and the eigenvalues are

$$\lambda_{\pm} = \frac{-\frac{b}{m} \pm \sqrt{\left(\frac{b}{m}\right)^2 - 4\frac{k}{m}}}{2} = \frac{b}{2m} \left( -1 \pm \sqrt{1 - \frac{4km}{b^2}} \right).$$

When  $4km/b^2 = 1$ , critically damped, two equal eigenvalues,  $A$  is not diagonalizable.

# Jordan canonical form

$$A = SJS^{-1}$$

$$A = \begin{bmatrix} -\frac{b}{m} & -\frac{k}{m} \\ 1 & 0 \end{bmatrix}, \quad 4km = b^2,$$

$$J = \begin{bmatrix} -\frac{b}{2m} & 1 \\ 0 & -\frac{b}{2m} \end{bmatrix}, \quad S = \begin{bmatrix} -\frac{b}{2m} & 1 - \frac{b}{2m} \\ 1 & 1 \end{bmatrix}$$

# Jordan canonical form

It is undesirable to compute Jordan form, because

- Jordan block is discontinuous

$$J(0) = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix} \quad J(\epsilon) = \begin{bmatrix} \epsilon & 1 \\ 0 & 2\epsilon \end{bmatrix},$$

while  $J(0)$  has an eigenvalue of multiplicity two,  $J(\epsilon)$  has two simple eigenvalues.

- In general, computing Jordan form is unstable, that is, there is no guarantee that  $\widehat{S}\widehat{J}\widehat{S}^{-1} = A + E$  for a small  $E$ .

# Schur canonical form

$$A = QTQ^H$$

$Q$ : unitary

$T$ : upper triangular

The eigenvalues of  $A$  are the diagonal elements of  $T$ .

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Real case

$$A = QTQ^T$$

$Q$ : orthogonal

$T$ : quasi-upper triangular, 1-by-1 or 2-by-2 blocks on the diagonal.

# Conditioning

Let  $\lambda$  be a simple eigenvalue of  $A$  with unit right eigenvector  $\mathbf{x}$  and left eigenvector  $\mathbf{y}$ .

$\lambda + \epsilon$  be the corresponding eigenvalue of  $A + E$ , then

$$\epsilon = \frac{\mathbf{y}^H E \mathbf{x}}{\mathbf{y}^H \mathbf{x}} + O(\|E\|^2)$$

or

$$|\epsilon| \leq \frac{1}{|\mathbf{y}^H \mathbf{x}|} \|E\| + O(\|E\|^2).$$

Condition number for finding a simple eigenvalue

$$1/|\mathbf{y}^H \mathbf{x}|$$

# Computing the Schur decomposition

$$A = QTQ^T, \quad T : \text{quasi-upper triangular}$$

Step 1: Reduce  $A$  to upper Hessenberg

$$A = Q_1 H Q_1^T, \quad h_{ij} = 0, \quad i > j + 1$$

Step 2: Compute the Schur decomposition of  $H$

$$H = Q_2 T Q_2^T$$



# Introducing zeros into a vector

Householder transformation

$$H = I - 2\mathbf{u}\mathbf{u}^T \quad \text{with } \mathbf{u}^T\mathbf{u} = 1$$

$H$  is symmetric and orthogonal ( $H^2 = I$ ).

Goal:  $H\mathbf{a} = \alpha\mathbf{e}_1$ .

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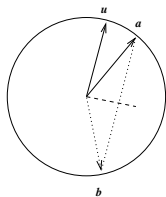
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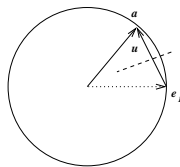
Choose

$$\mathbf{u} = \mathbf{a} \pm \|\mathbf{a}\|_2 \mathbf{e}_1$$

# A geometric interpretation



(a)



(b)

Figure (a) shows the image  $\mathbf{b} = (I - 2\mathbf{u}\mathbf{u}^T)\mathbf{a}$  for an arbitrary  $\mathbf{u}$ ,  
 in figure (b),  $\mathbf{u} = \mathbf{a} - \|\mathbf{a}\|_2 \mathbf{e}_1$ .

# Computing Householder transformations

Given a vector  $\mathbf{x}$ , it computes scalars  $\sigma$ ,  $\alpha$ , and vector  $\mathbf{u}$  such that

$$(I - \sigma^\dagger \mathbf{u} \mathbf{u}^T) \mathbf{x} = -\alpha \mathbf{e}_1$$

where  $\sigma^\dagger = 0$  if  $\sigma = 0$  and  $\sigma^{-1}$  otherwise

- $\alpha = \text{sign} x_1 \|\mathbf{x}\|_2$
- $\mathbf{u} = \mathbf{x} + \alpha \mathbf{e}_1$
- $\|\mathbf{u}\|_2^2 = 2(\alpha^2 + \alpha x_1) = 2\alpha u_1$

# house.m

```
function [u,sigma,alpha] = house(x)
```

```
u = x;
```

```
alpha = sign(u(1))*norm(u);
```

```
u(1) = u(1) + alpha;
```

```
sigma = alpha*u(1);
```

# Reducing $A$ to upper Hessenberg

```
n = length(A(1,:));
Q = eye(n);

for j=1:(n-2)
    [u, sigma, alpha] = myhouse(A(j+1:n, j));
    if sigma ~= 0.0
        for k = j:n
            A(j+1:n, k) = A(j+1:n, k) - ((u'*A(j+1:n, k))/sigma)*u;
        end %for k
        for i=1:n
            A(i, j+1:n) = A(i, j+1:n) - ((A(i, j+1:n)*u)/sigma)*u';
            Q(i, j+1:n) = Q(i, j+1:n) - ((Q(i, j+1:n)*u)/sigma)*u';
        end %for i
    end %if
end %for j
```

# Computing eigenvalues and eigenvectors

Suppose  $A$  has distinct eigenvalues  $\lambda_i$ ,  $i = 1, \dots, n$ , where  $|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|$ , and  $\mathbf{x}_i$  are the eigenvectors (linear independent).

An arbitrary vector  $\mathbf{u}$  can be expressed as

$$\mathbf{u} = \mu_1 \mathbf{x}_1 + \mu_2 \mathbf{x}_2 + \dots + \mu_n \mathbf{x}_n$$

If  $\mu_1 \neq 0$ ,  $A^k \mathbf{u}$  has almost the same direction as  $\mathbf{x}_1$  when  $k$  is large and thus  $(\lambda_i/\lambda_1)^k$  ( $i > 1$ ) is small.

Thus the Rayleigh quotient

$$\frac{(A^k \mathbf{u})^T A (A^k \mathbf{u})}{(A^k \mathbf{u})^T (A^k \mathbf{u})} \approx \lambda_1.$$

# Power method

Initial  $\mathbf{u}_0$ ;  $i = 0$ ;

repeat

$$\mathbf{v}_{i+1} = \mathbf{A}\mathbf{u}_i;$$

$$\mathbf{u}_{i+1} = \mathbf{v}_{i+1} / \|\mathbf{v}_{i+1}\|_2;$$

$$\tilde{\lambda}_{i+1} = \mathbf{u}_{i+1}^T \mathbf{A} \mathbf{u}_{i+1};$$

$$i = i + 1$$

until convergence



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## Problems

- Computes only  $(\lambda_1, \mathbf{x}_1)$
- Converges slowly when  $|\lambda_1| \approx |\lambda_2|$
- Does not work when  $|\lambda_1| = |\lambda_2|$

# Inverse power method

Suppose that  $\mu$  is an estimate of  $\lambda_k$ , then  $(\lambda_k - \mu)^{-1}$  is the dominant eigenvalue of  $(A - \mu I)^{-1}$ . Applying the power method to  $(A - \mu I)^{-1}$ , we can compute  $\mathbf{x}_k$  and  $\lambda_k$ .

Example.

Eigenvalues of  $A$ :  $-1, -0.2, 0.5, 1.5$

Shift  $\mu$ :  $-0.8$

Eigenvalues of  $(A - \mu I)^{-1}$ :  $-5.0, 1.7, 0.78, 0.43$

Very effective when we have a good estimate for an eigenvalue.

# QR method

Goal: Generate a sequence

$$A_0 = A, A_1, \dots, A_{k+1}$$

$$A_{i+1} = Q_i^T A Q_i = \begin{bmatrix} B & \mathbf{u} \\ \mathbf{s}^T & \mu \end{bmatrix}$$

where  $\mathbf{s}$  is small and  $Q_i$  is orthogonal, i.e.,  $Q_i^T = Q_i^{-1}$  (so  $A_{k+1}$  and  $A$  have the same eigenvalues).

- Since  $\mathbf{s}$  is small,  $\mu$  is an approximation of an eigenvalue of  $A_{k+1}$  ( $A$ );
- Deflate  $A_{k+1}$  and repeat the procedure on  $B$  when  $\mathbf{s}$  is sufficiently small. The problem size is reduced by one.

# QR method

What does  $Q_k$  look like?

If the last column of  $Q_k$  is a left eigenvector  $\mathbf{y}$  of  $A$ , then

$$\begin{aligned} Q_k^T A Q_k &= \begin{bmatrix} P_k^T \\ \mathbf{y}^T \end{bmatrix} A [P_k \mathbf{y}] \\ &= \begin{bmatrix} P_k^T \\ \mathbf{y}^T \end{bmatrix} [A P_k \ A \mathbf{y}] \\ &= \begin{bmatrix} B & \mathbf{u} \\ 0^T & \lambda \end{bmatrix} \end{aligned}$$

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One step of the inverse power method: Solve for  $\mathbf{q}$  in  $(A - \mu I)^T \mathbf{q} = \mathbf{e}_n$ , where  $\mu$  is an estimate for an eigenvalue of  $A$ . (How? Later.)

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How do we construct an orthogonal  $Q$  whose last column is  $\mathbf{q}$ ?  
If  $(A - \mu I) = QR$  is the QR decomposition and  $\mathbf{q}$  is the last column of  $Q$ , then

$$\mathbf{q}^T (A - \mu I) = \mathbf{q}^T QR = r_{n,n} \mathbf{e}_n^T.$$

Thus, after normalizing,

$$(A - \mu I)^T \mathbf{q} = \mathbf{e}_n.$$



## QR method

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$Q$  can be obtained from the QR decomposition  $(A - \mu I) = QR$ .

# QR decomposition

QR decomposition of an upper Hessenberg matrix using the Givens rotations.

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Givens rotation

$$G = \begin{bmatrix} \cos \theta & \sin \theta \\ -\sin \theta & \cos \theta \end{bmatrix}$$

Introducing a zero into a 2-vector:

$$G \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} \times \\ 0 \end{bmatrix}$$

i.e., rotate  $\mathbf{x}$  onto  $x_1$ -axis.

# Computing the Givens rotations

Given a vector  $[a \ b]^T$ , compute  $\cos \theta$  and  $\sin \theta$  in the Givens rotation.

$$\cos \theta = \frac{a}{\sqrt{a^2 + b^2}} \quad \sin \theta = \frac{b}{\sqrt{a^2 + b^2}}$$

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```
function [c, s] = grotate(a, b)

if (b = 0) c = 1.0; s = 0.0; return; end;

if (abs(b) >= abs(a))
    ct = a/b;
    s = 1/sqrt(1 + ct*ct); c = s*ct;
else
    t = b/a;
    c = 1/sqrt(1 + t*t); s = c*t;
end
```

# QR decomposition

Compute the QR decomposition  $H = QR$  of an upper Hessenberg matrix  $H$  using the Givens rotations.

```
function [R, Q] = hqrd(H)

n = length(H(1,:));
R = H; Q = eye(n);

for j=1:n-1
    [c, s] = grotate(R(j,j), R(j+1,j));
    R(j:j+1, j:n) = [c s; -s c]*R(j:j+1, j:n);
    Q(:, j:j+1) = Q(:, j:j+1)*[c -s; s c];
end
```

# QR method

But, we want

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- similarity transformations of  $A$ , not  $A - \mu I$ ;
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$$A - \mu I = QR.$$

$$RQ = Q^T(A - \mu I)Q = Q^T A Q - \mu I.$$

$$RQ + \mu I = Q^T A Q \text{ is similar to } A.$$



# QR method

## One step of QR method

```
repeat
  choose a shift  $\mu$ ;
  QR decomposition  $A - \mu \cdot I = QR$ ;
   $A = RQ + \mu \cdot I$ ;
until convergence ( $A(n, 1:n-1)$  small)
```

# QR method

If  $A$  has been reduced to the upper Hessenberg form, the structure is maintained during the iteration.

$H_0$  is upper Hessenberg;

$H_0 - \mu I$  is upper Hessenberg;

$H_0 - \mu I = QR$ ,  $R$  is upper triangular and  $Q$  is upper Hessenberg;

$H_1 = RQ + \mu I$  is upper Hessenberg;

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Implication: The QR decomposition is cheap (only eliminate the subdiagonal).

## Choosing the shift

Since the last element converges to an eigenvalue, it is reasonable to choose  $h_{n,n}$  as the shift. But, it doesn't always work. A more general method is to choose the eigenvalue of the trailing 2-by-2 submatrix

$$\begin{bmatrix} h_{n-1,n-1} & h_{n-1,n} \\ h_{n,n-1} & h_{n,n} \end{bmatrix}$$

that is close to  $h_{n,n}$ . Heuristically, it is more effective than choosing  $h_{n,n}$  especially in the beginning.

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What if the trailing 2-by-2 submatrix has a complex conjugate pair of eigenvalues? The double shift strategy can be used to overcome the difficulty. In general, the Francis QR method using double implicit shift strategy can reduce a real Hessenberg matrix into the real Schur form.

## Symmetric case

A symmetric matrix is diagonalizable:  $A = Q\Lambda Q^T$ ,  
 $\Lambda = \text{diag}(\lambda_1, \dots, \lambda_n)$ .

QR method. This method is very efficient if only all eigenvalues are desired or all eigenvalues and eigenvectors are desired and the matrix is small ( $n \leq 25$ ).

- 1 Reduce  $A$  to symmetric tridiagonal. This costs  $\frac{4}{3}n^3$  or  $\frac{8}{3}n^3$  if eigenvectors are also desired.
- 2 Apply the QR iteration to the tridiagonal. On average, it takes two QR steps per eigenvalue. Finding all eigenvalues takes  $6n^2$ . Finding all eigenvalues and eigenvectors requires  $6n^3$ .

# Example

$$A = \begin{bmatrix} 1 & 2 & 3 & 4 \\ 2 & 1 & 2 & 3 \\ 3 & 2 & 1 & 2 \\ 4 & 3 & 2 & 1 \end{bmatrix}$$

After tridiagonalization

$$\begin{bmatrix} 1.0000 & -5.3852 & 0 & 0 \\ -5.3852 & 5.1379 & -1.9952 & 0 \\ 0 & -1.9952 & -1.3745 & 0.2895 \\ 0 & 0 & 0.2895 & -0.7634 \end{bmatrix}$$

# Example

	$\mu$	$\beta_1$	$\beta_2$	$\beta_3$
1	-0.6480	3.8161	0.2222	-0.0494
2	-0.5859	1.2271	0.0385	$10^{-5}$
3	-0.5858	0.3615	0.0070	converge
4	-1.0990	0.0821	$10^{-10}$	
5	-1.0990	0.0186	converge	



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# Introduction

$$A = U\Sigma V^T$$

$A$ :  $m$ -by- $n$  real matrix ( $m \geq n$ )

$U$ :  $m$ -by- $m$  orthogonal

$V$ :  $n$ -by- $n$  orthogonal

$\Sigma$ : diagonal,  $\text{diag}(\sigma_i)$ ,

$\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_n \geq 0$

Singular values:  $\sigma_i$

Left singular vectors: columns of  $U$

Right singular vectors : columns of  $V$

# Introduction

SVD reveals many important properties of a matrix  $A$ . For example,

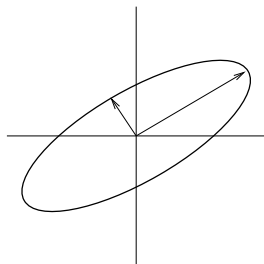
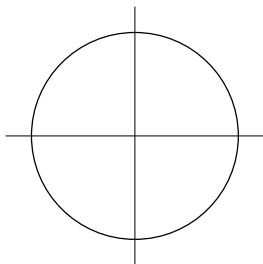
- The number of nonzero singular values is the rank of  $A$ .  
Suppose  $\sigma_k > 0$  and  $\sigma_{k+1} = 0$ , then  $\text{rank}(A) = k$ . If  $k < n$ , the columns of  $A$  are linearly dependent. ( $A$  is rank deficient.)
- If  $\sigma_n > 0$  ( $A$  is of full rank),

$$\text{cond}(A) = \frac{\sigma_1}{\sigma_n}$$

# A geometric interpretation

Transformation  $A: \mathbf{x} \rightarrow A\mathbf{x}$

$$\sigma_1 \geq \frac{\|A\mathbf{x}\|}{\|\mathbf{x}\|} \geq \sigma_n$$



# Application: Linear least-squares problem

$$\min_{\mathbf{x}} \|\mathbf{Ax} - \mathbf{b}\|_2^2$$

also called linear regression problem in statistics.

$$\text{SVD: } A = U\Sigma V^T$$

$$\|\mathbf{Ax} - \mathbf{b}\|_2^2 = \|\Sigma\mathbf{z} - \mathbf{d}\|_2^2$$

where

$$\mathbf{d} = U^T\mathbf{b} \quad \mathbf{z} = V^T\mathbf{x}$$

# Application: Linear least-squares problem

## Solution

$$z_j = \frac{d_j}{\sigma_j} \quad \text{if } \sigma_j \neq 0$$
$$z_j = \text{anything} \quad \text{if } \sigma_j = 0$$

Usually, we set

$$z_j = 0 \quad \text{if } \sigma_j = 0$$

for minimum norm solution.

This allows us to solve the linear least-squares problems with singular  $A$ .

## Application: Principal component analysis

Suppose that  $A$  is a data matrix. For example, each column contains samples of a variable. It is frequently standardized by subtracting the means of the columns and dividing by their standard deviations.

If  $A$  is standardized, then  $A^T A$  is the correlation matrix.

If the variables are strongly correlated, there are few components, fewer than the number of variables, can predict all the variables.

# Application: Principal component analysis

In terms of SVD, let

$$A = U\Sigma V^T$$

be the SVD of  $A$ . If  $A$  is  $m$ -by- $n$  ( $m \geq n$ ), we partition  $U = [\mathbf{u}_1, \dots, \mathbf{u}_m]$  and  $V = [\mathbf{v}_1, \dots, \mathbf{v}_n]$ . Then we can write

$$A = \sigma_1 \mathbf{u}_1 \mathbf{v}_1^T + \dots + \sigma_n \mathbf{u}_n \mathbf{v}_n^T.$$

If the variables are strongly correlated, there are few singular values that are significantly larger than the others.



## Application: Principal component analysis

In other words, we can find an  $r$  such that  $\sigma_r \gg \sigma_{r+1}$ . We can use the rank  $r$  matrix

$$A_r = \sigma_1 \mathbf{u}_1 \mathbf{v}_1^T + \cdots + \sigma_r \mathbf{u}_r \mathbf{v}_r^T$$

to approximate  $A$ . In fact,  $A_r$  is the closest (in Frobenius norm) rank  $r$  approximation to  $A$ . Usually, we can find  $r \ll n$ .

It is also called low-rank approximation.

Principal component analysis is used in a wide range of fields. In image processing, for example,  $A$  is a 2D image.

# Computing SVD

Note that

- the columns of  $U$  are the eigenvectors of  $AA^T$  (symmetric and positive semi-definite);
- the columns of  $V$  are eigenvectors of  $A^T A$ .

The algorithm is parallel to the QR method for symmetric eigenvalue decomposition.

We work on  $A$  instead of  $A^T A$ .

# Computing the SVD

- 1 Bidiagonalize  $A$  using Householder transformations ( $A \rightarrow B$  is upper bidiagonal and  $B^T B$  tridiagonal);
- 2 Implicit QR iteration
  - 1 Find a Givens rotation  $G_1$  from the first column of  $B^T B - \mu I$ ;
  - 2 Apply  $G_1$  to  $B$ ;
  - 3 Apply a sequence of rotations, left and right, to restore the bidiagonal structure of  $B$ ;

The shift  $\mu$  is obtained by calculating the eigenvalues of the 2-by-2 trailing submatrix of  $B^T B$ .

# Outline

- 1 Eigenvalue Problems
- 2 Singular Value Decomposition
- 3 Software Packages**

# Software packages

**NETLIB** LAPACK: sgees (Schur form), sgeev (eigenvalues and eigenvectors), ssyev (symmetric and dense eigenproblems), sstev (symmetric and tridiagonal eigenproblems), sgesvd (SVD), sbdsqr (small and dense SVD)

**IMSL** evcrg, evcsf, lsvrr

**MATLAB/Octave** schur, eig, svd

**NAG** f02agf, f02abf, f02wef

# Summary

- Eigenvalue decomposition, Jordan and Schur canonical forms
- Condition number for eigenvalue
- Householder transformation, Givens rotation
- QR method
- Singular value decomposition
- Linear least-squares problem
- Low-rank approximation

# References

- [1] ] Z. Bai, J. Demmel, J. Dongarra, A. Ruhe, and H. van der Vorst, editors. *Templates for the Solution of Algebraic Eigenvalue Problems: A Practical Guide*. SIAM, Philadelphia, 2000.
- [2] ] James W. Demmel. *Applied Numerical Linear Algebra*. SIAM, Philadelphia, 1997.  
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- [3] ] G. H. Golub and C. F. Van Loan. *Matrix Computations*, 3rd Ed. The Johns Hopkins University Press, Baltimore, MD, 1996.  
Ch 7, Ch 8