

## A Symmetric Rank-Revealing Toeplitz Matrix Decomposition

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**Abstract.** In signal and image processing, regularization often requires a rank-revealing decomposition of a symmetric Toeplitz matrix with a small rank deficiency. In this paper, we present an efficient factorization method that exploits symmetry as well as the rank and Toeplitz properties of the given matrix.

**Keywords:** Toeplitz matrix, regularization, symmetric rank-revealing decomposition

### 1. Introduction

In signal and image processing applications [5], [6], a noisy and distorted signal vector  $\hat{x}$  is given by

$$\hat{x} = Tx + w, \quad (1)$$

where  $x$  and  $w$  represent an unknown original signal vector and a noise vector, respectively, and  $T$  is a predetermined matrix describing the spread of signals. This problem arises often in array processing, where the matrix  $T$  may be real, symmetric, and Toeplitz. Assuming the dimensions of  $T$  to be  $n \times n$ , we have

$$T = \begin{pmatrix} t_1 & t_2 & t_3 & \dots & t_n \\ t_2 & t_1 & t_2 & \dots & t_{n-1} \\ t_3 & t_2 & t_1 & \dots & t_{n-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ t_n & t_{n-1} & t_{n-2} & \dots & t_1 \end{pmatrix}.$$

To restore the original signal vector  $x$  from  $\hat{x}$ , we should invert  $T$ . Quite frequently, cf. [1],  $T$  is ill-conditioned and regularization is required. One popular method is the truncated singular value decomposition (TSVD) [4], which requires  $O(n^3)$  floating-point operations (flops) because it is not known how to compute the singular value decomposition (SVD) of a Toeplitz matrix in fewer flops.

For general matrices, less expensive rank-revealing methods like the  $URV$  decomposition have been developed by Stewart [7] and others to replace the SVD. But  $O(n^3)$  flops are still required. If the matrix is Toeplitz and banded with bandwidth  $b$ , Nagy [6] proposed an  $O(bn^2)$  method for computing an approximate  $URV$  decomposition. However, no one has shown how to exploit the symmetry of  $T$ . We suggest a possible approach in Section 2.

In addition, assume that  $T$  has a small rank deficiency, viz.,

$$\text{rank}(T) = n - k,$$

where  $k$  denotes a small integer. We will present an  $O(kn^2)$  method in Sections 3 to 5 for computing a rank-revealing factorization. Our other contribution is to show how to avoid complex arithmetic in the computation. The examples in Section 6 show that our new method restores the signal almost as accurately as the TSVD approach.

The paper is organized as follows. We present an extension of rank-revealing factorizations to symmetric matrices in Section 2. In Section 3 we sketch a fast  $O(n^2)$  triangularization scheme based on hyperbolic rotations, and in Section 4 we show how to avoid the use of complex arithmetic when these rotations are applied. Lastly, in Sections 5 and 6, we discuss rank-revealing techniques and present simulation results.

## 2. Use Symmetry

Many popular regularization techniques are based on an SVD of  $T$ :

$$T = U\Sigma V^T,$$

where  $U$  and  $V$  are orthogonal, and  $\Sigma = \text{diag}(\sigma_i)$ , with

$$\sigma_1 \geq \cdots \geq \sigma_n \geq 0.$$

Select a tolerance  $\tau$  to find  $k$  such that

$$\sigma_{n-k} > \tau \geq \sigma_{n-k+1}.$$

We may say that  $T$  has  $k$  small singular values, or that  $T$  has a numerical rank of  $n - k$ . Let

$$\hat{\Sigma} = \text{diag}(\sigma_1, \dots, \sigma_{n-k}, 0, \dots, 0).$$

We restore the original signal  $x$  via

$$x = V\hat{\Sigma}^+U^T\hat{x},$$

where  $\hat{\Sigma}^+ = \text{diag}(\sigma_i^+)$  denotes the pseudo-inverse of  $\hat{\Sigma}$ .

Although the TSVD is stable, it is expensive to compute. Since regularization does not require the diagonalization of  $T$ , we may pick the less costly  $URV$  decomposition:

$$T = URV^T,$$

where  $U$  and  $V$  are orthogonal, and  $R$  is upper triangular, viz.,

$$R = \begin{pmatrix} \bar{R} & E \\ 0 & G \end{pmatrix}.$$

The three submatrices  $\bar{R}$ ,  $E$  and  $G$  ( $k \times k$ ) possess the special properties that  $\|\bar{R}\|$  is large, and  $\|E\|$  and  $\|G\|$  are small:

$$\begin{cases} \sigma_{\min}(\bar{R}) \approx \sigma_{n-k}, \\ \|E\|_F^2 + \|G\|_F^2 \approx \sigma_{n-k+1}^2 + \cdots + \sigma_n^2. \end{cases} \quad (2)$$

A regularized solution to (1) is obtained from

$$x = V \begin{pmatrix} \bar{R}^{-1} & 0 \\ 0 & 0 \end{pmatrix} U^T \hat{x}.$$

We wish to exploit the symmetry of  $T$  to save on storage and work. Start by computing a symmetric eigenvalue decomposition:

$$T = V\Lambda V^T,$$

where  $V$  is orthogonal, and  $\Lambda = \text{diag}(\lambda_i)$ , with

$$|\lambda_1| \geq \cdots \geq |\lambda_n| \geq 0.$$

Choose a tolerance  $\tau$  to find  $k$  such that

$$|\lambda_{n-k}| > \tau \geq |\lambda_{n-k+1}|,$$

i.e.,  $T$  has a numerical rank of  $n - k$ . Let

$$\hat{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_{n-k}, 0, \dots, 0).$$

The original signal is restored by the formula:

$$x = V\hat{\Lambda}^+V^T\hat{x}.$$

where  $\hat{\Lambda}^+ = \text{diag}(\lambda^+)$ .

We propose to generalize the  $URV$  decomposition as follows:

$$T = VSV^T, \quad (3)$$

where  $V$  is orthogonal and  $S$  symmetric. We shall call (3) a  $VSV$  decomposition. Partition  $S$ :

$$S \equiv \begin{pmatrix} \bar{S} & E \\ E^T & G \end{pmatrix}, \quad (4)$$

where the three submatrices  $\bar{S}$ ,  $E$  and  $G$  ( $k \times k$ ) possess similar norm properties as in (2):

$$\begin{cases} \lambda_{\min}(\bar{S}) \approx \lambda_{n-k}, \\ \|E\|_F^2 + \|\text{triu}(G)\|_F^2 \approx \lambda_{n-k+1}^2 + \cdots + \lambda_n^2. \end{cases} \quad (5)$$

We use  $\text{triu}(G)$  to denote the upper triangular part of  $G$ . The signal vector  $x$  can be restored by a truncated  $VSV$  decomposition:

$$x = V \begin{pmatrix} \bar{S}^{-1} & 0 \\ 0 & 0 \end{pmatrix} V^T \hat{x}.$$

Note that (5) is not an obvious extension of (2) to the symmetric case. Why do we use  $\text{triu}(G)$ ? Roughly speaking, since  $E$ ,  $E^T$  and  $G$  in (4) are small and  $G$  is symmetric, it is reasonable to exclude the redundant data and consider only  $E$  and  $\text{triu}(G)$ . A more rigorous argument will be given at the end of Section 5.

### 3. Fast Triangular Factorization

In this section we discuss a fast  $O(n^2)$  triangularization of an  $n \times n$  symmetric Toeplitz matrix.

First, assume that  $t_1 \neq 0$  and  $t_1 > 0$ ; otherwise consider  $-T$ . Use a displacement representation for  $T$  [2]:

$$T = R_1^T R_1 - R_2^T R_2, \quad (6)$$

where  $R_1$  and  $R_2$  are Toeplitz matrices:

$$R_1 = \frac{1}{\sqrt{t_1}} \begin{pmatrix} t_1 & t_2 & t_3 & \dots & t_{n-1} & t_n \\ 0 & t_1 & t_2 & \dots & t_{n-2} & t_{n-1} \\ 0 & 0 & t_1 & \dots & t_{n-3} & t_{n-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & t_1 & t_2 \\ 0 & 0 & 0 & \dots & 0 & t_1 \end{pmatrix}$$

and

$$R_2 = \frac{1}{\sqrt{t_1}} \begin{pmatrix} 0 & t_2 & t_3 & \dots & t_{n-1} & t_n \\ 0 & 0 & t_2 & \dots & t_{n-2} & t_{n-1} \\ 0 & 0 & 0 & \dots & t_{n-3} & t_{n-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & 0 & t_2 \\ 0 & 0 & 0 & \dots & 0 & 0 \end{pmatrix}.$$

Rewrite (6) to get

$$T = (R_1^T \ R_2^T) \begin{pmatrix} I & 0 \\ 0 & -I \end{pmatrix} \begin{pmatrix} R_1 \\ R_2 \end{pmatrix},$$

and apply  $2 \times 2$  real transformations (details in Section 4) to eliminate  $R_2$ :

$$T = (R^T \ 0) \begin{pmatrix} D_1 & 0 \\ 0 & D_2 \end{pmatrix} \begin{pmatrix} R \\ 0 \end{pmatrix} = R^T D R, \quad (7)$$

where  $R$  is upper triangular, and  $D_1$  and  $D_2$  are signature matrices (diagonal matrices with  $\pm 1$  on the diagonal). We take care to eliminate  $R_2$  in a special order, so as to maintain the Toeplitz structures of  $R_1$  and  $R_2$ . Rotating the second row of  $R_1$  against the first row of  $R_2$ , we zero out the (1,2)-entry of  $R_2$ . But a whole super-diagonal of  $R_2$  can be annihilated by applying this same rotation to the  $(i+1)$ -st row of  $R_1$  and the  $i$ -th row of  $R_2$ , for  $i = 1, \dots, n-1$ . The key is that we store and operate upon  $R_1$  and  $R_2$  as vectors, in view of

their Toeplitz structures. The calculation of the  $RDR$  decomposition defined by (7) requires only  $O(n^2)$  flops.

It is well known that the  $RDR$  decomposition without pivoting is numerically unstable. Let

$$T = \begin{pmatrix} 1.00 & 0.999 & -0.602 \\ 0.999 & 1.00 & 0.999 \\ -0.602 & 0.999 & 1.00 \end{pmatrix}.$$

Using three-decimal-digit arithmetic with rounding, the procedure computes  $D = \text{diag}(1, 1, -1)$ ,

$$R = \begin{pmatrix} 1.00 & 0.999 & -0.602 \\ 0 & 0.0447 & 35.8 \\ 0 & 0 & 35.8 \end{pmatrix},$$

and

$$T = R^T D R + \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0.638 \end{pmatrix}.$$

The matrix  $T$  has a small condition number of 2.88; its eigenvalues are 2.14, 1.60, and  $-0.746$ . The problem is that its leading  $2 \times 2$  principal submatrix is ill-conditioned. Pivoting can improve the stability, but it destroys the Toeplitz and symmetry structure, resulting in a slow algorithm. Since we are primarily interested in separating small eigenvalues from large ones, we can apply a moderately small shift to  $T$  to improve the numerical stability. For the above example, we may apply the procedure to a shifted  $T$ , e.g.,  $T + 0.1I$ . Then

$$\hat{R} = \begin{pmatrix} 1.05 & 0.951 & -0.573 \\ 0 & 0.444 & 3.47 \\ 0 & 0 & 3.36 \end{pmatrix}$$

and

$$T + 0.1I = \hat{R}^T D \hat{R} + \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0.004 \\ 0 & 0.004 & 0.02 \end{pmatrix}.$$

So the factorization error is much smaller.

Second, when  $t_1 = 0$ , the  $RDR$  decomposition does not exist. We may use the above shifting technique to overcome the difficulty. However, the choice of an appropriate shift can be a delicate matter.

#### 4. Avoid Complex Arithmetic

Since the given matrix  $T$  is real and the resultant matrices  $R$  and  $D$  are real, we want to restrict the computation to real arithmetic. In this section, we present details on how to construct a sequence of  $2 \times 2$  real transformations  $Y$  to eliminate  $R_2$ .

Given two real quantities  $\alpha$  and  $\beta$ , and two scalars  $d_1 = \pm 1$  and  $d_2 = \pm 1$ , consider the problem of finding a real transformation  $Y$  so that

$$Y \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \gamma \\ 0 \end{pmatrix} \quad (8)$$

and

$$Y^{-T} \begin{pmatrix} d_1 & 0 \\ 0 & d_2 \end{pmatrix} Y^{-1} = \begin{pmatrix} \hat{d}_1 & 0 \\ 0 & \hat{d}_2 \end{pmatrix}, \quad (9)$$

where  $\hat{d}_1 = \pm 1$  and  $\hat{d}_2 = \pm 1$ . Then we get

$$\begin{aligned} (\alpha \ \beta) Y^T Y^{-T} \begin{pmatrix} d_1 & 0 \\ 0 & d_2 \end{pmatrix} Y^{-1} Y \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \\ = (\gamma \ 0) \begin{pmatrix} \hat{d}_1 & 0 \\ 0 & \hat{d}_2 \end{pmatrix} \begin{pmatrix} \gamma \\ 0 \end{pmatrix}. \end{aligned}$$

For simplicity, we will assume that

$$d_1 = 1.$$

If needed, we can factor out a minus sign to force  $d_1 = 1$ .

1. **Q-Rotate** ( $\alpha, \beta, d_2$ )
2. check that  $\beta \neq 0$  and  $|\alpha| \neq -d_2|\beta|$ ;
3. if  $|\alpha| > |\beta|$  then
4.  $h := 1 + d_2(\beta/\alpha)^2$ ;
5.  $r := \sqrt{h} \cdot |\alpha|$ ;
6. else
7.  $h := (\alpha/\beta)^2 + d_2$ ;
8.  $r := \sqrt{|h|} \cdot |\beta|$ ;
9.  $c := \alpha/r$ ;  $s := \beta/r$ ;
10. if  $h > 0$ ;
11.  $\gamma := r$ ;  $\hat{d}_1 := 1$ ;  $\hat{d}_2 := d_2$ ;
12. else
13.  $\gamma := -r$ ;  $\hat{d}_1 := -1$ ;  $\hat{d}_2 := -d_2$ ;
14. return( $\gamma, \hat{d}_1, \hat{d}_2, c, s$ ).
15. End Q-Rotate

*Fig. 1.* Generate  $Y$  for relations (8) and (9).

To zero out  $\beta$  in (8), we choose

$$Z = \begin{pmatrix} c & d_2 s \\ -s & c \end{pmatrix}, \quad (10)$$

with  $c = \alpha/r$  and  $s = \beta/r$ , where

$$r = \sqrt{\alpha^2 + d_2 \beta^2}.$$

So  $c^2 + d_2 s^2 = 1$ , and

$$Z^{-1} = \begin{pmatrix} c & -d_2 s \\ s & c \end{pmatrix}.$$

We see that (9) is satisfied:

$$Z^{-T} \begin{pmatrix} 1 & 0 \\ 0 & d_2 \end{pmatrix} Z^{-1} = \begin{pmatrix} 1 & 0 \\ 0 & d_2 \end{pmatrix}.$$

When  $d_2 = 1$ , we get a circular (also known as Givens) rotation  $Z$ , where

$$Z = \begin{pmatrix} c & s \\ -s & c \end{pmatrix}, \quad Z^{-1} = Z^T,$$

and  $c^2 + s^2 = 1$ . When  $d_2 = -1$ , we get a hyperbolic rotation  $Z$  satisfying

$$Z = \begin{pmatrix} c & -s \\ -s & c \end{pmatrix}, \quad Z^{-1} = \begin{pmatrix} c & s \\ s & c \end{pmatrix},$$

and  $c^2 - s^2 = 1$ . Thus, our new transformation  $Z$  unifies the circular and hyperbolic rotations. We refer to  $Z$  as a quadratic rotation.

When  $Z$  is real, viz.,  $\alpha^2 + d_2 \beta^2 > 0$ , the choice is clear:

$$Y = Z.$$

But  $Z$  is complex when  $d_2 = -1$  and  $|\alpha| < |\beta|$ , in which case we calculate

$$\tilde{r} = \sqrt{-(\alpha^2 + d_2 \beta^2)} > 0,$$

as well as  $\tilde{c} = \alpha/\tilde{r}$  and  $\tilde{s} = \beta/\tilde{r}$ . Defining

$$Y = \begin{pmatrix} \tilde{c} & d_2 \tilde{s} \\ -\tilde{s} & \tilde{c} \end{pmatrix},$$

we see that  $Y$  is real and satisfies (8). So  $\tilde{r} = ir$ , where  $i = \sqrt{-1}$ . Thus,

$$Y = -iZ,$$

and  $\hat{d}_1 = -1$ ,  $\hat{d}_2 = -d_2$ .

We summarize our results on quadratic rotations in a procedure called *Q-Rotate* in Figure 1.

## 5. Reveal Rank

Suppose that we have decomposed  $T$  into the triangular factorization of (7):

$$T = R^T D R.$$

In this section, we show how to use it to compute a rank-revealing factorization.

We start by finding a normalized vector  $z$  that approximates the eigenvector  $z_n$  corresponding to the smallest (in magnitude) eigenvalue  $\lambda_n$ . Then

$$z = z_n + u,$$

where  $u$  denotes an error vector satisfying

$$\|u\|_2 < \epsilon$$

for some small quantity  $\epsilon$ . Use the technique in [7] to find circular rotations (call the product  $V^T$ ) that transform  $z$  into the  $n$ -th unit vector  $e_n$ :

$$V^T z = e_n.$$

Apply the transpose of these rotations from the right on  $R$ . However, when a circular rotation is applied to the  $i$ -th and  $(i+1)$ -st columns of  $R$ , it creates a nonzero  $(i+1, i)$  subdiagonal entry in  $R$ . To restore the triangular structure of  $R$ , apply a quadratic rotation from the left to annihilate the newly created nonzero element. We have

$$T = V S V^T,$$

where

$$S = \bar{R}^T \bar{D} \bar{R},$$

with  $\bar{R}$  and  $\bar{D}$  denoting, respectively, the resultant triangular and signature matrices. To reveal rank, we partition

$$S = \begin{pmatrix} \bar{S} & y \\ y^T & \lambda \end{pmatrix} \quad (11)$$

where  $\bar{S}$  is  $(n-1) \times (n-1)$ ; we also let

$$V = (V_1 \ z),$$

where  $V_1$  is  $n \times (n-1)$ . Now, we show that the matrix  $S$  satisfies the norm properties (5) with  $k = 1$ . First, since  $\lambda = e_n^T S e_n$ , we get

$$|\lambda_n| \leq |\lambda|.$$

Also, since  $S = V^T T V$  and  $V e_n = z$ , we have

$$\lambda = z^T T z = \lambda_n + 2\lambda_n u^T z_n + u^T T u,$$

and so

$$|\lambda| \leq |\lambda_n| + 2|\lambda_n| \epsilon + \|T\|_2 \epsilon^2.$$

Second, check the vector:

$$y = V_1^T T z = V_1^T T (z_n + u) = \lambda_n V_1^T z_n + V_1^T T u.$$

Since  $V_1^T z = 0$ , we have

$$V_1^T z_n = -V_1^T u.$$

It follows that

$$y = V_1^T (T - \lambda_n I) u$$

and

$$\|y\|_2 \leq \|V_1^T (T - \lambda_n I)\|_2 \epsilon \leq \|T - \lambda_n I\|_2 \epsilon.$$

Consequently, when  $z \approx z_n$ , we get

$$\|y\|_2^2 + \lambda^2 \approx \lambda_n^2.$$

Third, check  $\lambda_{\min}(\bar{S})$ . Let

$$\bar{S} = \tilde{V} \tilde{D} \tilde{V}^T$$

be an eigenvalue decomposition of  $\bar{S}$ . Then

$$T = V S V^T = V \begin{pmatrix} \bar{S} & y \\ y^T & \lambda \end{pmatrix} V^T$$

and so

$$T = V \begin{pmatrix} \tilde{V} & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} \tilde{D} & \tilde{V}^T y \\ y^T \tilde{V} & \lambda \end{pmatrix} \begin{pmatrix} \tilde{V}^T & 0 \\ 0 & 1 \end{pmatrix} V^T.$$

Gerschgorin theorem states that if  $\epsilon$  is sufficiently small and if the eigenvalues of  $T$  are distinct, then

$$|\lambda_{\min}(\bar{S}) - \lambda_{n-1}| \leq \|\tilde{V}^T y\|_2 \leq \|T - \lambda_n I\|_2 \epsilon,$$

showing that

$$\lambda_{\min}(\bar{S}) \approx \lambda_{n-1}.$$

In summary, if  $z \approx z_n$  and the eigenvalues of  $T$  are distinct, then

$$\begin{cases} \|y\|_2^2 + \lambda^2 \approx \lambda_n^2, \\ \lambda_{\min}(\bar{S}) \approx \lambda_{n-1}, \end{cases}$$

which are simply the properties (5) when  $k = 1$ .

If the estimated smallest eigenvalue is less than the preset tolerance, we deflate  $T$  and repeat the procedure on  $\bar{S}$  in (11). As before, we get

$$\bar{S} = \tilde{V} \begin{pmatrix} \tilde{S} & \tilde{y} \\ \tilde{y}^T & \tilde{\lambda} \end{pmatrix} \tilde{V}^T \quad (12)$$

and

$$\begin{cases} \|\tilde{y}\|_2^2 + \tilde{\lambda}^2 \approx \lambda_{n-1}^2, \\ \lambda_{\min}(\tilde{S}) \approx \lambda_{n-2}. \end{cases}$$

Now,

$$\|y\|_2^2 + \|\tilde{y}\|_2^2 + \lambda^2 + \tilde{\lambda}^2 \approx \lambda_{n-1}^2 + \lambda_n^2.$$

Combining (11) and (12), we get

$$T = \hat{V} \begin{pmatrix} \tilde{S} & E \\ E^T & G \end{pmatrix} \hat{V}^T$$

where  $G$  is  $2 \times 2$ . Then we have

$$\begin{cases} \|E\|_F^2 + \|\text{triu}(G)\|_F^2 \approx \lambda_{n-1}^2 + \lambda_n^2, \\ \lambda_{\min}(\tilde{S}) \approx \lambda_{n-2}, \end{cases}$$

which are precisely (5) for the case where  $k = 2$ . This also justifies the use of  $\text{triu}(G)$  in (5).

We continue the deflation procedure until the estimated eigenvalue exceeds the tolerance. Consequently, we obtain the desired  $VSV$  decomposition of (3). This factorization, including the eigenvector estimation, costs  $O(kn^2)$  flops.

## 6. Examples

We present three examples to show how our new method performs as well as the SVD approach. The three models of  $T$  are adopted from [6]. We wrote a program in MATLAB and ran it on a SUN/Sparc2000 computer. Each estimated eigenvector was computed using seven inverse power iterations [3], and the singular value decomposition was computed using the MATLAB function SVD.

In our examples, we calculate the condition number of  $T$ , viz.  $\kappa_2(T)$ , and determine its numerical rank  $n - k$  using the tolerance

$$\tau = 10^{-3}.$$

The matrices are ill-conditioned (calling for regularization), and our new technique always calculates the numerical rank correctly. After computing the rank-revealing decomposition

$$T = VSV^T,$$

where

$$S = \begin{pmatrix} \bar{S} & E \\ E^T & G \end{pmatrix},$$

we partition

$$V = (V_S, V_N),$$

where  $V_N$  is  $n \times k$ . Denote by  $x_S$  the restored signal vector computed by our new symmetric decomposition. Then

$$x_S = V_S \bar{S}^{-1} V_S^T \hat{x}.$$

Correspondingly, in the SVD

$$T = U\Sigma W^T,$$

we partition

$$W = (W_S, W_N) \text{ and } U = (U_S, U_N),$$

where  $W_N$  and  $U_N$  are  $n \times k$  matrices; similarly,

$$\Sigma = \text{diag}(\Sigma_S, \Sigma_N),$$

where

$$\begin{cases} \Sigma_S = \text{diag}(\sigma_1, \dots, \sigma_{n-k}), \\ \Sigma_N = \text{diag}(\sigma_{n-k+1}, \dots, \sigma_n). \end{cases}$$

*Fig. 2.* The 48 smallest eigenvalues in Example 3.

Since  $T$  is symmetric,  $U$  and  $W$  are the same except for signs in their columns. Let  $x_T$  denote the restored signal vector computed by the TSVD method. Then

$$x_T = W_S \Sigma_S^{-1} U_S^T \hat{x}.$$

We use the parameter

$$\epsilon = \|V_N^T W_S\|_2$$

to measure the distance between the two subspaces  $\text{Range}(V_S)$  and  $\text{Range}(W_S)$  [3].

*Example 1.* We pick a  $250 \times 250$  banded symmetric Toeplitz matrix  $T$ :

$$t_m = \begin{cases} 1, & \text{if } m = 1; \\ \frac{\sin^2((m-1)/\omega)}{((m-1)/\omega)^2}, & \text{if } 2 \leq m \leq 5; \\ 0, & \text{otherwise;} \end{cases}$$

with  $\omega = 4.0$ . The matrix is ill-conditioned:

$$\kappa_2(T) \approx 2.2 \times 10^6.$$

Its numerical rank equals 248, as detected by both methods; hence

$$k = 2,$$

a small rank deficiency. After regularization, we get a much better conditioned matrix  $\bar{S}$ :

$$\kappa_2(\bar{S}) \approx 3.5 \times 10^2.$$

In Table 1, we present the three ( $= k + 1$ ) smallest (in magnitude) eigenvalues  $\lambda_i$  computed by the MATLAB SVD routine and  $\hat{\lambda}_i$  computed by our new algorithm.

Since the eigenvalues are well isolated, the estimates are accurate. As predicted by the analysis in Section 5, the off-diagonal blocks  $E$  and  $\text{triu}(G)$  are small:

$$\begin{cases} \|E\|_F = 7.45 \times 10^{-12}, \\ \|\text{triu}(G)\|_F = 1.19 \times 10^{-4} \approx \sqrt{\lambda_{249}^2 + \lambda_{250}^2}. \end{cases}$$

Furthermore, the error  $\epsilon$  equals  $4.0 \times 10^{-9}$ , and so the subspace approximation of  $\text{Range}(W_S)$  by  $\text{Range}(V_S)$  is also very good.

*Example 2.* Consider a band Toeplitz matrix  $T$  of order 150:

$$t_m = \begin{cases} 2\omega, & \text{if } m = 1; \\ \frac{\sin(2\pi\omega(m-1))}{\pi(m-1)}, & \text{if } 2 \leq m \leq 9; \\ 0, & \text{otherwise;} \end{cases}$$

with  $\omega = 0.05$ . While the condition number of the given matrix  $T$  is large:

$$\kappa_2(T) \approx 1.2 \times 10^6,$$

the condition number of the regularized matrix  $\bar{S}$  is acceptable:

$$\kappa_2(\bar{S}) \approx 9.5 \times 10^2.$$

Both methods computed the rank deficiency as

$$k = 5.$$

Table 2 lists the six smallest (in magnitude) eigenvalues as computed by the MATLAB routine and as estimated by our procedure.

Since the three smallest eigenvalues are closely clustered, the approximate eigenvalues  $\hat{\lambda}_{148}$  and  $\hat{\lambda}_{149}$  are inaccurate and

$$\begin{cases} \|E\|_F = 1.36 \times 10^{-1}, \\ \|\text{triu}(G)\|_F = 7.46 \times 10^{-3}. \end{cases}$$

The error  $\epsilon$  equals  $1.4 \times 10^{-2}$  and so the estimation of  $\text{Range}(W_S)$  by  $\text{Range}(V_S)$  is off.

*Example 3.* We choose a  $120 \times 120$  positive definite Toeplitz matrix  $T$ :

$$t_m = \begin{cases} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(m-1)^2}{2\sigma^2}}, & \text{if } 1 \leq m \leq 8; \\ 0, & \text{otherwise;} \end{cases}$$

with  $\sigma = 2.0$ . The matrix  $T$  is ill-conditioned:

$$\kappa_2(T) \approx 1.6 \times 10^6.$$

Note that

$$k = 47$$

with both methods and

$$\kappa_2(\bar{S}) \approx 1.1 \times 10^3.$$

*Table 1.* Three smallest eigenvalues in Example 1.

$i$	248	249	250
$\lambda$	$2.26e-2$	$1.19e-4$	$3.51e-6$
$\hat{\lambda}_i$	$2.17e-2$	$1.19e-4$	$3.51e-6$

*Table 2.* Six smallest eigenvalues in Example 2.

$i$	145	146	147
$\lambda_i$	$1.19e-3$	$4.88e-4$	$1.71e-4$
$\hat{\lambda}_i$	$1.19e-3$	$4.86e-4$	$1.50e-4$
$i$	148	149	150
$\lambda_i$	$2.76e-5$	$1.00e-5$	$9.37e-6$
$\hat{\lambda}_i$	$5.83e-5$	$5.94e-5$	$9.43e-6$

Figure 2 plots the forty-eight ( $= k + 1$ ) smallest eigenvalues computed by the MATLAB SVD routine (represented by a solid line) against those esti-

mated by our symmetric rank-revealing program (represented by a dashed line).

As shown in Figure 2, the eigenvalues decrease gradually and  $\lambda_{n-k}$  is close to  $\lambda_{n-k+1}$ . Just as in Example 2, we observe sizable errors:

$$\begin{cases} \|E\|_F = 5.26 \times 10^{-2}, \\ \|\text{triu}(G)\|_F = 3.97 \times 10^{-2}. \end{cases}$$

The error  $\epsilon$  equals  $4.8 \times 10^{-1}$ . So,  $\text{Range}(V_S)$  does not estimate  $\text{Range}(W_S)$  very well.

What have we learned from these three examples? The TSVD and our method perform well when there is a significant gap between the large eigenvalues and the small ones (such as in Example 1), and may not work well otherwise. In Example 4, we show that even in the case where the eigenvalues decrease gradually (as in Examples 2 and 3) both the TSVD and our new approach still do good jobs in restoring the original signal vector.

*Fig. 3.* Original signal vector  $x$  (left) and noisy signal vector  $\hat{x}$  (right) in Example 4.

*Example 4.* We stay with the same matrix as in Example 3. In addition, we choose an original signal vector  $x \in R^{120}$  given by Nagy [6]. The noise vector  $w$  is a random vector generated in MATLAB by a normal distribution with zero mean and unit variance. The noise is scaled so that

$$\|w\|_2 / \|Tx\|_2 = 0.001.$$

We use this ratio as the tolerance  $\tau$  for the numerical rank. Figure 3 depicts the original signal vector  $x$  on the left and the noisy signal vector  $\hat{x}$  on the right, and Figure 4 presents the signal vector  $x_T$  restored by the TSVD method (left) and vector  $x_S$  restored by our method (right). The results show that the restoration capabilities of both methods are comparable.

Fig. 4. Vector  $x_T$  restored via TSVD (left) and vector  $x_S$  restored via our new scheme (right) in Example 4.

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